

Adding Delimited and Composable Control to a Production Programming Environment



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Scheme provides first-class *continuations*

Continuations are useful for implementing a
web server

Web Servlet with Continuations

```
(define (paper-search-servlet)
  (let ([terms (get-search-terms)])
    (find-paper terms)))
```

Web Servlet with Continuations

send back HTML form, then
wait for answer as new request

```
(define (paper-search-servlet)
  (let ([terms (get-search-terms)])
    (find-paper terms)))
```

Implementing a Web Server

```
(serve out1 in1)
```

Implementing a Web Server

```
(reply out1 (generate-html in1))
```

Implementing a Web Server

`(reply out1 [])`

```
graph TD; A("reply out1 []") --> B[generate-html in1];
```

`(generate-html in1)`

Implementing a Web Server

`(reply out1 [])`

A diagram illustrating the implementation of a web server. At the top, the expression `(reply out1 [])` is enclosed in an oval. A downward-pointing arrow connects this oval to a rectangular box below it. Inside the box is the expression `(call/cc (λ (esc) (dispatch in1)))`.

`(call/cc
 (λ (esc) (dispatch in1)))`

Implementing a Web Server

`esc = (reply out1 [])`

`(reply out1 [])`

`(call/cc
 (λ (esc) (dispatch in1)))`

Implementing a Web Server

esc = (reply out₁ [])

(reply out₁ [])



(dispatch in₁)

Implementing a Web Server

esc = (reply out₁ [])

(reply out₁ [])

(let ([terms (get-search-terms)])
 (find-paper terms))

Implementing a Web Server

esc = (reply out₁ [])

(reply out₁ [])

(let ([terms []])
 (find-paper terms))

(get-search-terms)

Implementing a Web Server

`esc = (reply out1 [])`

`(reply out1 [])`

`(let ([terms []])
 (find-paper terms))`

`(call/cc
 (λ (web-k)
 (esc (build-form (cont->url web-k))))))`

Implementing a Web Server

esc = (reply out₁ [])

web-k = (reply out₁ [])

(reply out₁ [])

(let ([terms []])
 (find-paper terms))

(let ([terms []])
 (find-paper terms))

```
(call/cc  
  (λ (web-k)  
    (esc (build-form (cont->url web-k))))))
```

Implementing a Web Server

esc = (reply out₁ [])

web-k = (reply out₁ [])

(reply out₁ [])

(let ([terms []])
 (find-paper terms))

(let ([terms []])
 (find-paper terms))

(esc (build-form (cont->url web-k)))

Implementing a Web Server

esc = (reply out₁ [])

(reply out₁ [])

(let ([terms []])
 (find-paper terms))

(esc <html/>)

web-k = (reply out₁ [])

(let ([terms []])
 (find-paper terms))

Implementing a Web Server

esc = (reply out₁ [])

(reply out₁ <html/>)

web-k = (reply out₁ [])

(let ([terms []])
 (find-paper terms))

Implementing a Web Server

esc = (reply out₁ [])

(reply out₂ [])

(dispatch in₂)

web-k = (reply out₁ [])

(let ([terms []])
 (find-paper terms))

Implementing a Web Server

esc = (reply out₁ [])

(reply out₂ [])

(web-k form-answers)

web-k = (reply out₁ [])

(let ([terms []])
 (find-paper terms))

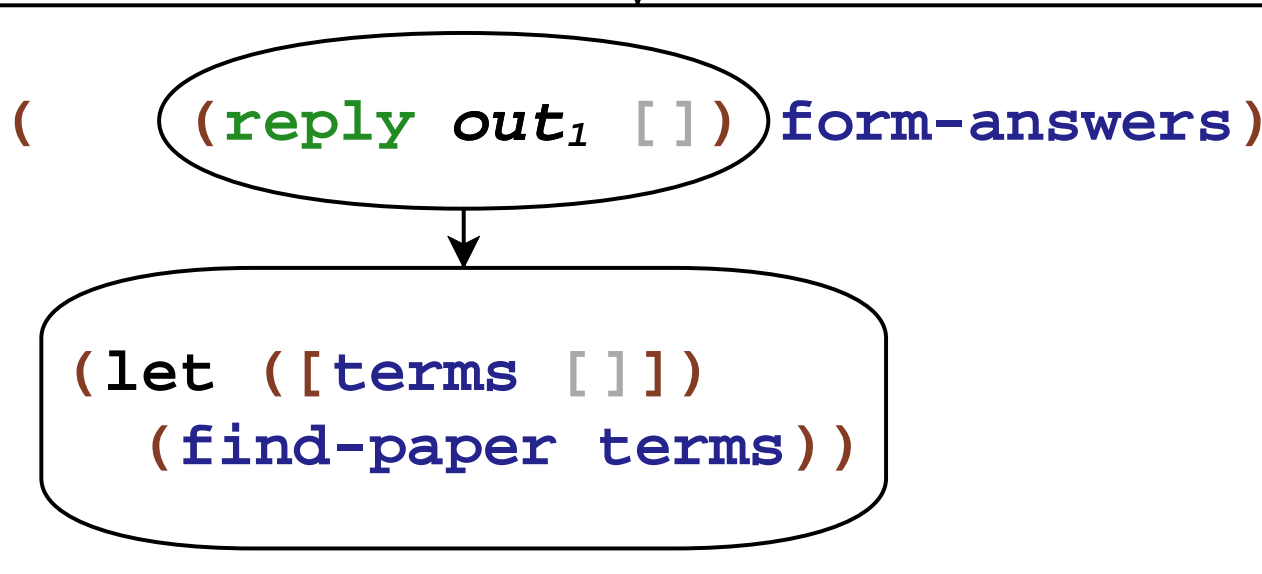
Implementing a Web Server

esc = (reply out₁ [])

web-k = (reply out₁ [])

(reply out₂ [])

(let ([terms []])
 (find-paper terms))



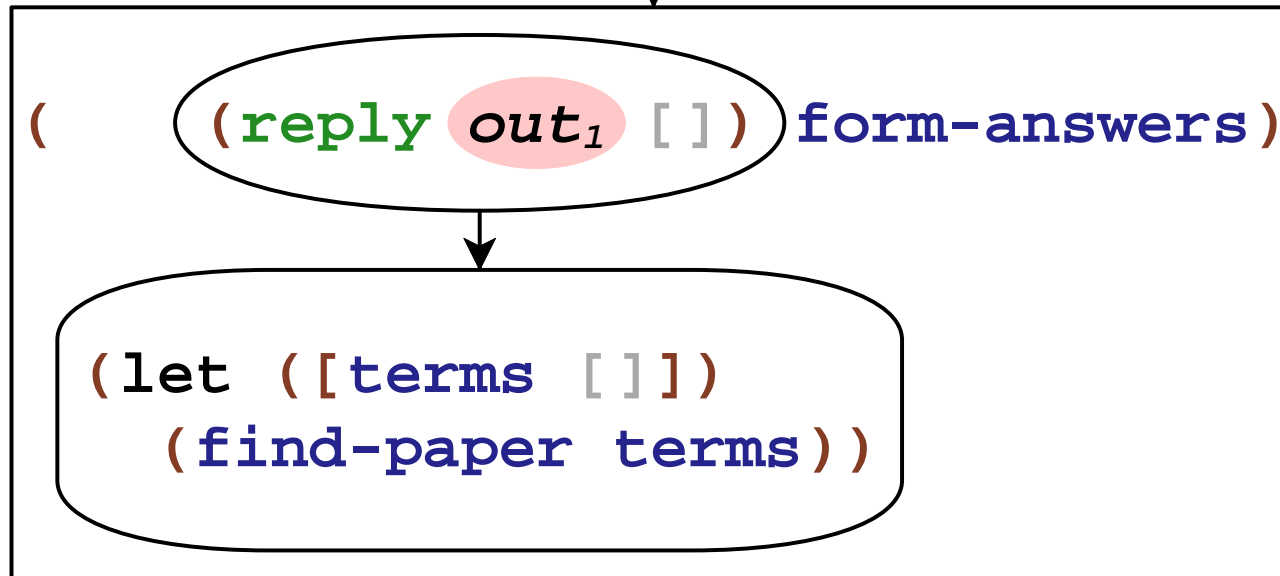
Implementing a Web Server

esc = (reply out₁ [])

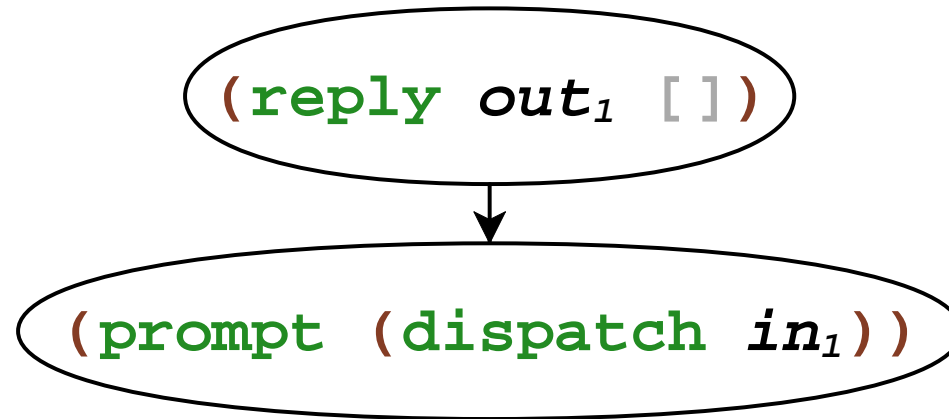
web-k = (reply out₁ [])

(reply out₂ [])

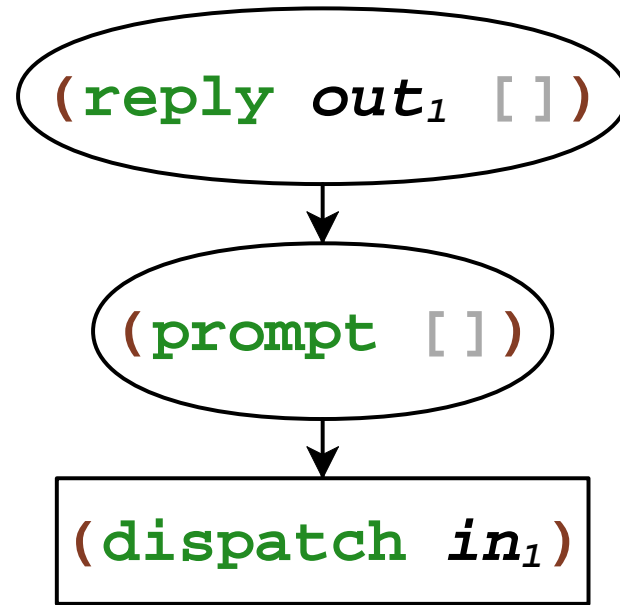
(let ([terms []])
 (find-paper terms))



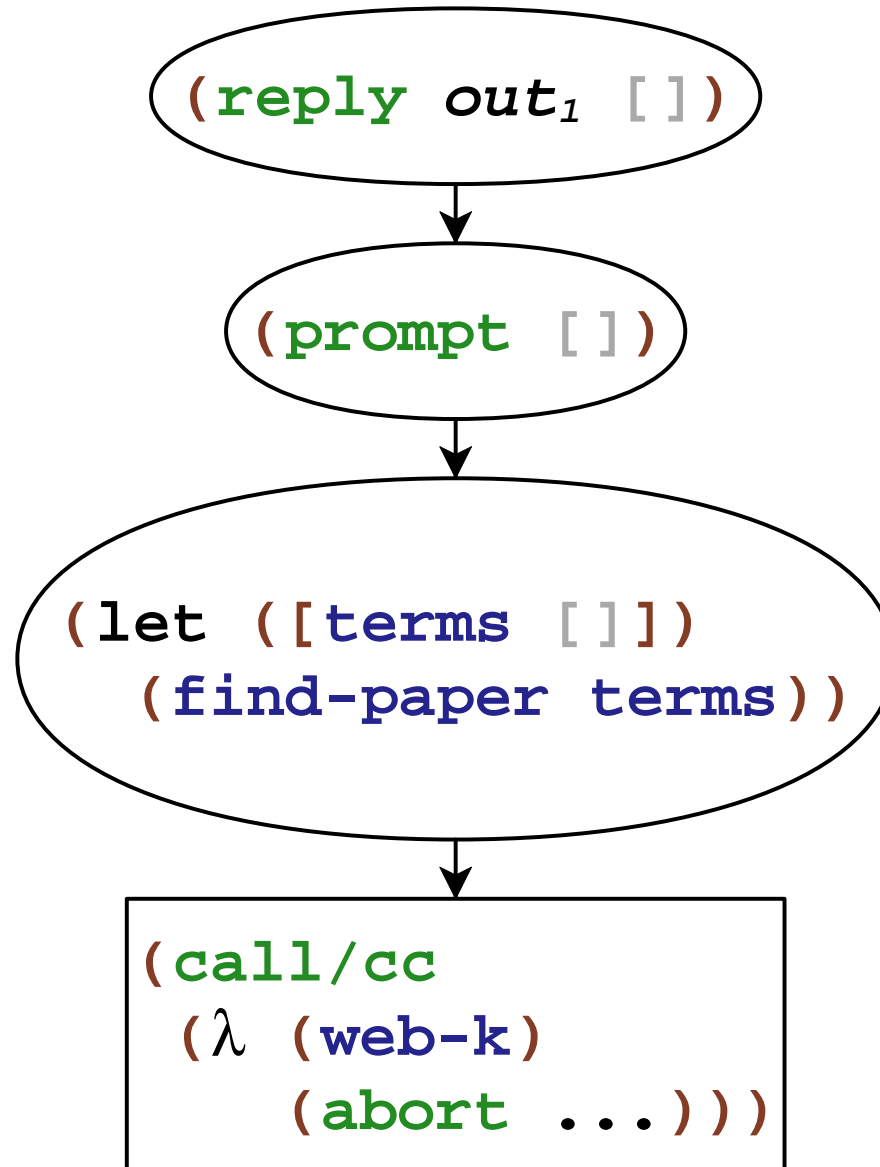
Web Server with Prompt



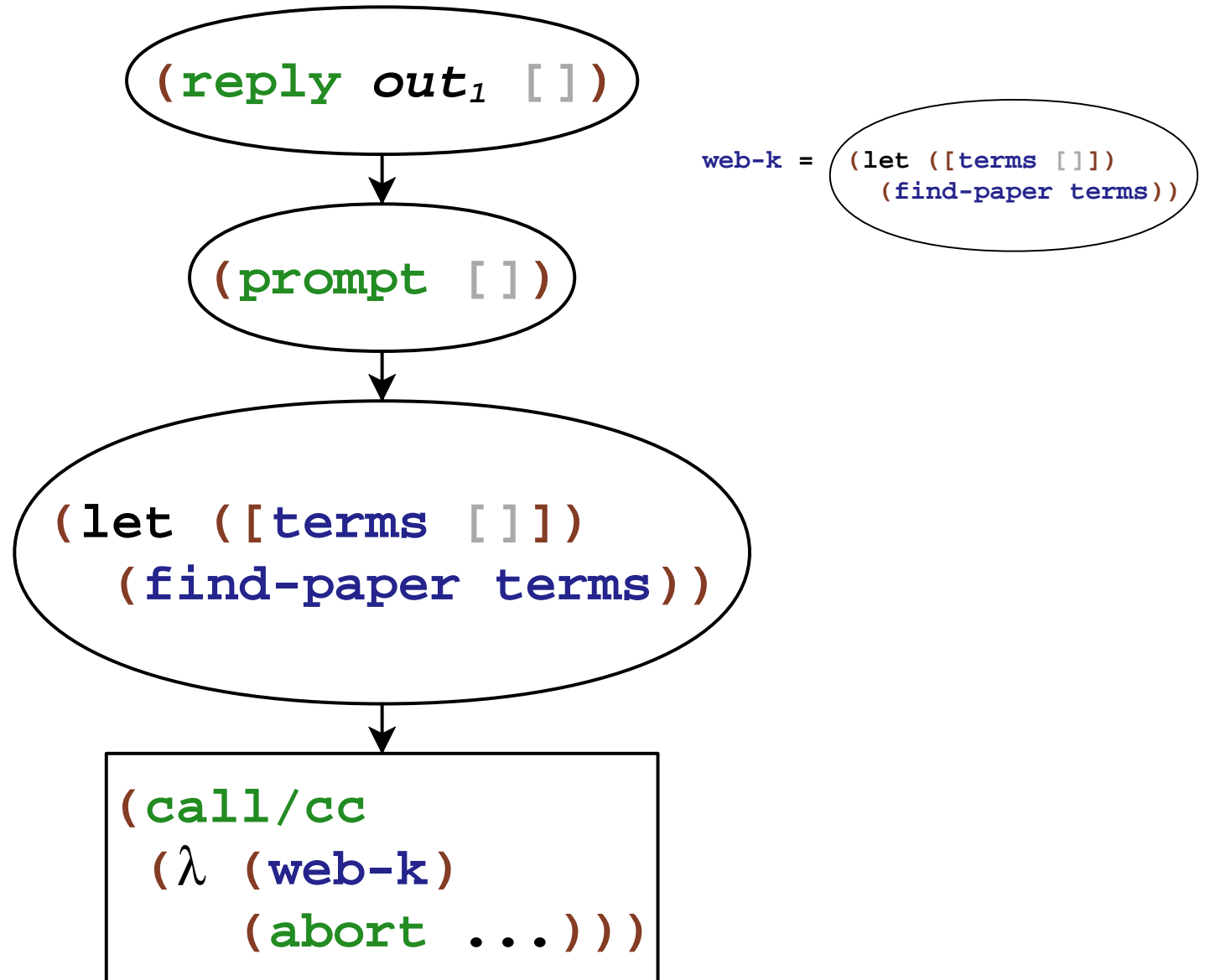
Web Server with Prompt



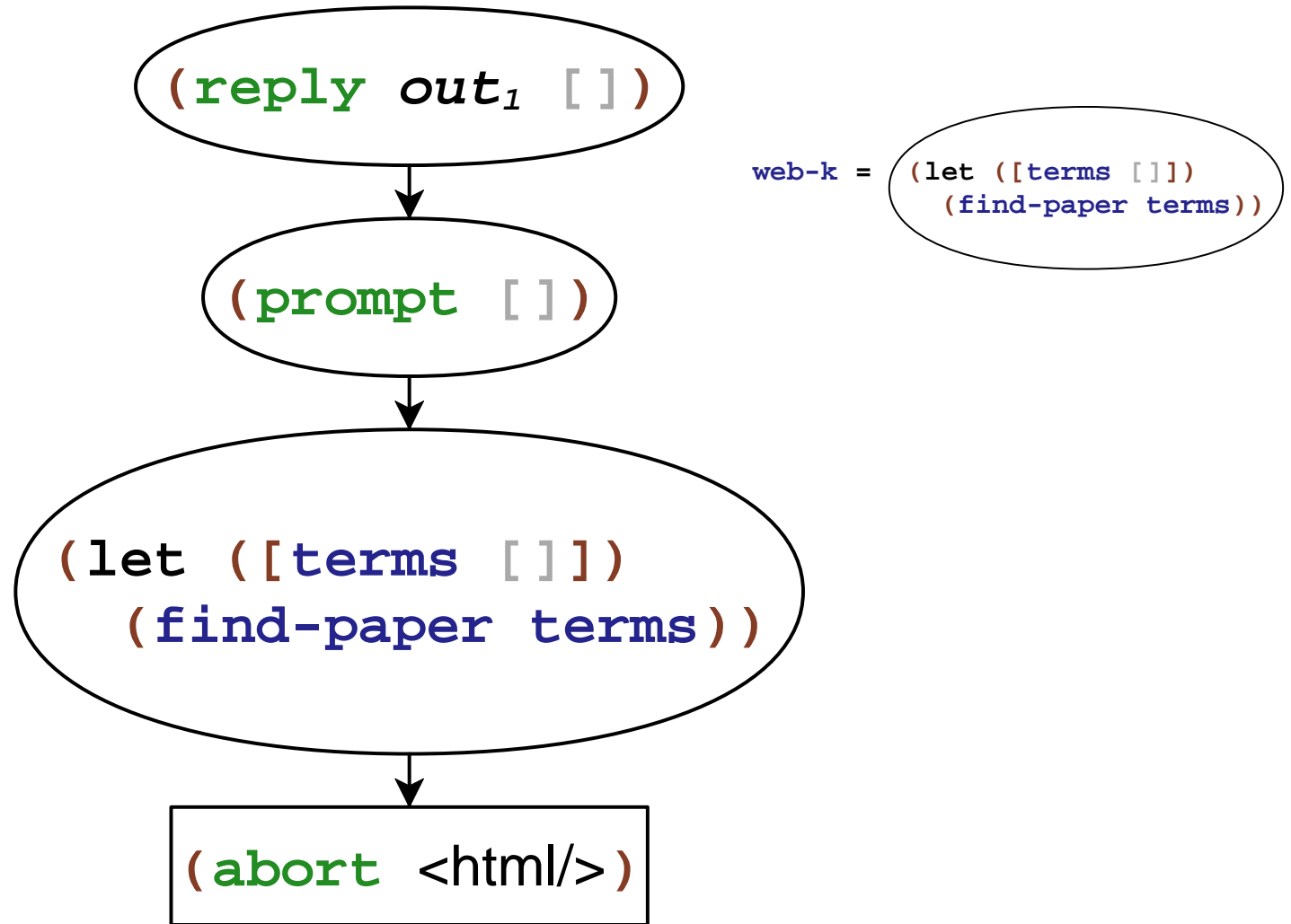
Web Server with Prompt



Web Server with Prompt



Web Server with Prompt

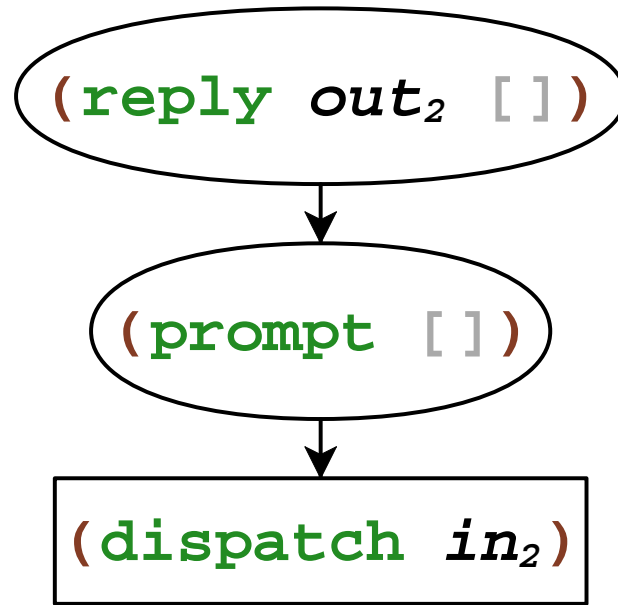


Web Server with Prompt

```
(reply out1 <html/>)
```

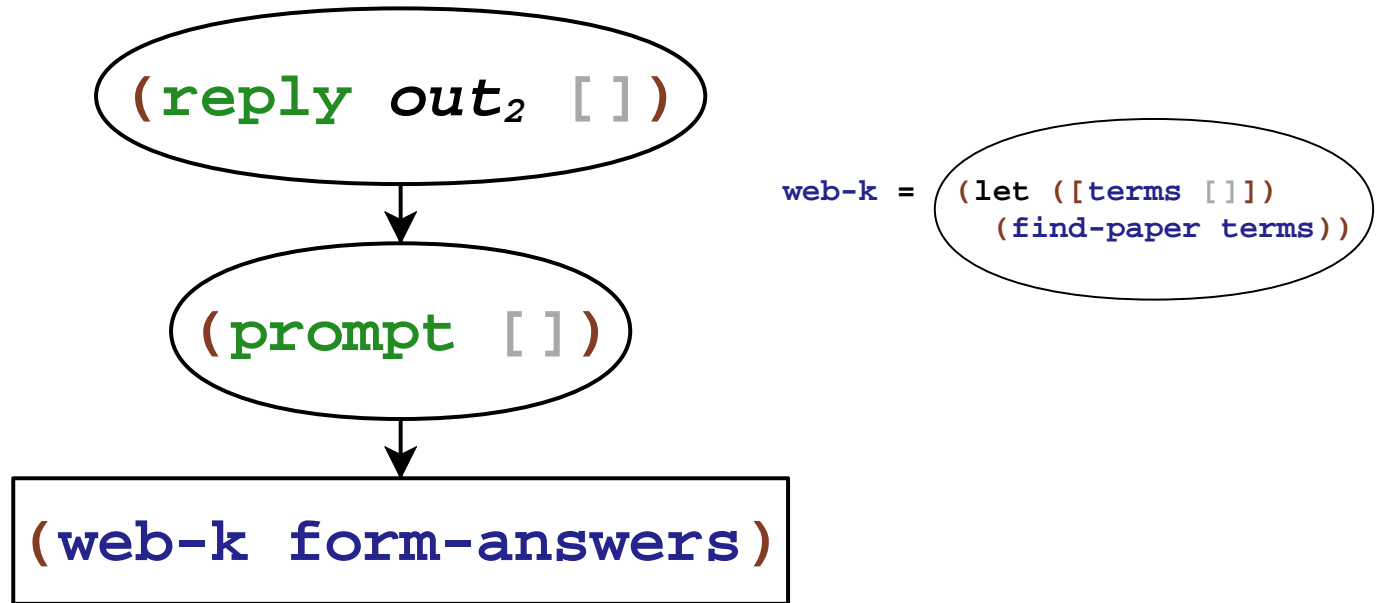
```
web-k = (let ([terms []])  
         (find-paper terms))
```

Web Server with Prompt

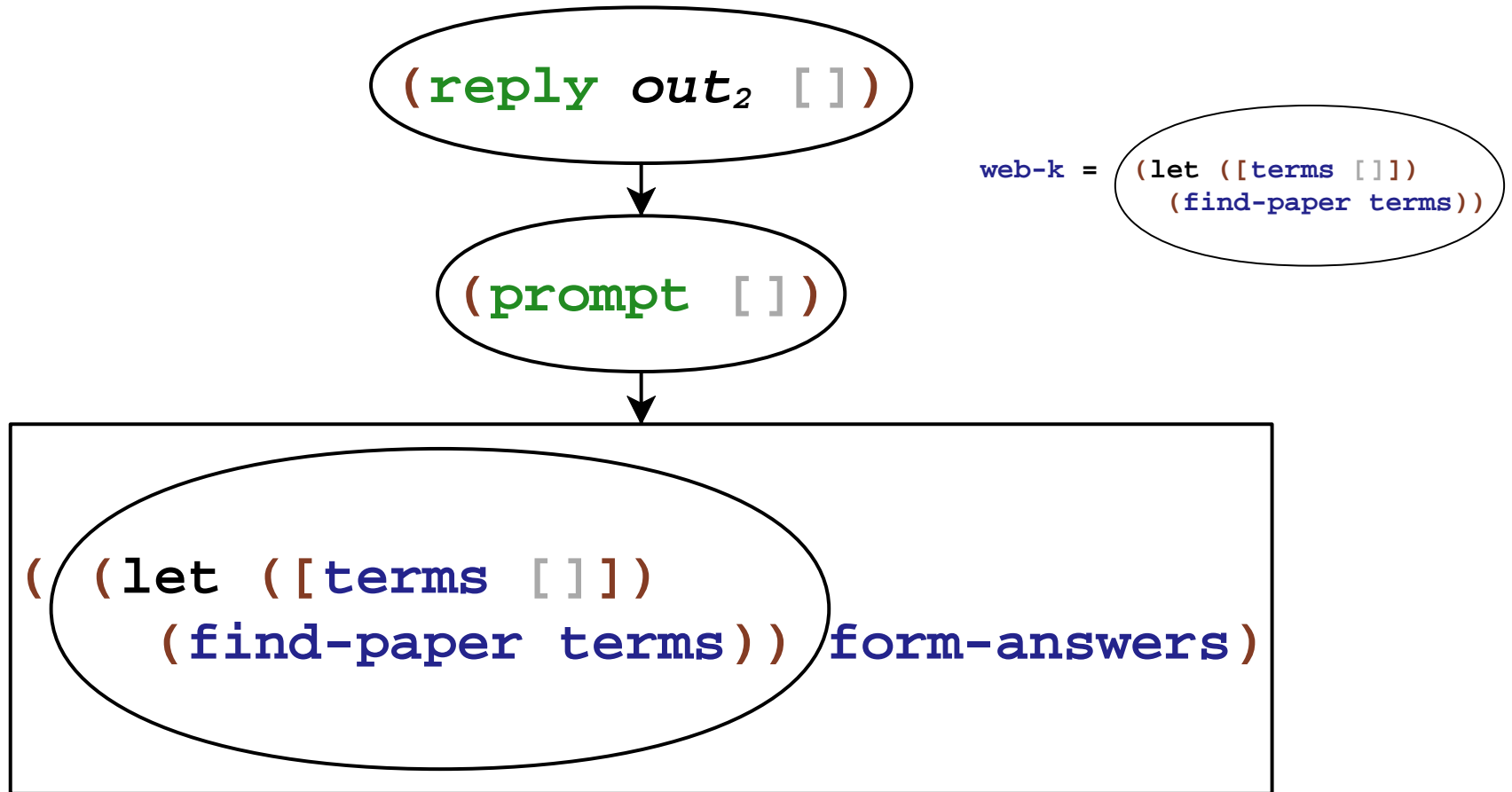


web-k = (`let` ([`terms` []])
(`find-paper terms`))

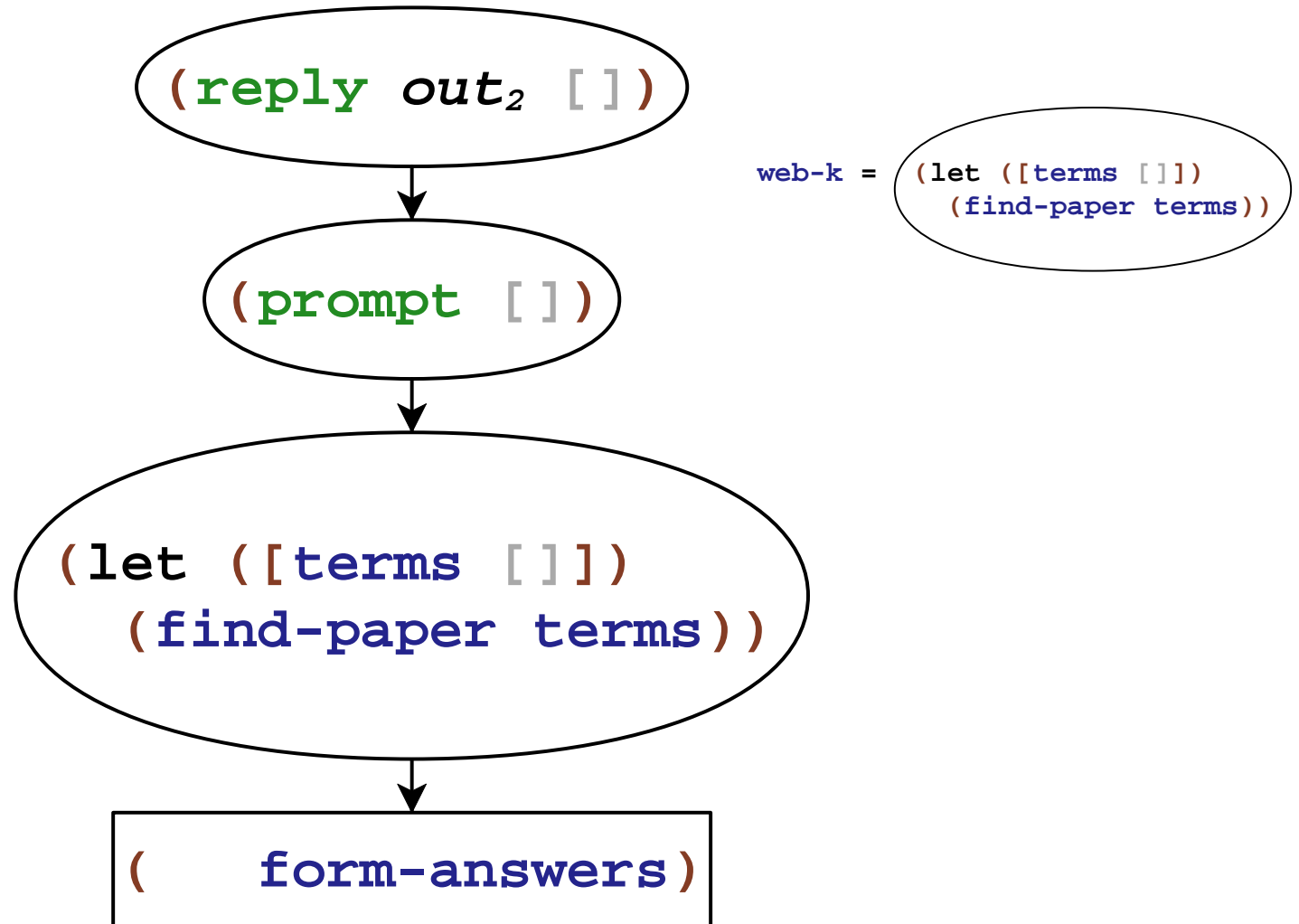
Web Server with Prompt



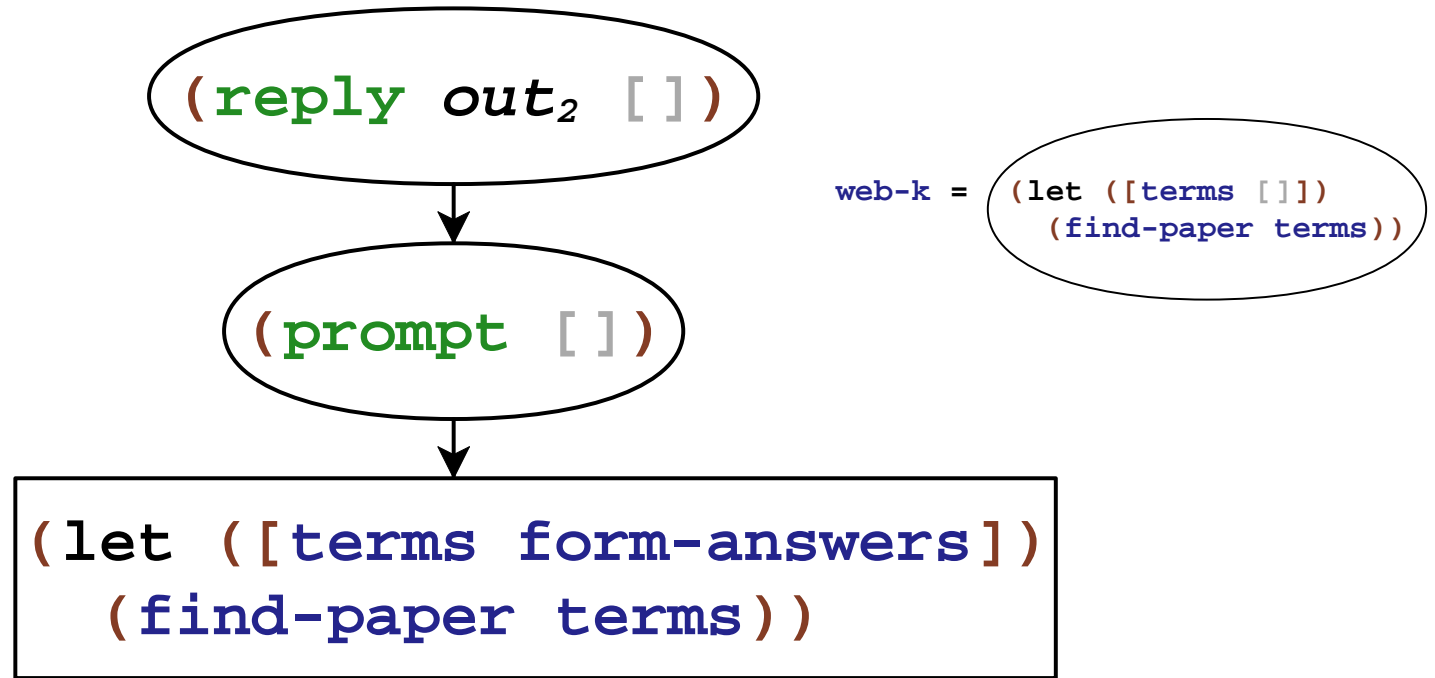
Web Server with Prompt



Web Server with Prompt



Web Server with Prompt



Rolling Your Own vs. Language Extension

Delimited control via `call/cc` doesn't work right with

- exceptions
- dynamic binding
- `dynamic-wind`

Our goal: delimited control integrated with existing constructs

Papers on Delimited Continuations

Felleisen 88

Hieb and Dybvig 90

Queinnec and Serpette 91

Queinnec 93

Wadler 94

Rehof and Sørensen 94

Gunter et al. 95

Rehof 01

Kameyama and Hasegawa 03

Shan 04

Kiselyov et al. 06

Biernacki et al. 06

Danvy and Filinski 90

Sitaram and Felleisen 90

Sitaram 93

Moreau and Queinnec 94

deGroote 94

Gunter et al. 95

Thielecke 97

Gasbichler and Sperber 02

Ariola et al. 04

Saurin 05

Dybvig et al. 06

Papers on Design

Felleisen 88

Danvy and Filinski 90

Hieb and Dybvig 90

Sitaram and Felleisen 90

Queinnec and Serpette 91

Sitaram 93

Moreau and Queinnec 94

Gunter et al. 95

Kiselyov et al. 06

Dybvig et al. 06

Papers on Implementation

Gasbichler and Sperber 02
Dybvig et al. 06

Implementations

Implementations (Now)



Balance

exceptions

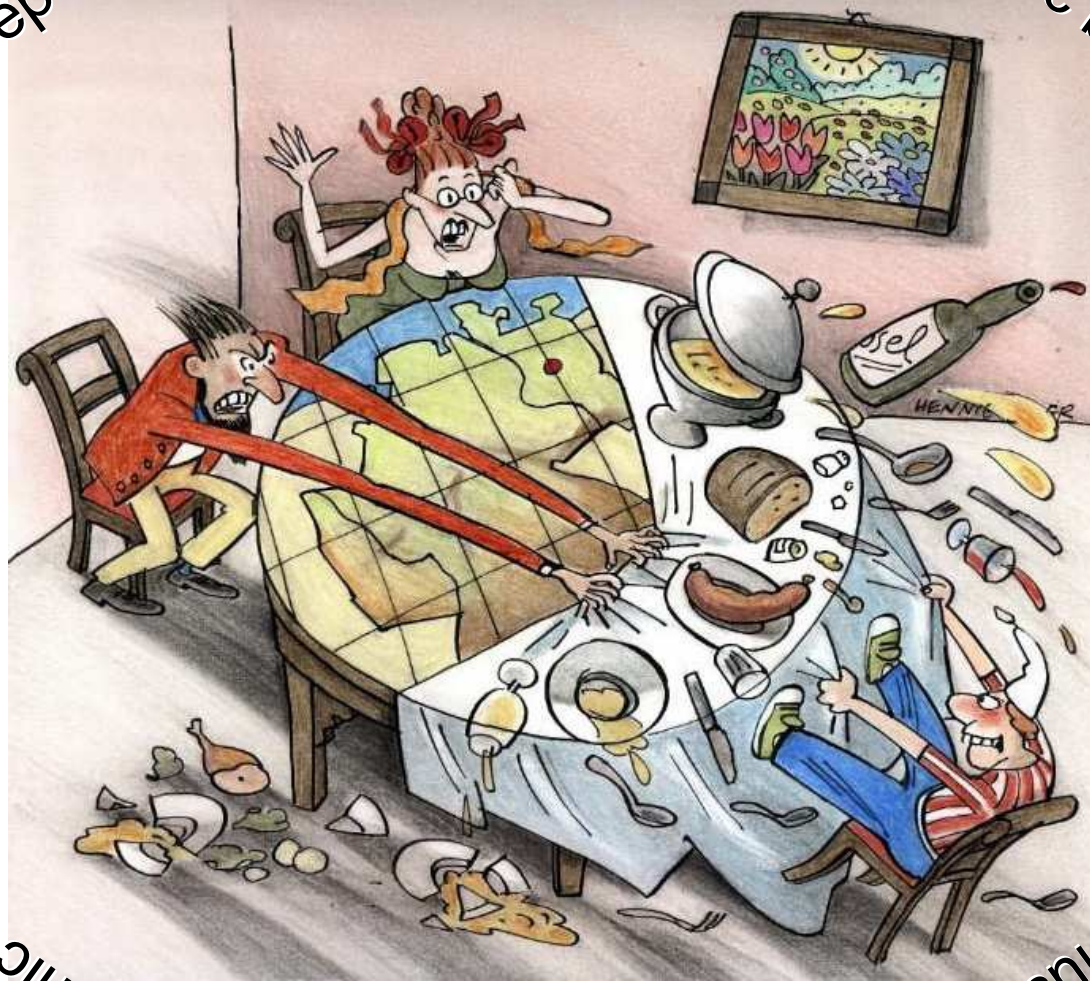
dynamic binding

prompts

abort

dynamic-wind

continuations



Contributions

- Comprehensive design
- Formal model
- Implementation

Contributions

- Comprehensive design


graphical intuition

- Formal model

- Implementation

Notation

$$(\mathbf{v}_1 \ ((\lambda \ (\mathbf{x}) \ \mathbf{x}) \ \mathbf{v}_3) \ \mathbf{v}_2))$$

Notation

$$(\mathbf{v}_1 \left((\lambda \mathbf{x} \mathbf{x}) \mathbf{v}_3 \right) \mathbf{v}_2))$$

Notation

$$(\mathbf{v}_1 \ ((\lambda \ (\mathbf{x}) \ \mathbf{x}) \ \mathbf{v}_3) \ \mathbf{v}_2))$$

$$((\lambda \ (\mathbf{x}) \ \mathbf{x}) \ \mathbf{v}_3)$$

Notation

$$(\mathbf{v}_1 \ ((\lambda \ (\mathbf{x}) \ \mathbf{x}) \ \mathbf{v}_3) \ \mathbf{v}_2))$$

$$((\lambda \ (\mathbf{x}) \ \mathbf{x}) \ \mathbf{v}_3)$$

Notation

$(\mathbf{v}_1 \ ((\lambda \ (\mathbf{x}) \ \mathbf{x}) \ \mathbf{v}_3) \ \mathbf{v}_2))$

$([\] \ \mathbf{v}_2)$

$((\lambda \ (\mathbf{x}) \ \mathbf{x}) \ \mathbf{v}_3)$

Notation

$(\mathbf{v}_1 \ ((\lambda \ (\mathbf{x}) \ \mathbf{x}) \ \mathbf{v}_3) \ \mathbf{v}_2))$

$([\] \ \mathbf{v}_2)$

$((\lambda \ (\mathbf{x}) \ \mathbf{x}) \ \mathbf{v}_3)$

Notation

$(\mathbf{v}_1 ((\lambda (\mathbf{x}) \mathbf{x}) \mathbf{v}_3) \mathbf{v}_2))$

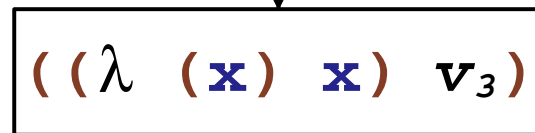
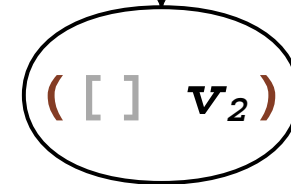
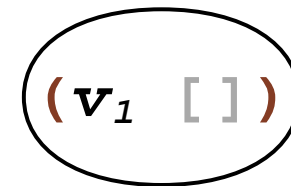
$(\mathbf{v}_1 [])$

$([] \mathbf{v}_2)$

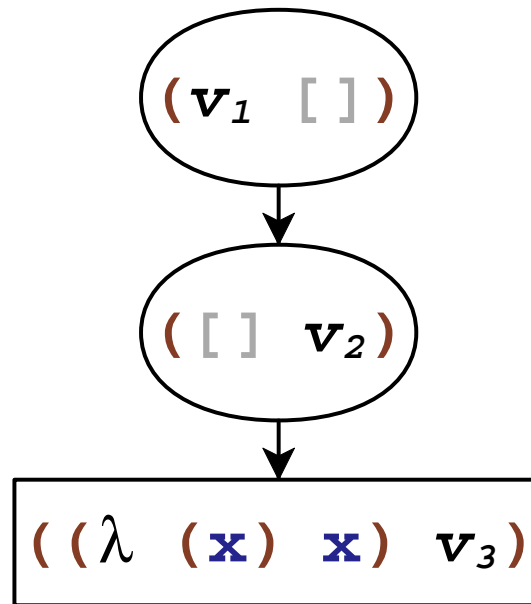
$((\lambda (\mathbf{x}) \mathbf{x}) \mathbf{v}_3)$

Notation

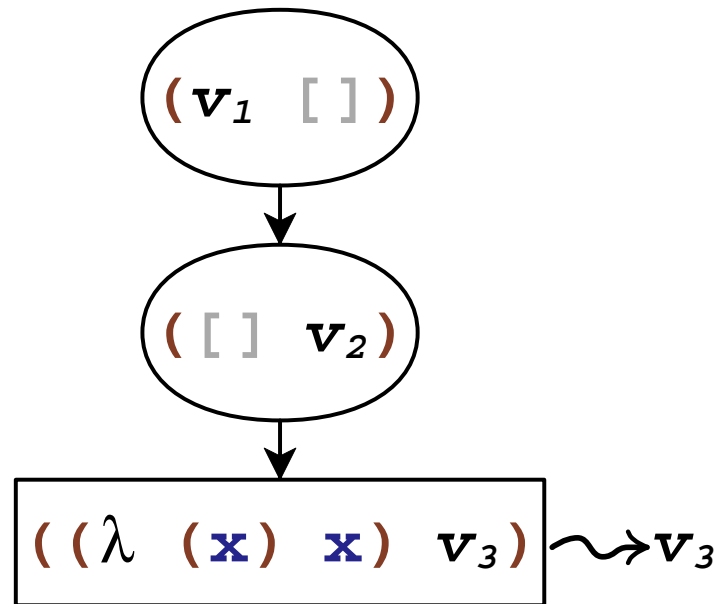
$$(\mathbf{v}_1 \ ((\lambda \ (\mathbf{x}) \ \mathbf{x}) \ \mathbf{v}_3) \ \mathbf{v}_2)) =$$



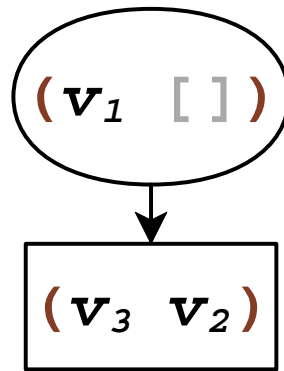
Reductions



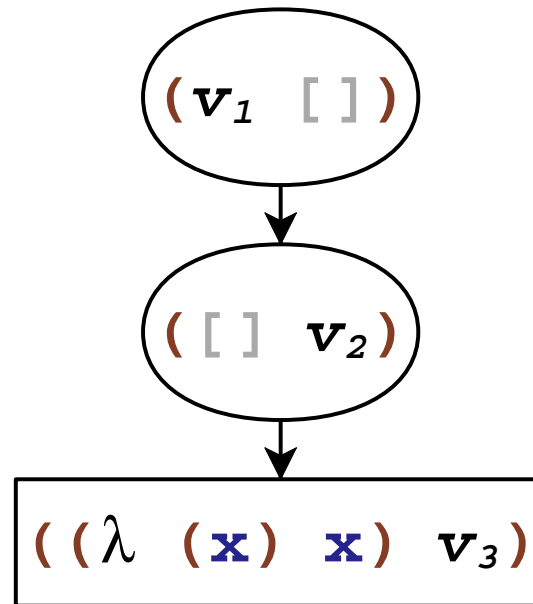
Reductions



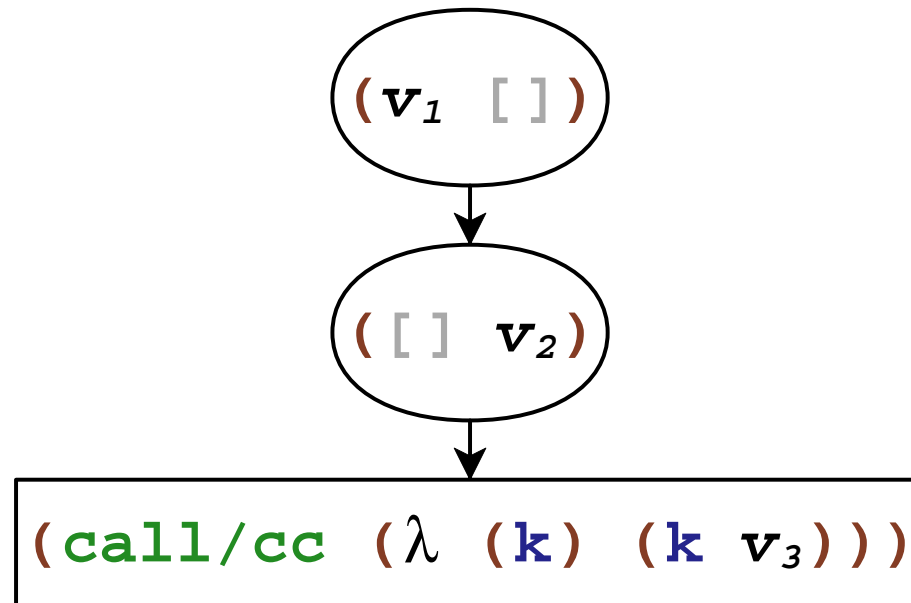
Reductions



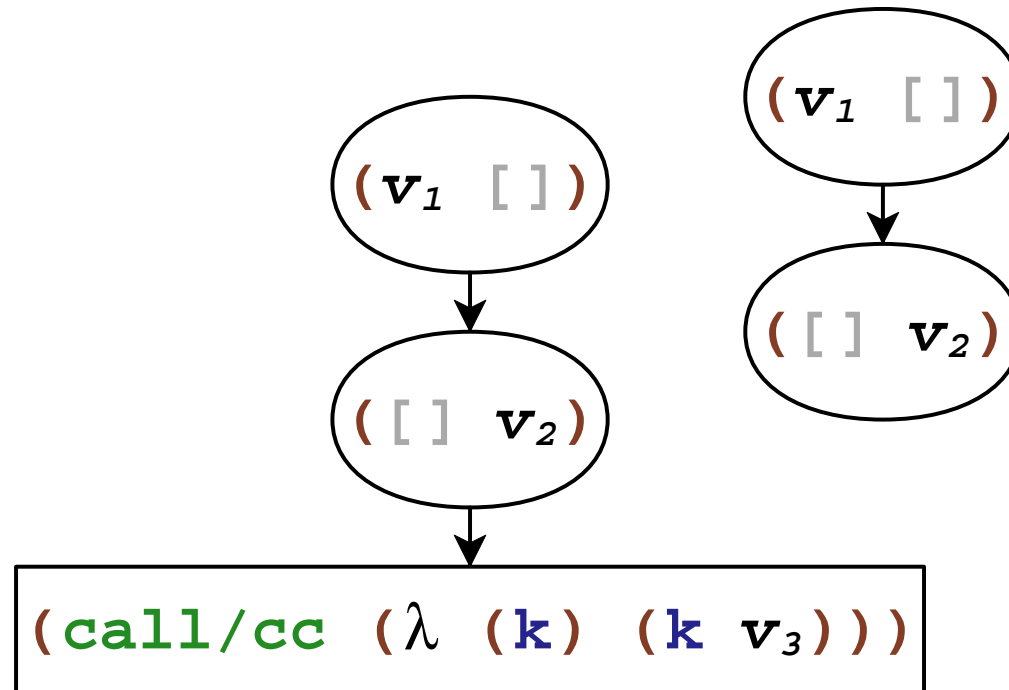
Reductions



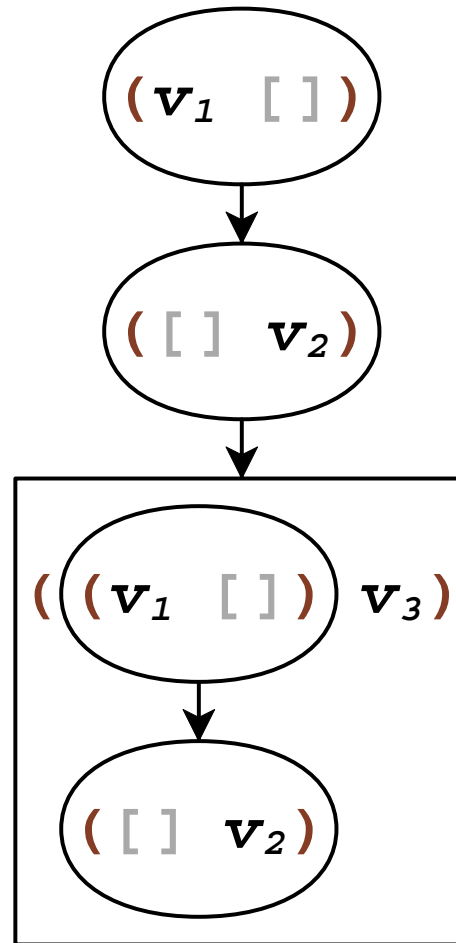
Continuations



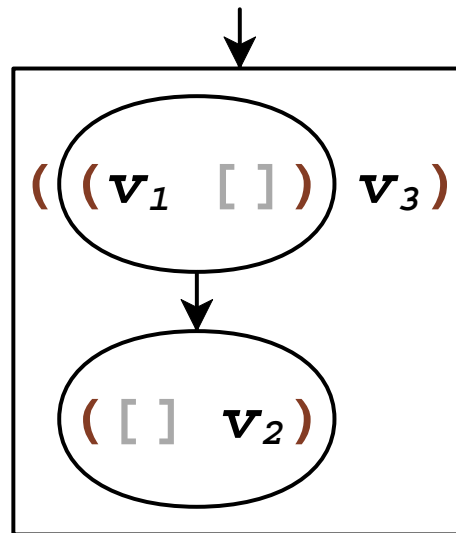
Continuations



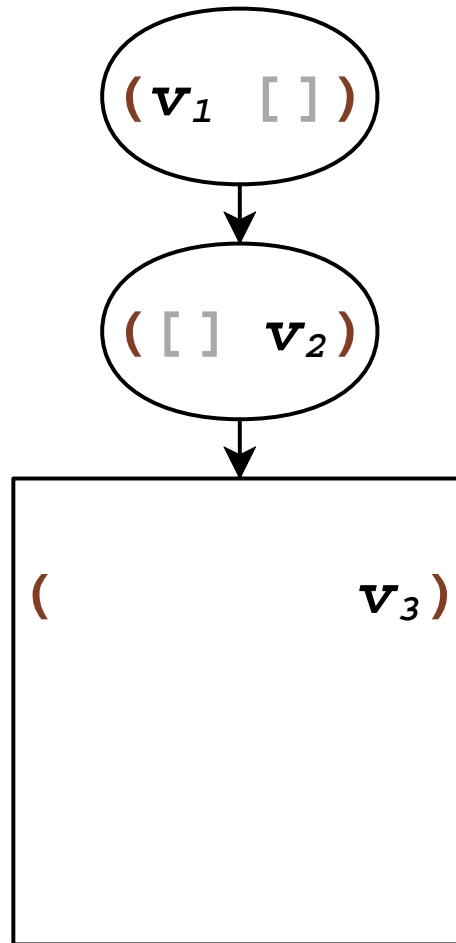
Continuations



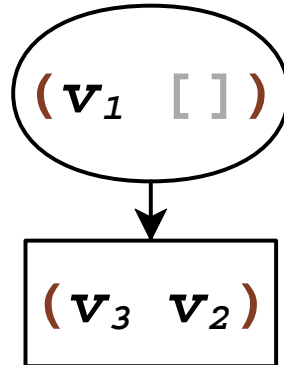
Continuations



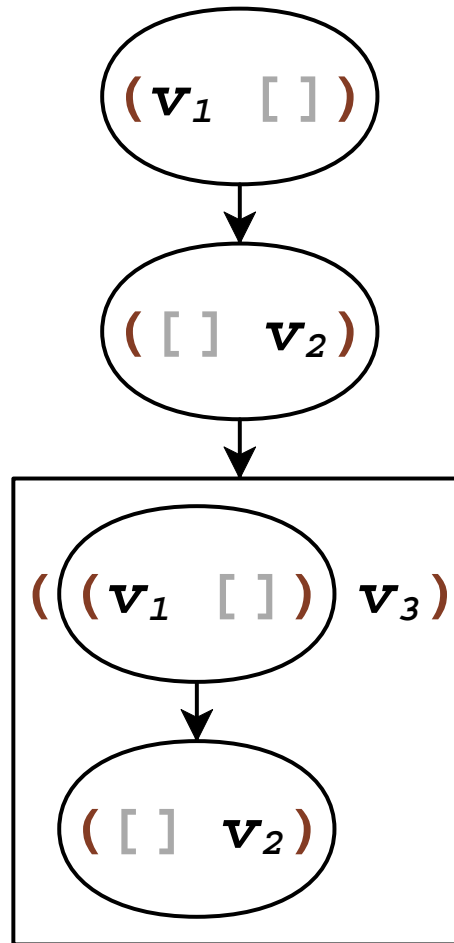
Continuations



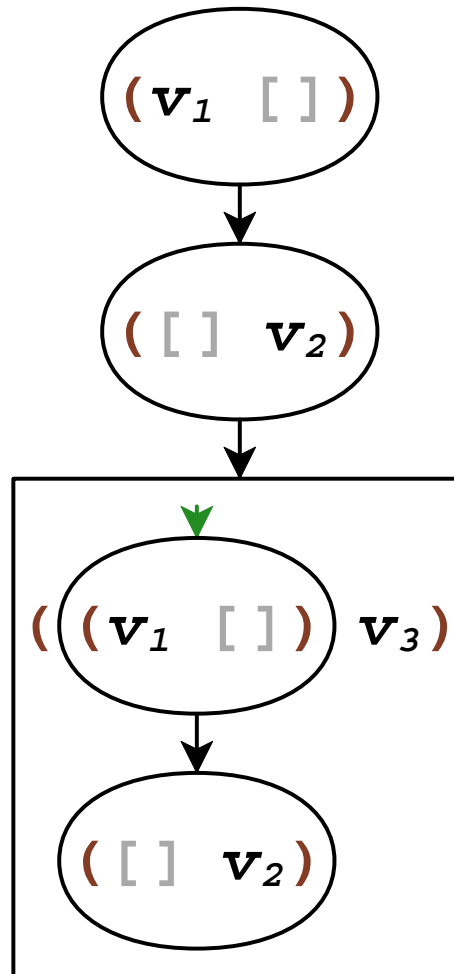
Continuations



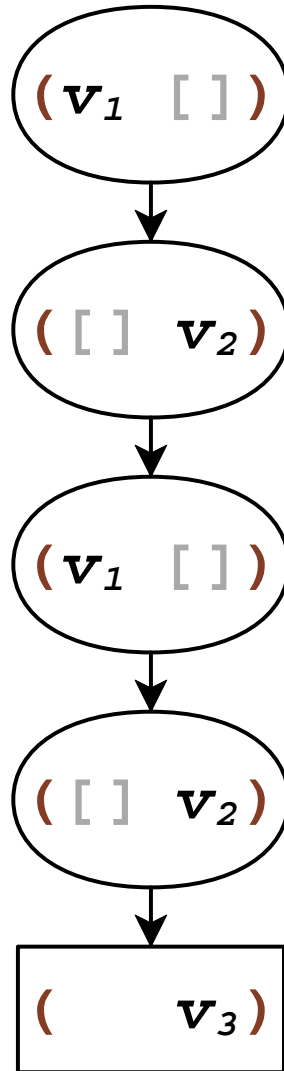
Continuations



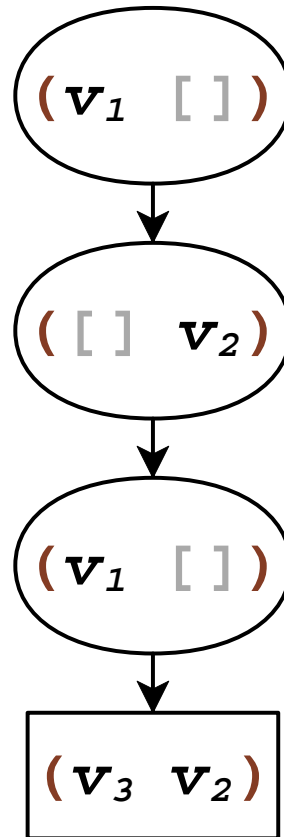
Composable Continuations



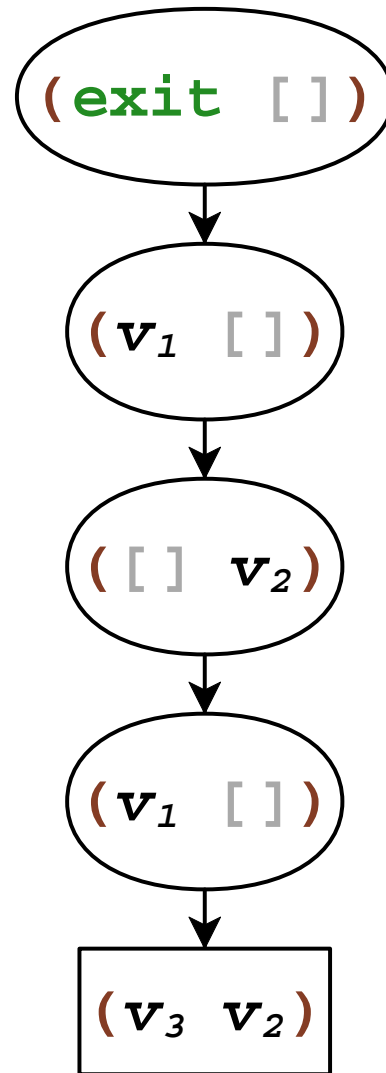
Composable Continuations



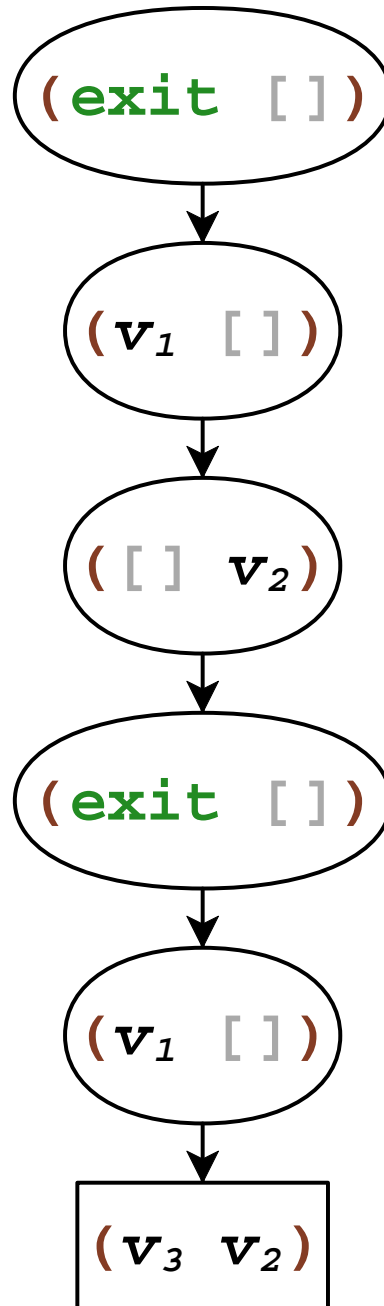
Composable Continuations



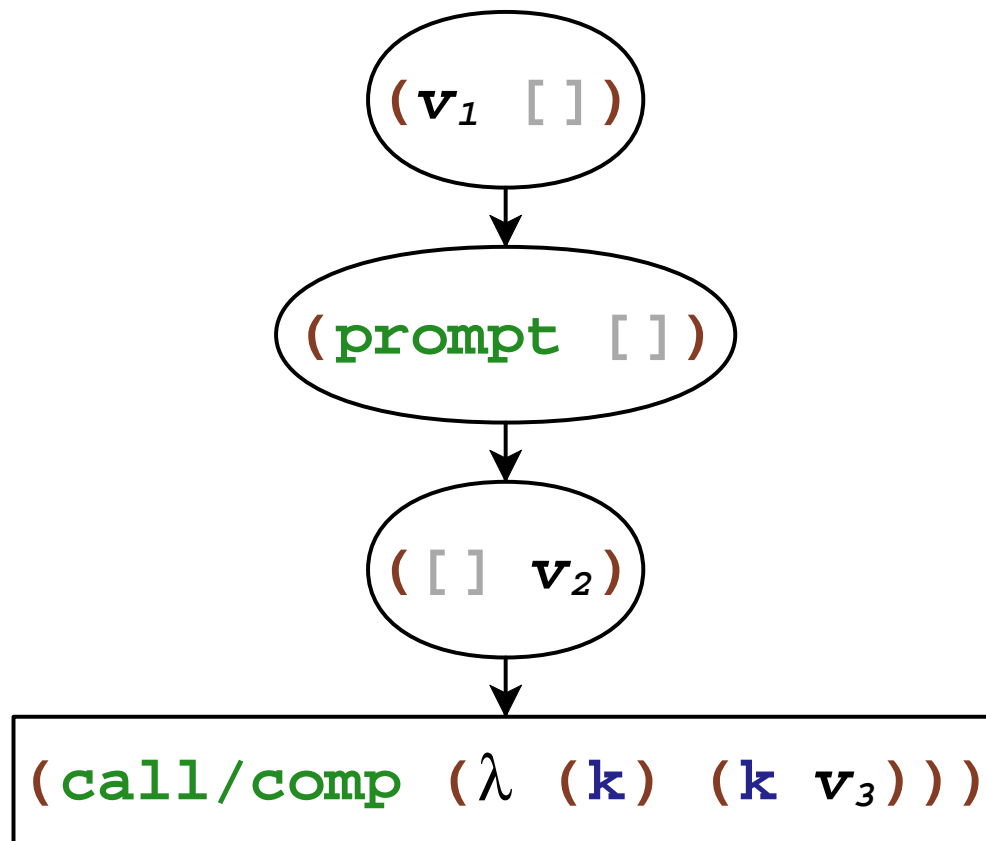
Composable Continuations



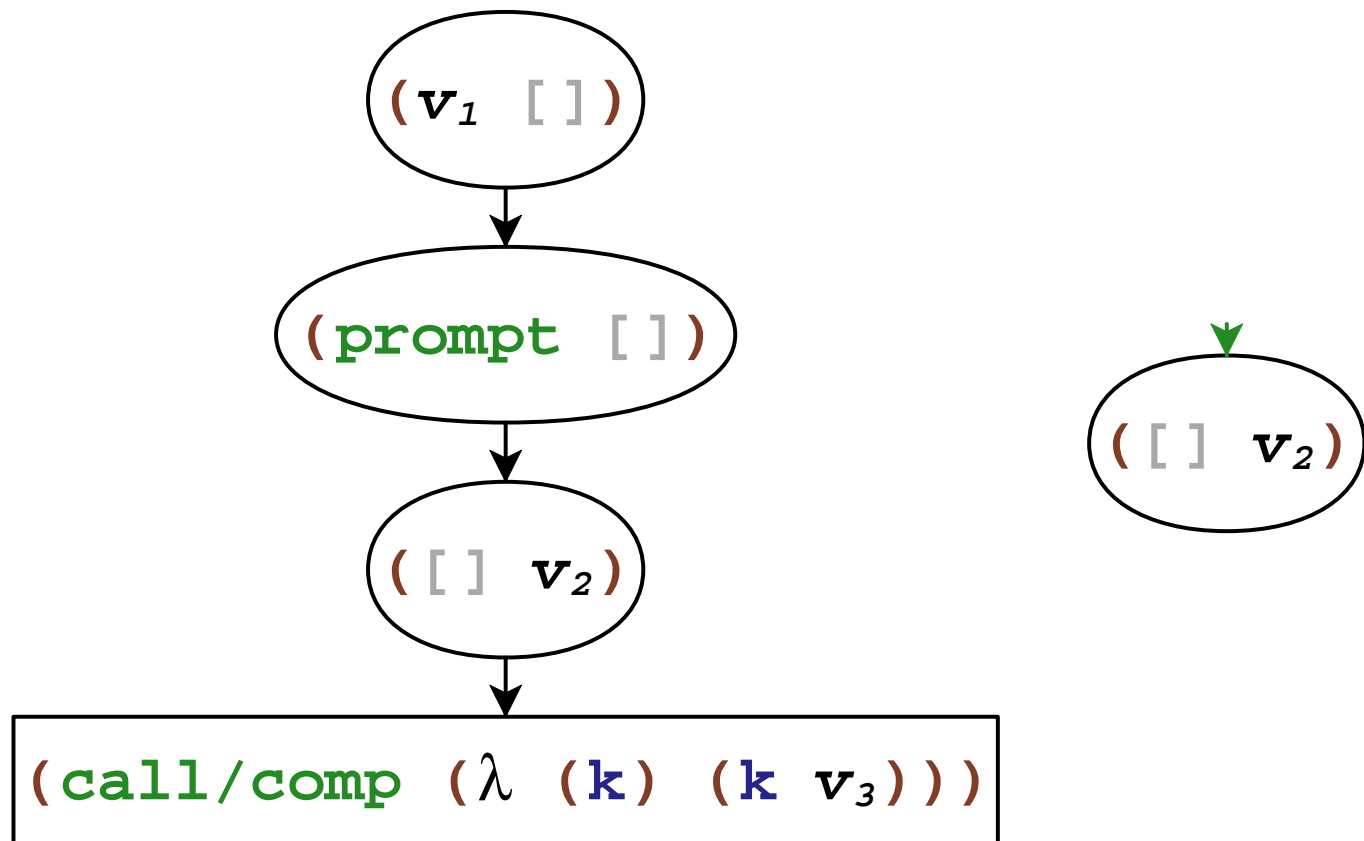
Composable Continuations



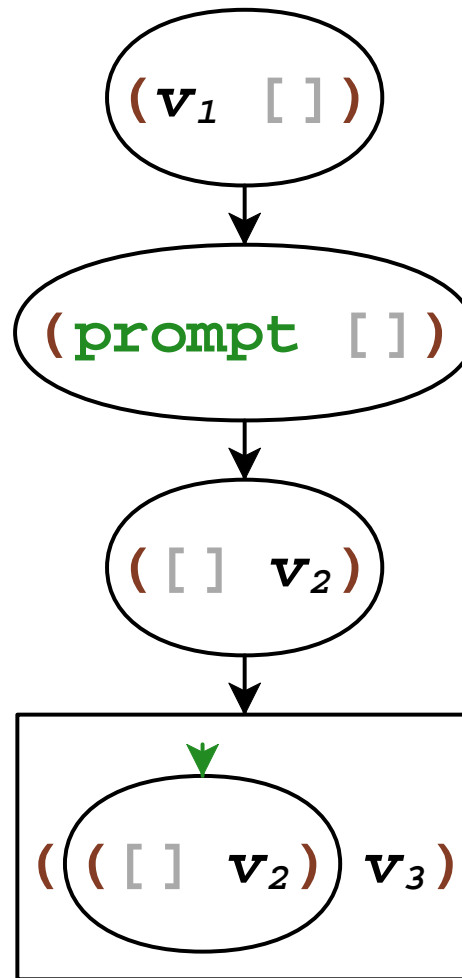
Delimited Capture



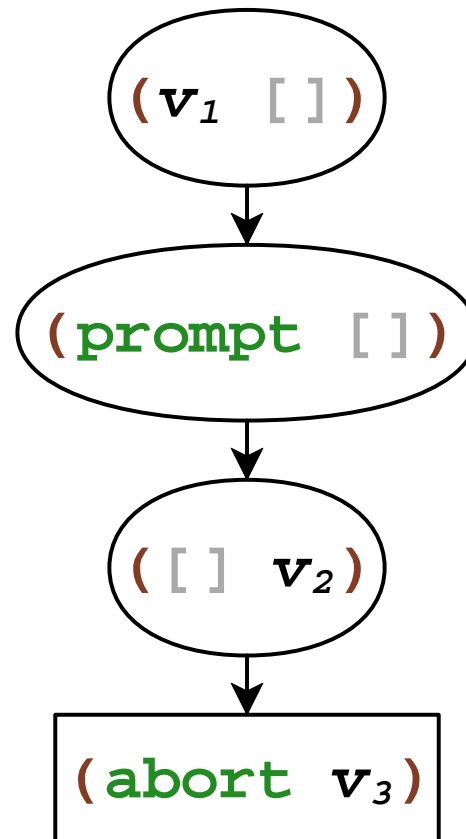
Delimited Capture



Delimited Capture



Delimited Abort



Delimited Abort

$(\mathbf{v}_1 \quad [\])$

\mathbf{v}_3

Delimited Abort

$$\boxed{(\mathbf{v}_1 \ \mathbf{v}_3)}$$

Splitting Capture and Abort

- **call/comp** : capture current continuation
- **abort** : abort current continuation
- \mathcal{F} : capture and abort current continuation

Splitting Capture and Abort

- **call/comp** : capture current continuation
- **abort** : abort current continuation
- \mathcal{F} : capture and abort current continuation

$$\mathcal{F} = (\lambda (\mathbf{f}) (\mathbf{call/comp} (\lambda (\mathbf{k}) (\mathbf{abort} (\lambda () (\mathbf{f} \mathbf{k}))))))$$

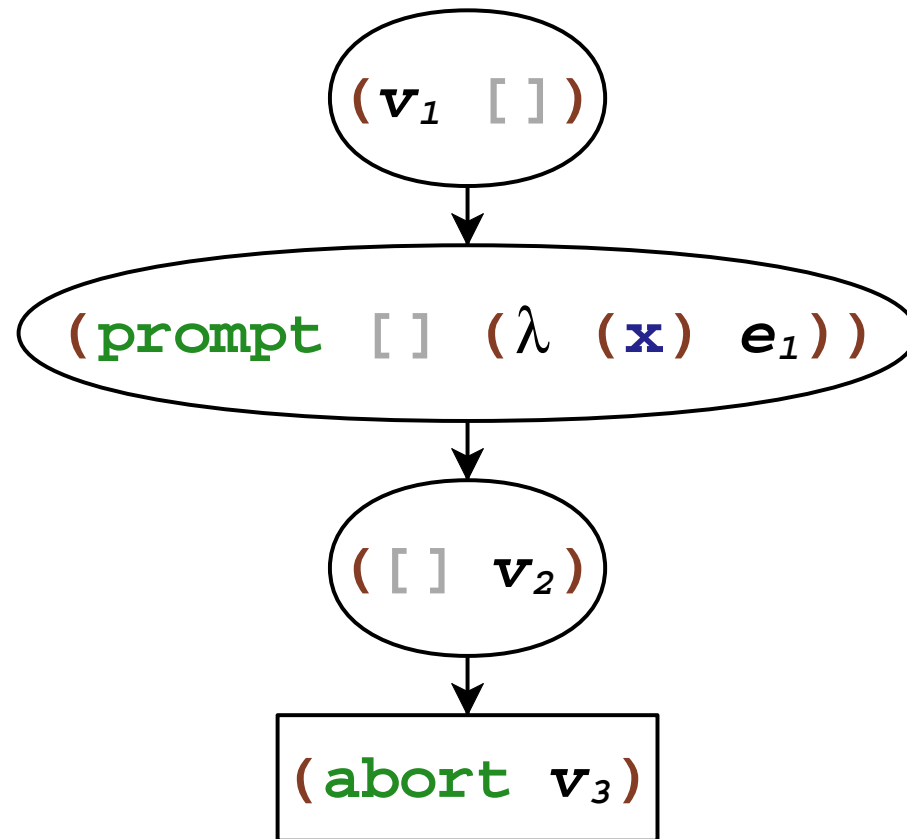
Splitting Capture and Abort

- **call/comp** : capture current continuation
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- \mathcal{F} : capture and abort current continuation

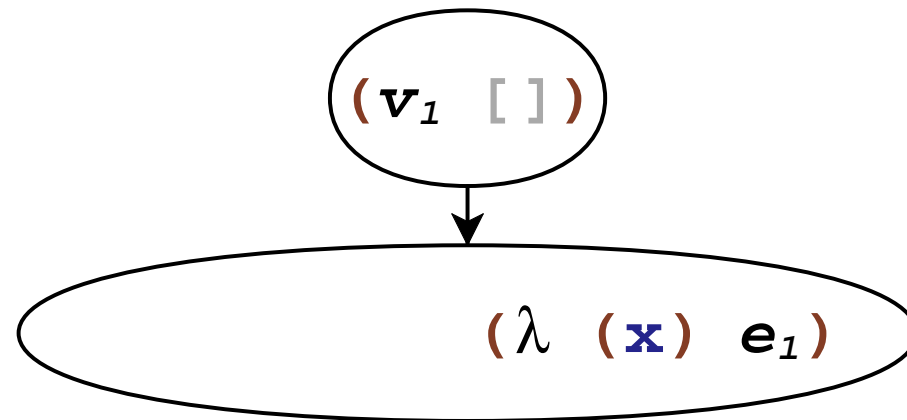
$\mathcal{F} = (\lambda (\mathbf{f}) (\mathbf{call/comp}$
 $\quad (\lambda (\mathbf{k}) (\mathbf{abort} (\lambda () (\mathbf{f} \mathbf{k}))))))$

delay ($\mathbf{f} \mathbf{k}$)
until after abort

Prompt with Handler

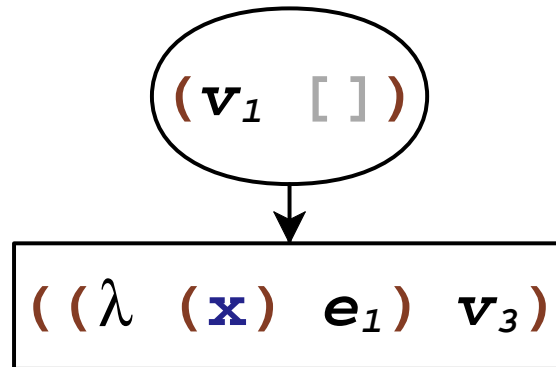


Prompt with Handler

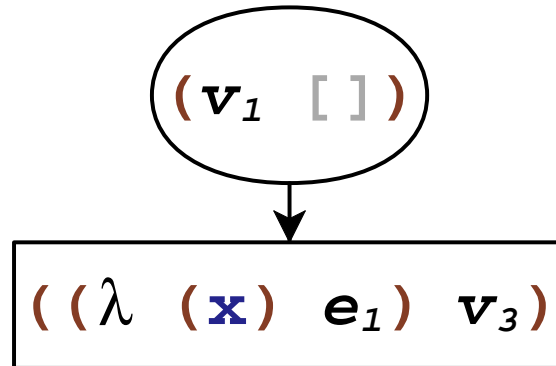


\mathbf{v}_3

Prompt with Handler



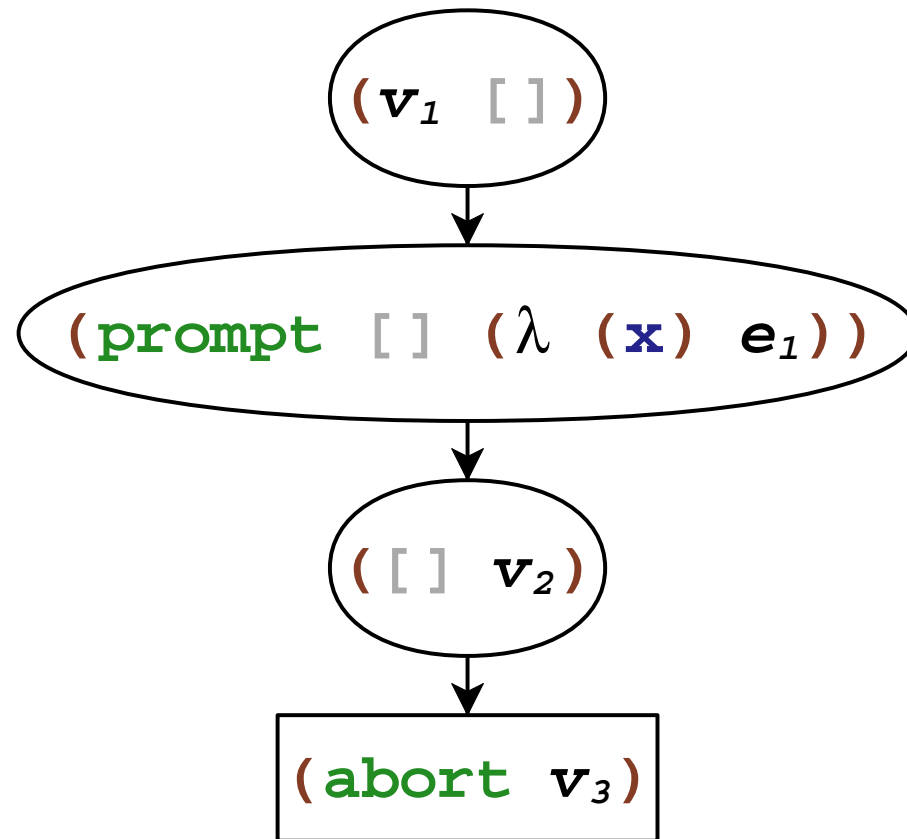
Prompt with Handler



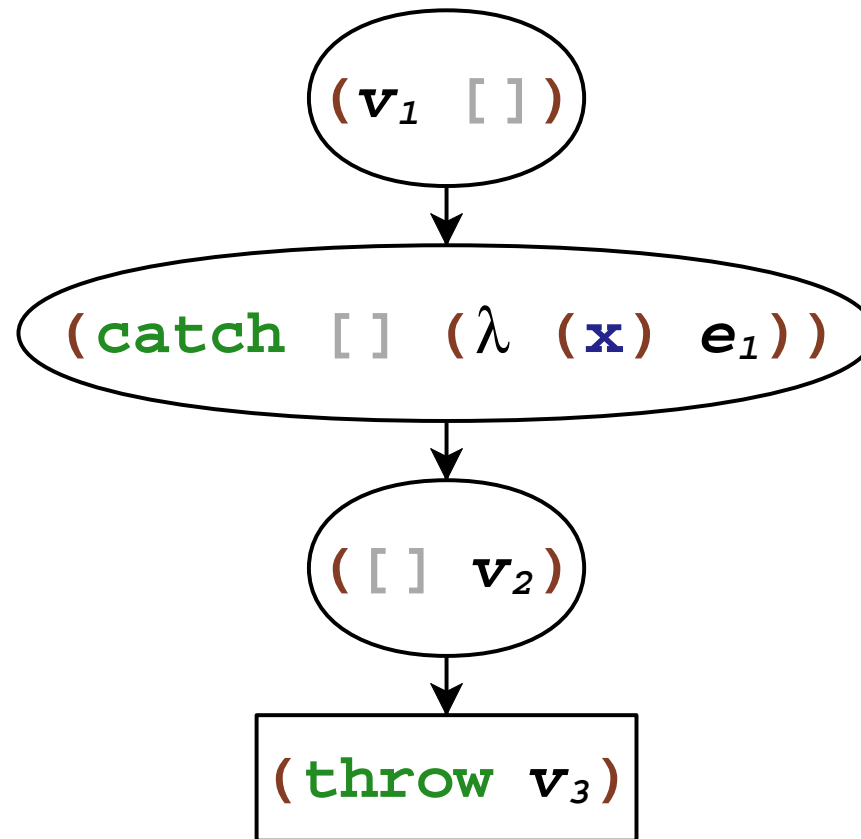
$(\text{prompt } e_1) = (\text{prompt } e_1 (\lambda (t) (t)))$

$\mathcal{F} = (\lambda (f) (\text{call/comp} (\lambda (k) (\text{abort } (\lambda () (f k))))))$

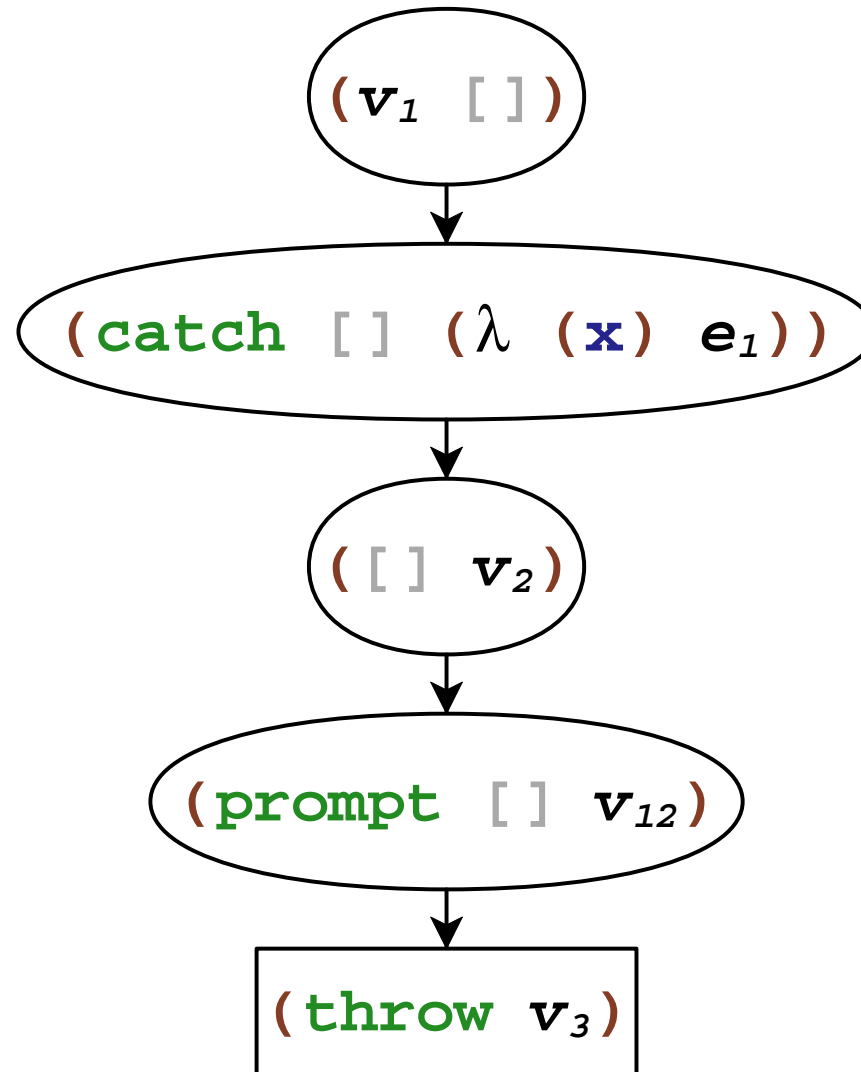
Prompt with Handler



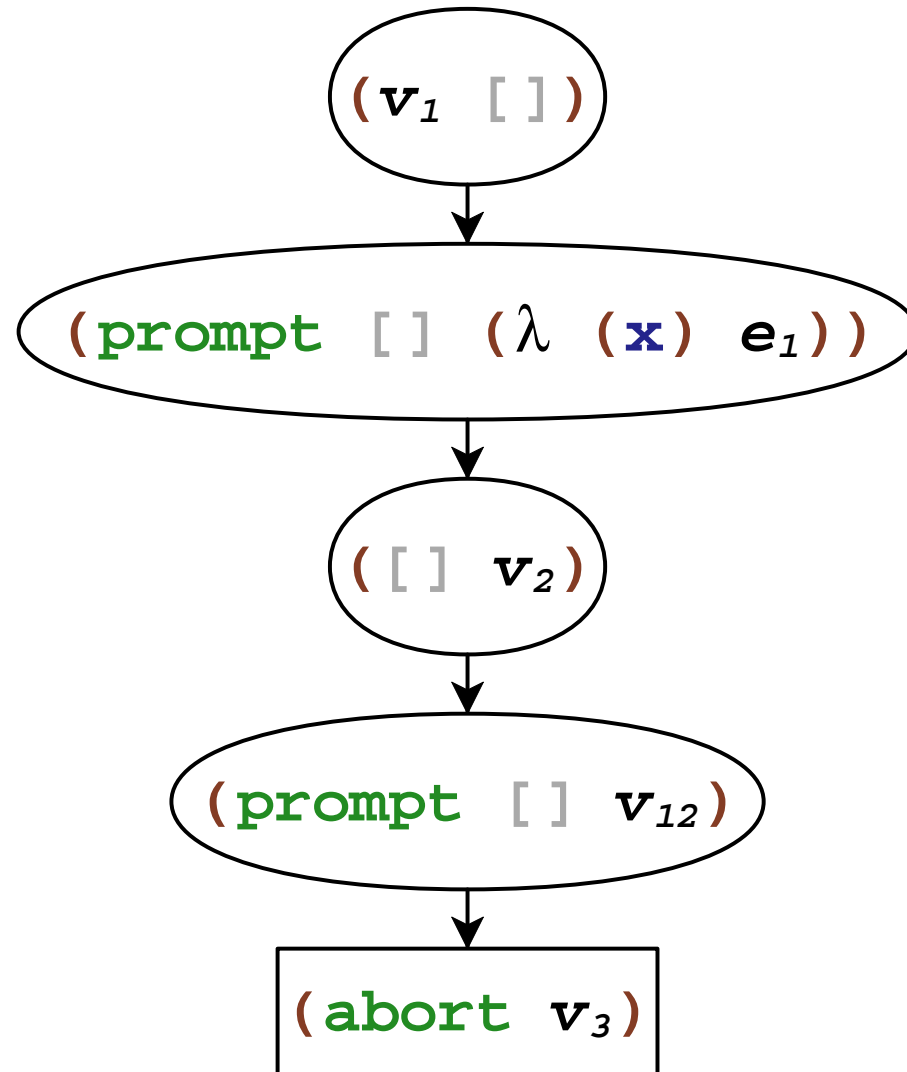
Catch and Throw



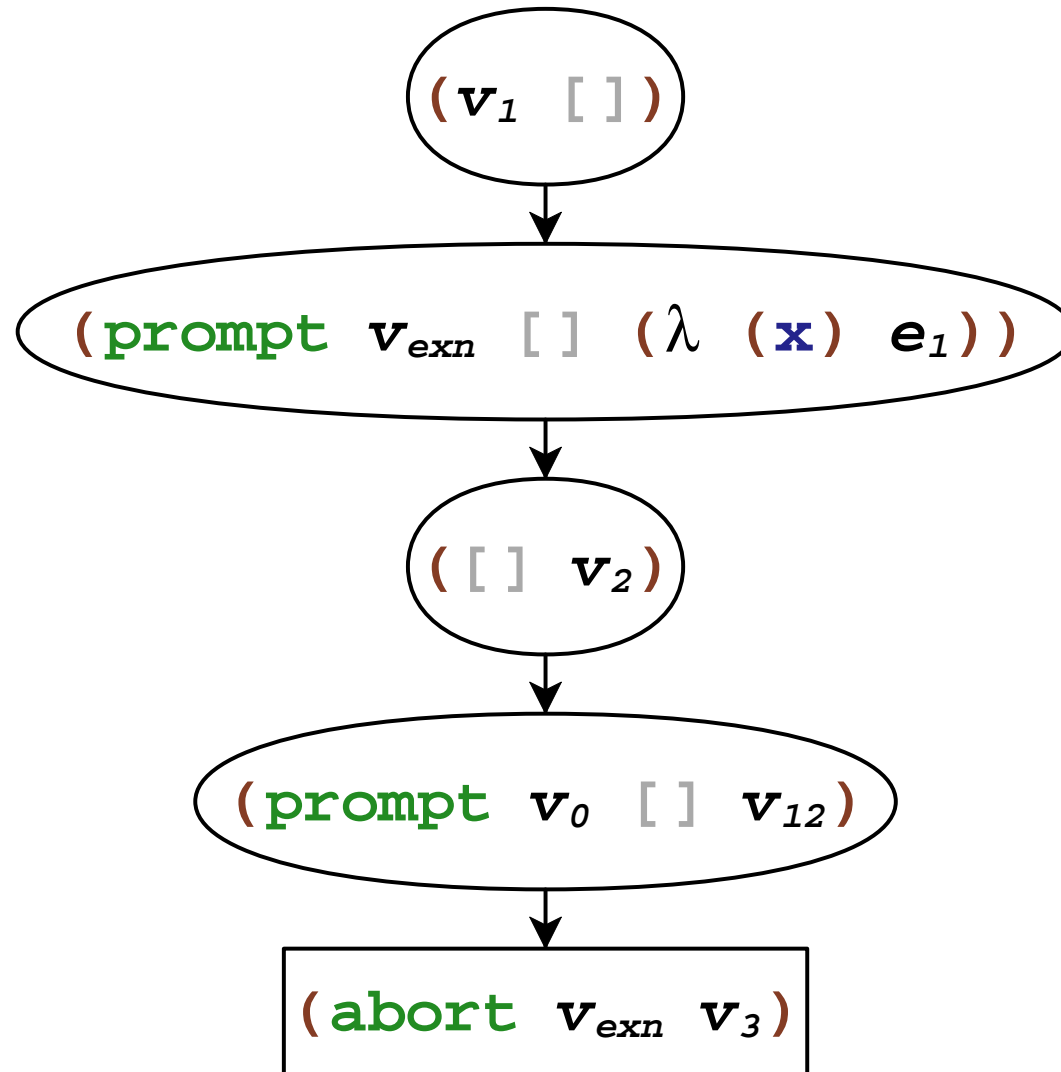
Catch and Throw



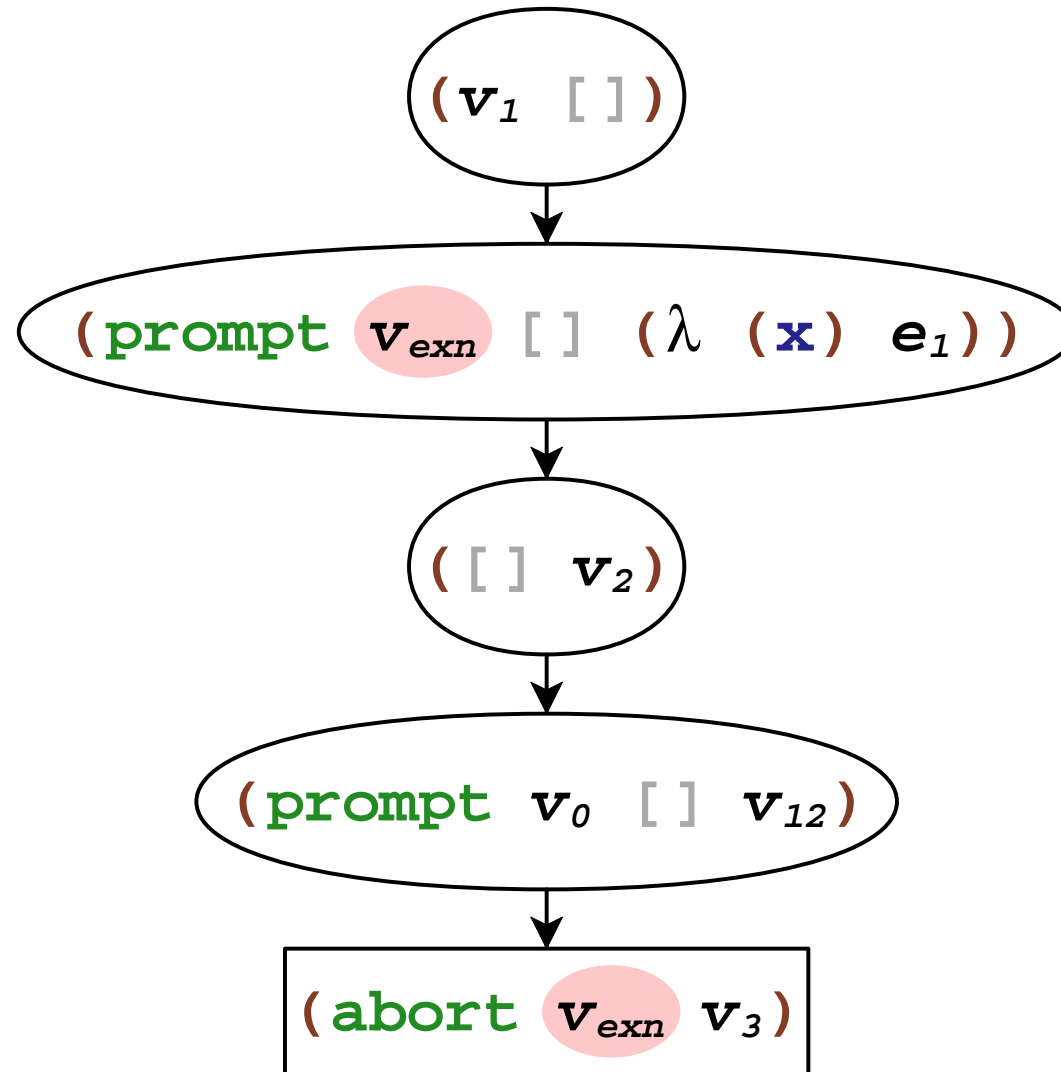
Catch and Throw



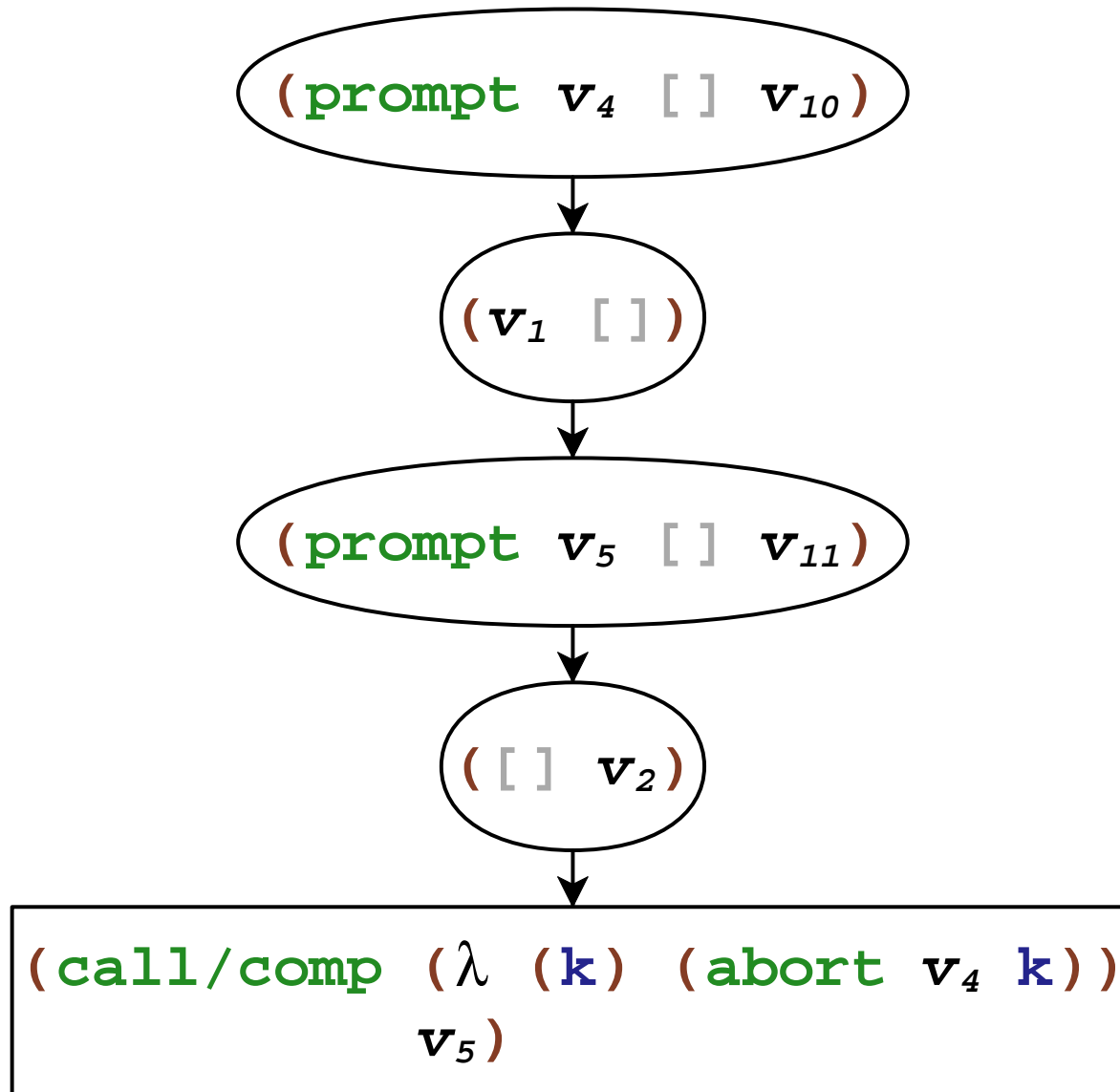
Catch and Throw



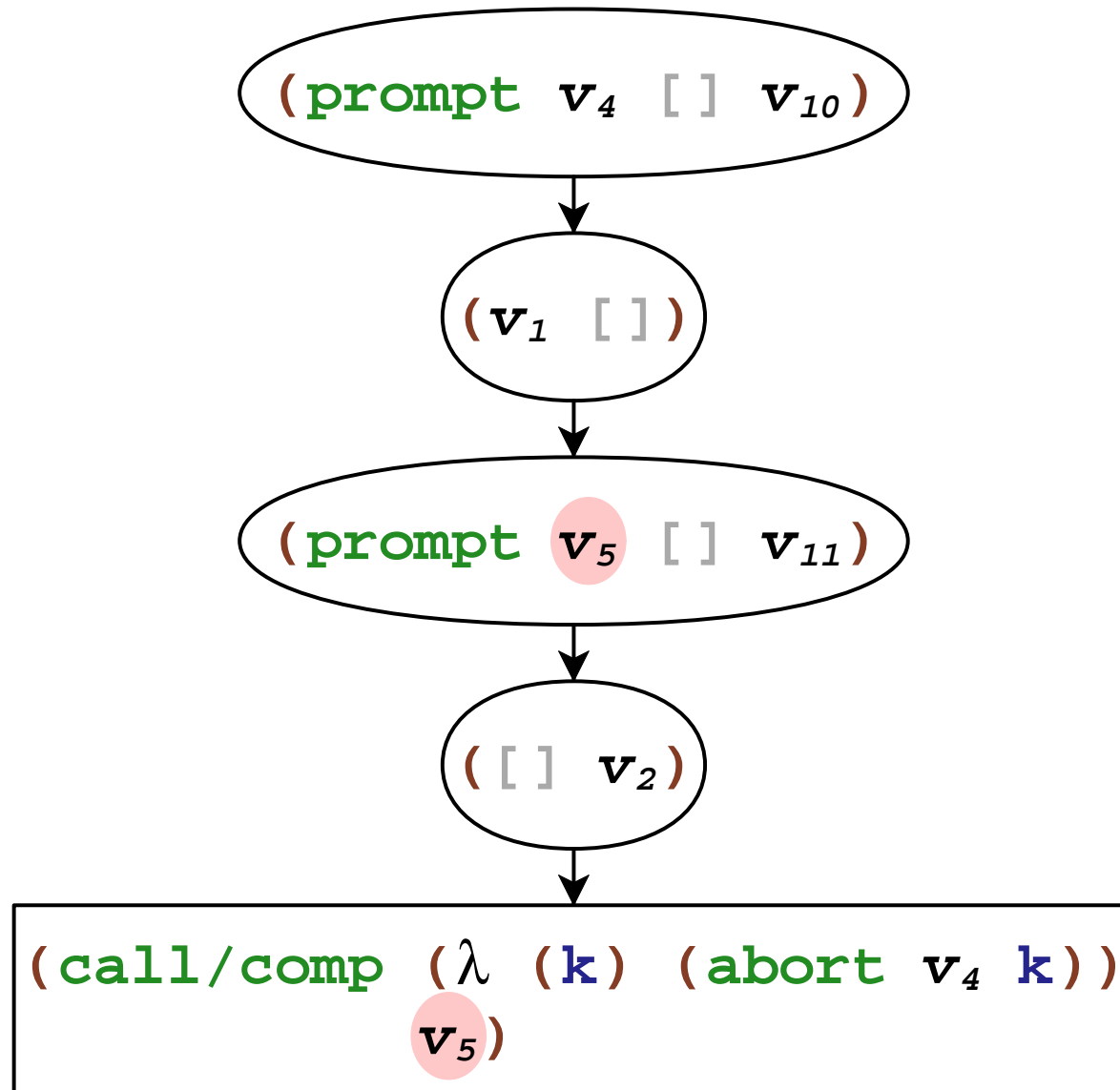
Tagged Prompts



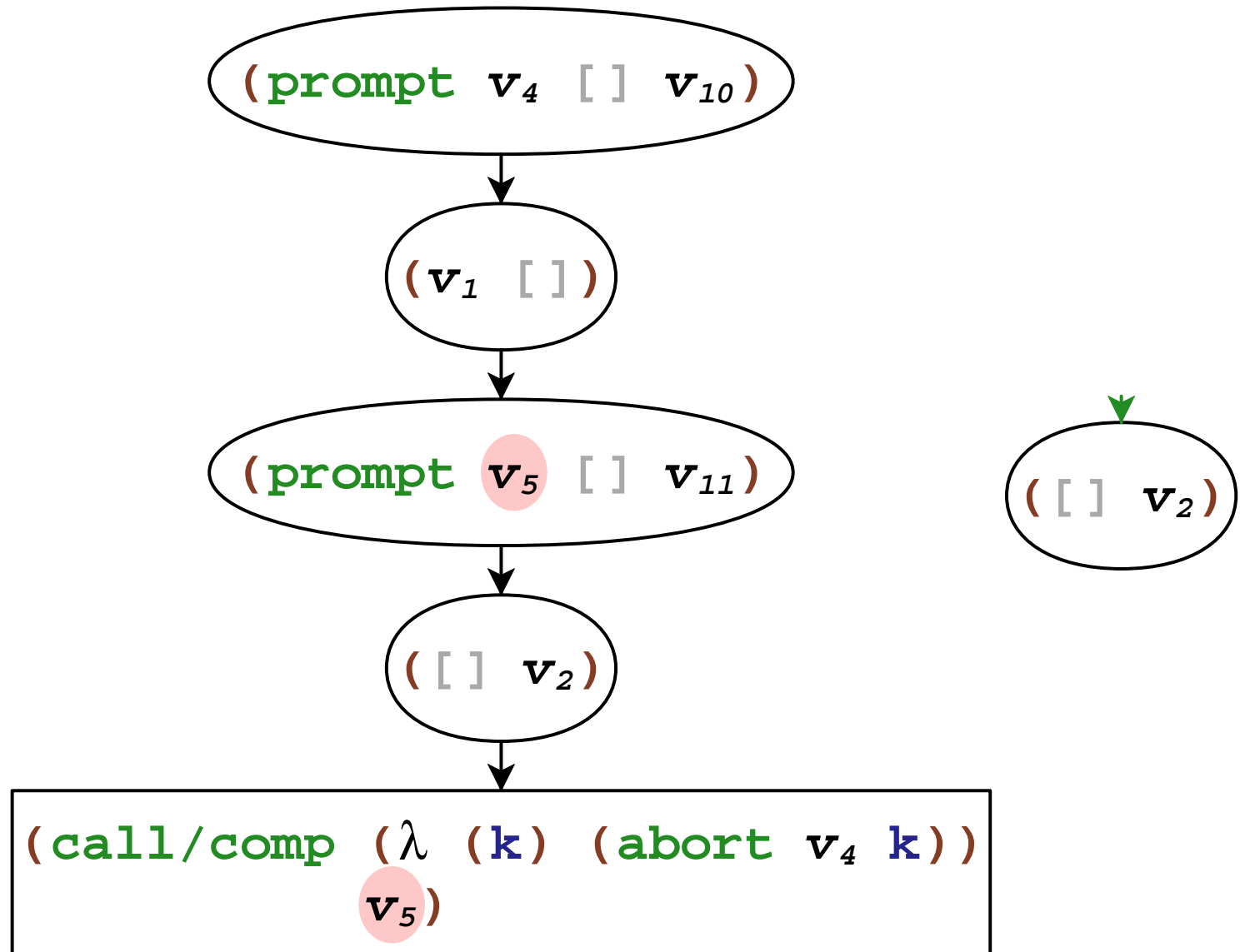
Tagged Prompts



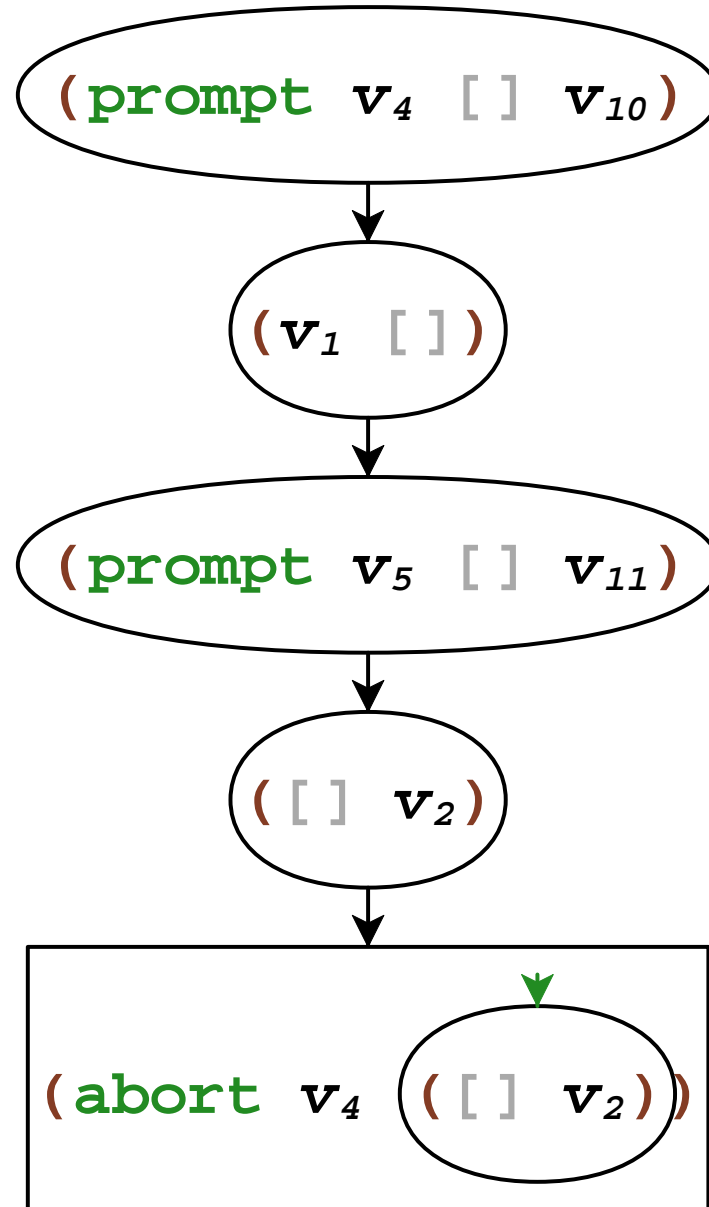
Tagged Prompts



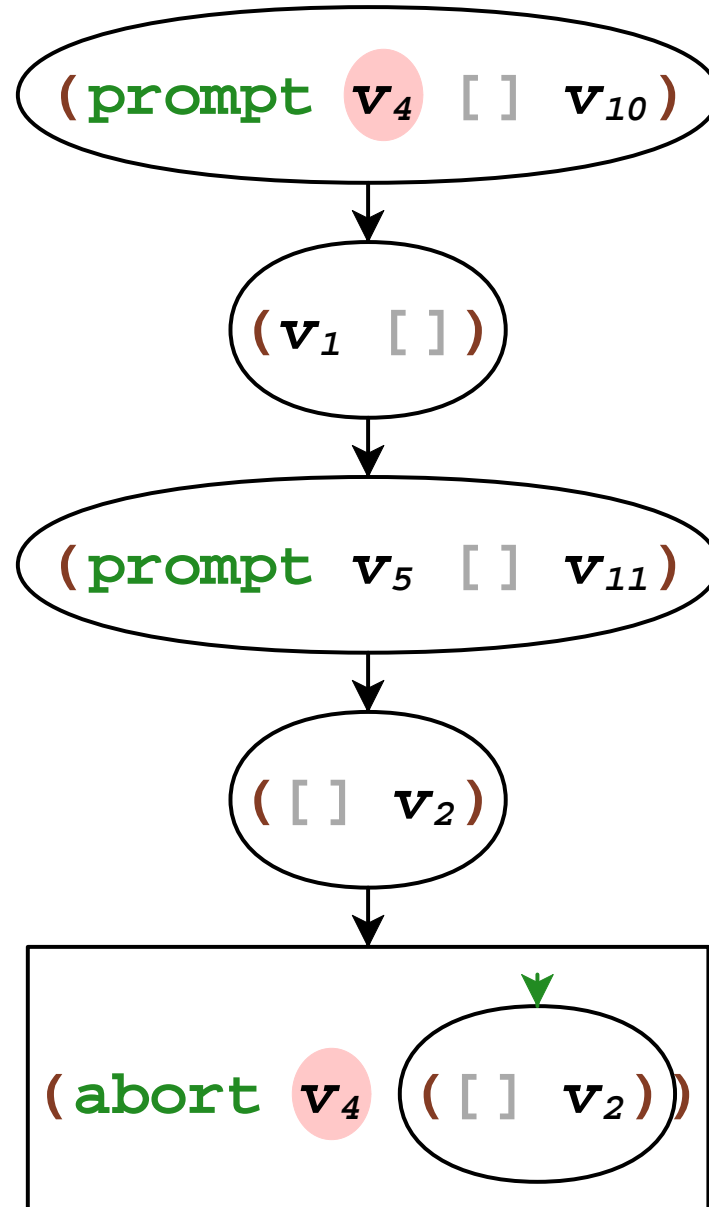
Tagged Prompts



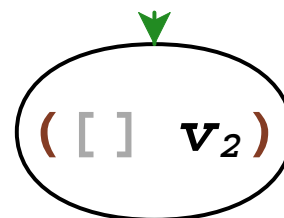
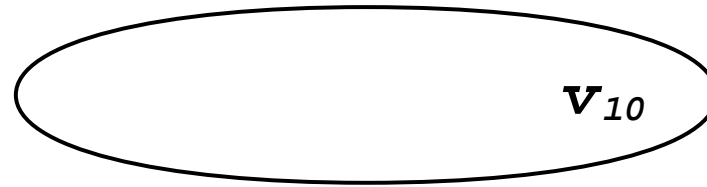
Tagged Prompts



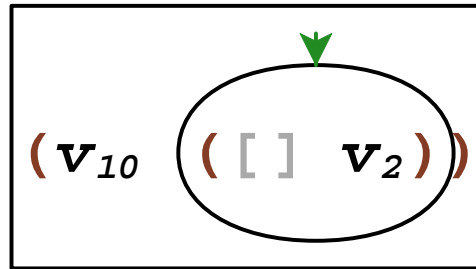
Tagged Prompts



Tagged Prompts



Tagged Prompts

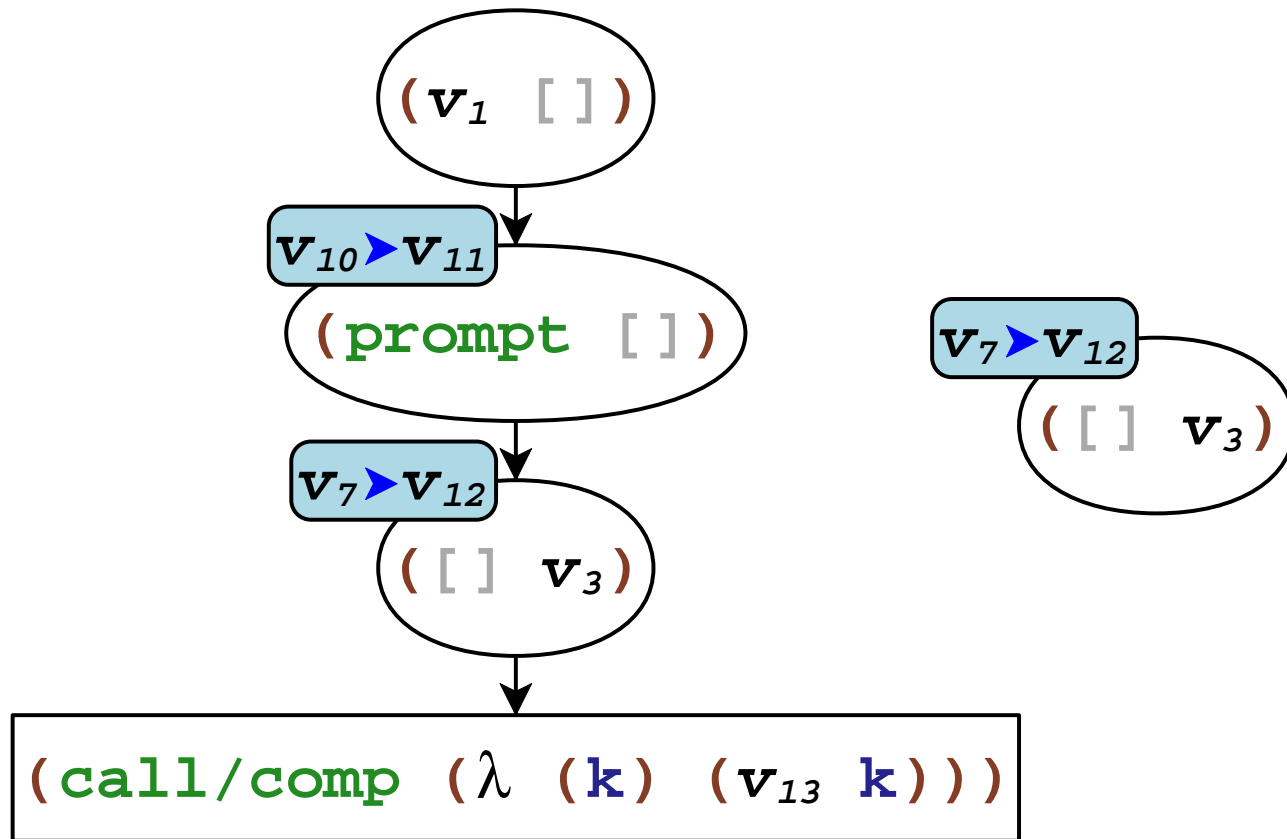


Continuations Summary

- `prompt` with handler and tag
- `abort` with tag
- `call/comp` with tag
- continuation composition
- plain `call/cc` uses default tag

[Sitaram PLDI'93]

Continuation Marks



Dynamic Binding Summary

- `call/cm` to add marks
- `current-marks` to get marks, up to a tag
- capture marks in `call/comp`
- splice marks in continuation composition

Side Effects and Control

```
(define (with-resource work-thunk)
  (begin
    (grab-resource)
    (work-thunk)
    (release-with-resource)))
```

Dynamic Wind

```
(define (with-resource work-thunk)
  (dynamic-wind
    (λ () (grab-resource))
    (λ () (work-thunk))
    (λ () (release-with-resource))))
```


Dynamic Wind

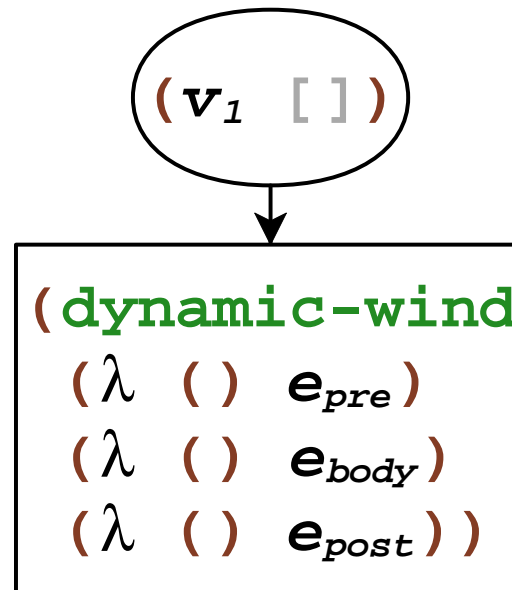
(dynamic-wind

(λ () e_{pre})

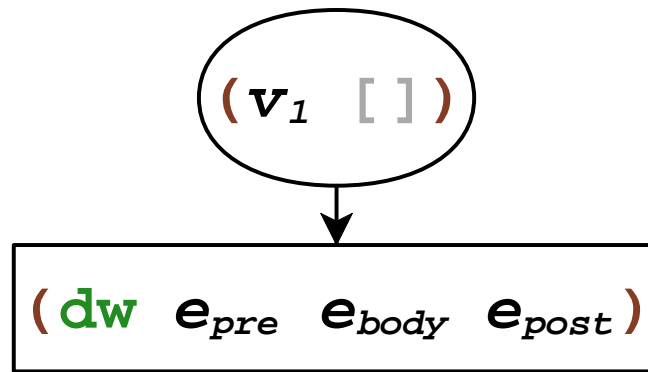
(λ () e_{body})

(λ () e_{post}))

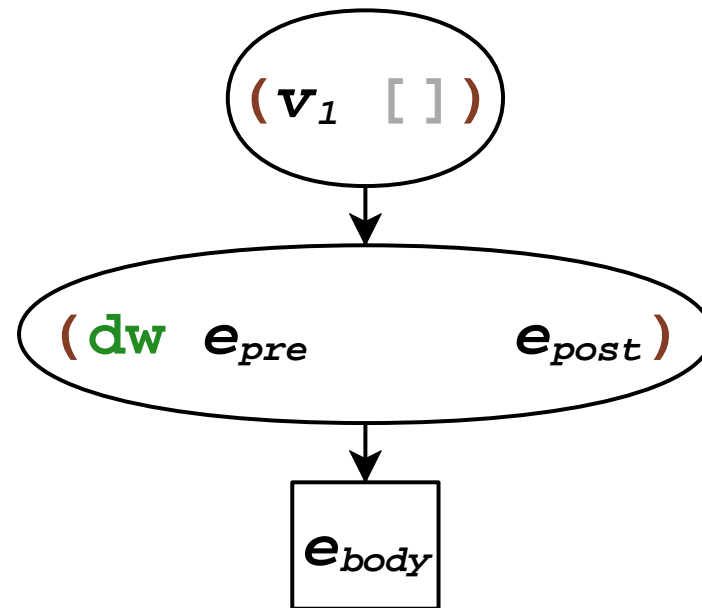
Dynamic Wind



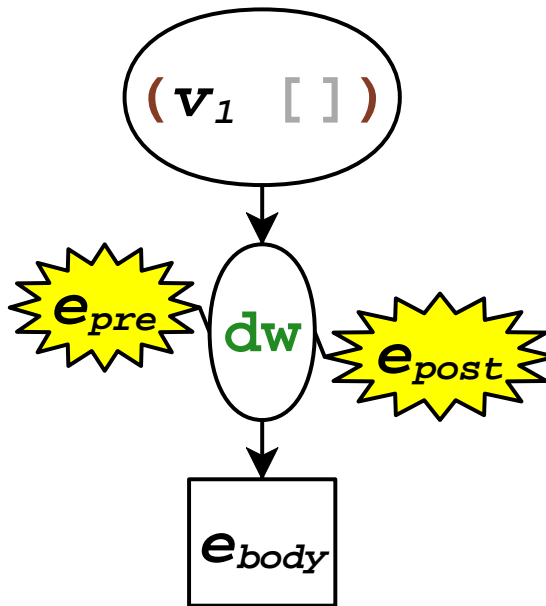
Dynamic Wind



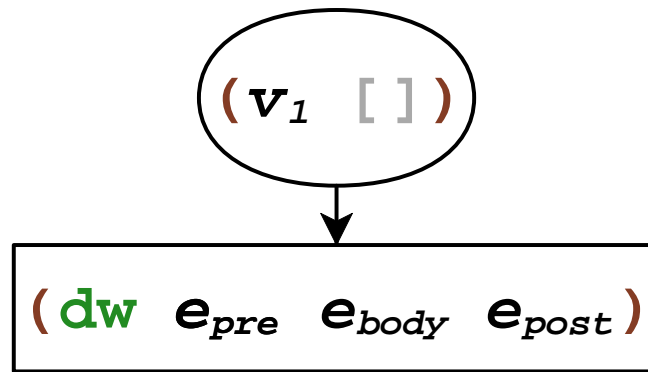
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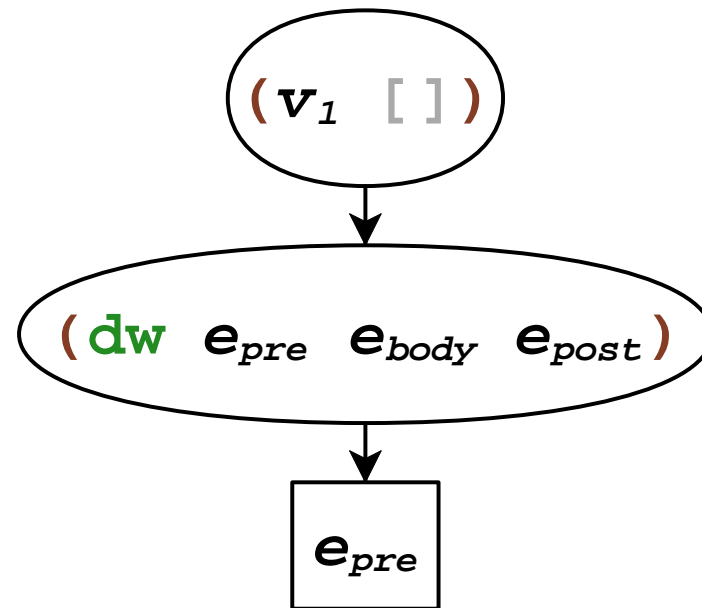
Dynamic Wind



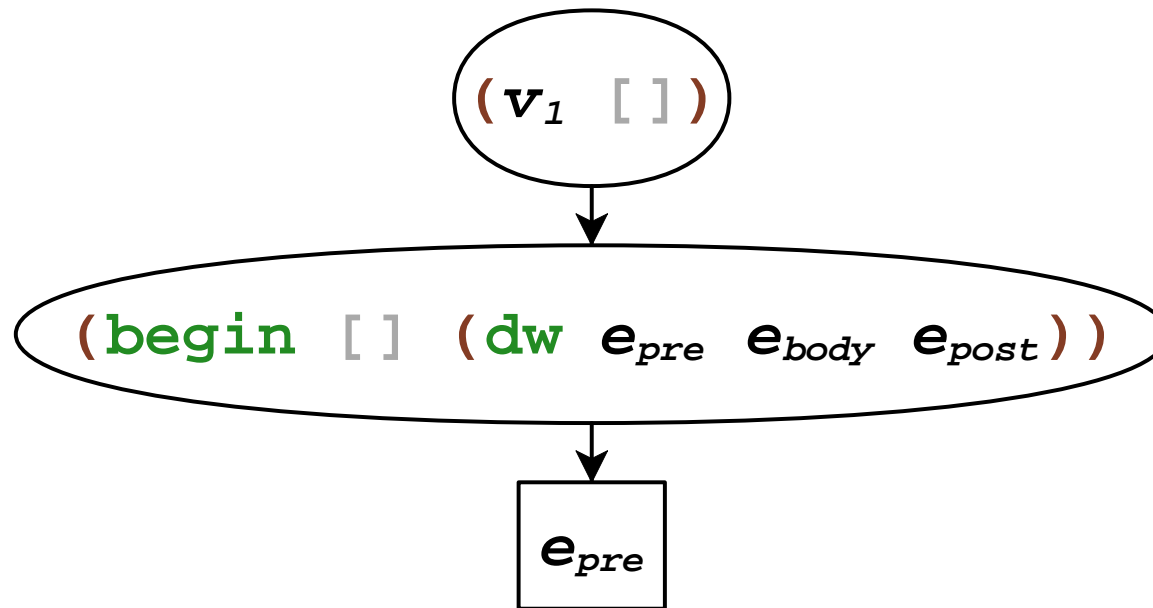
Dynamic Wind Evaluation



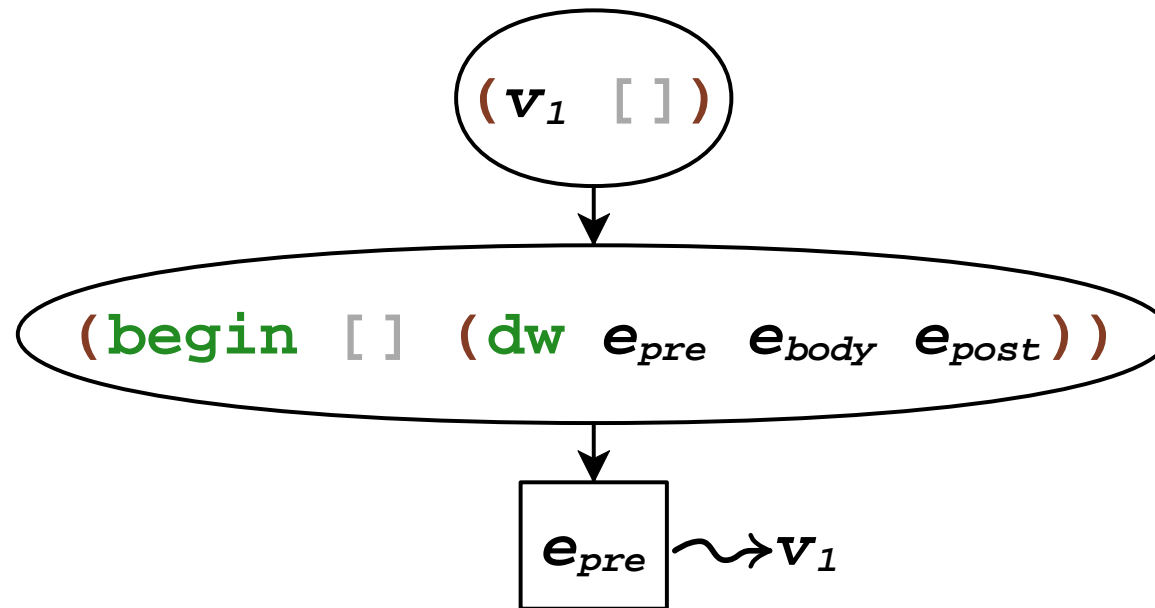
Dynamic Wind Evaluation



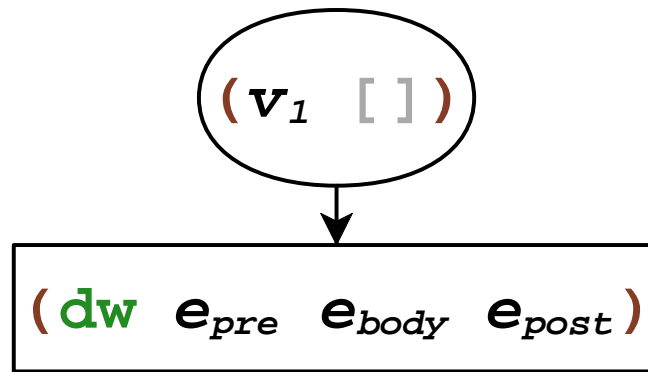
Dynamic Wind Evaluation



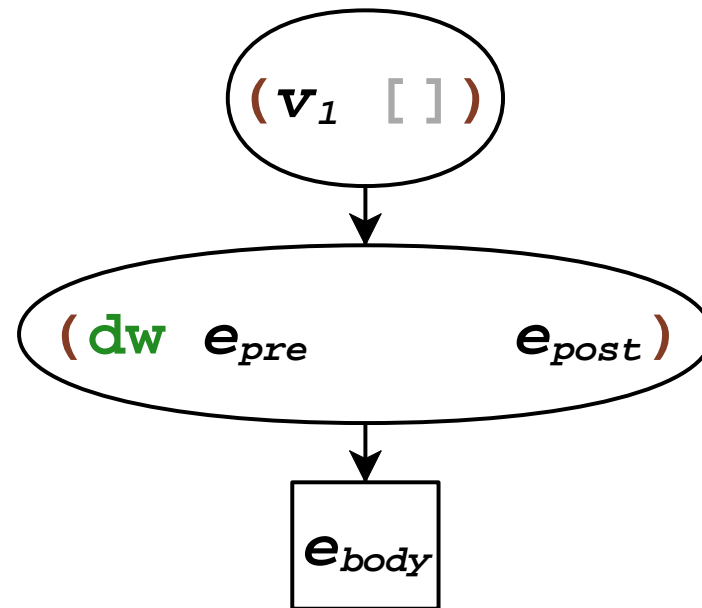
Dynamic Wind Evaluation



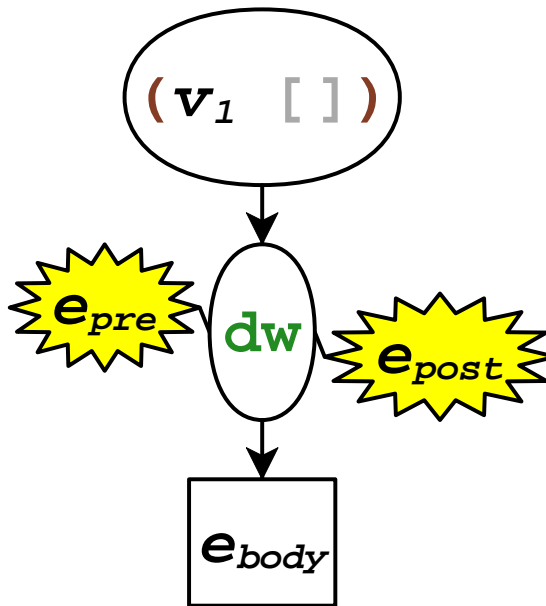
Dynamic Wind Evaluation



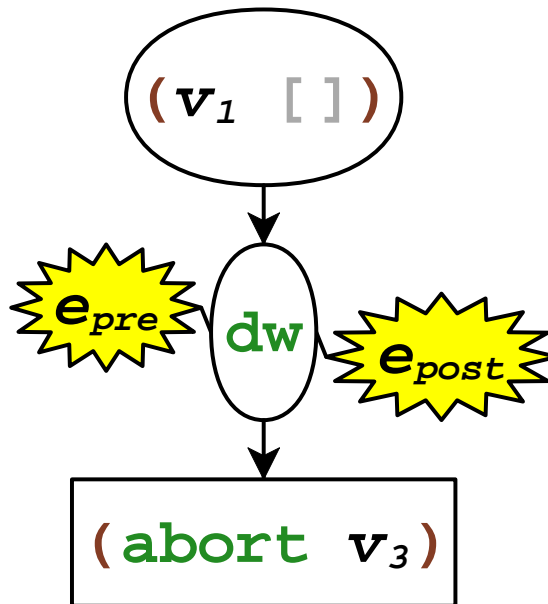
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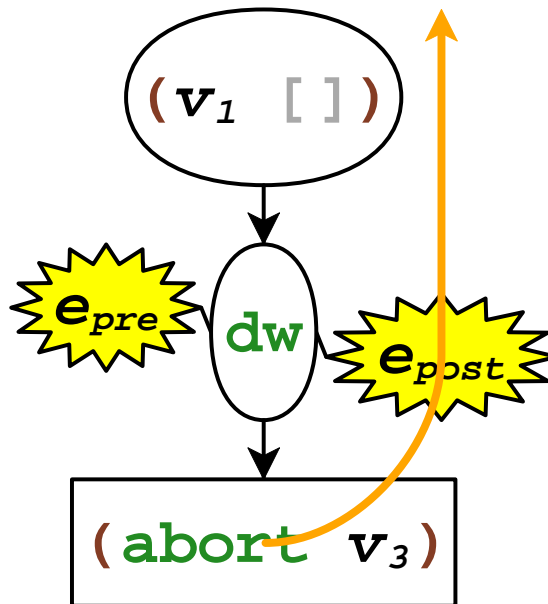
Dynamic Wind and Jumps



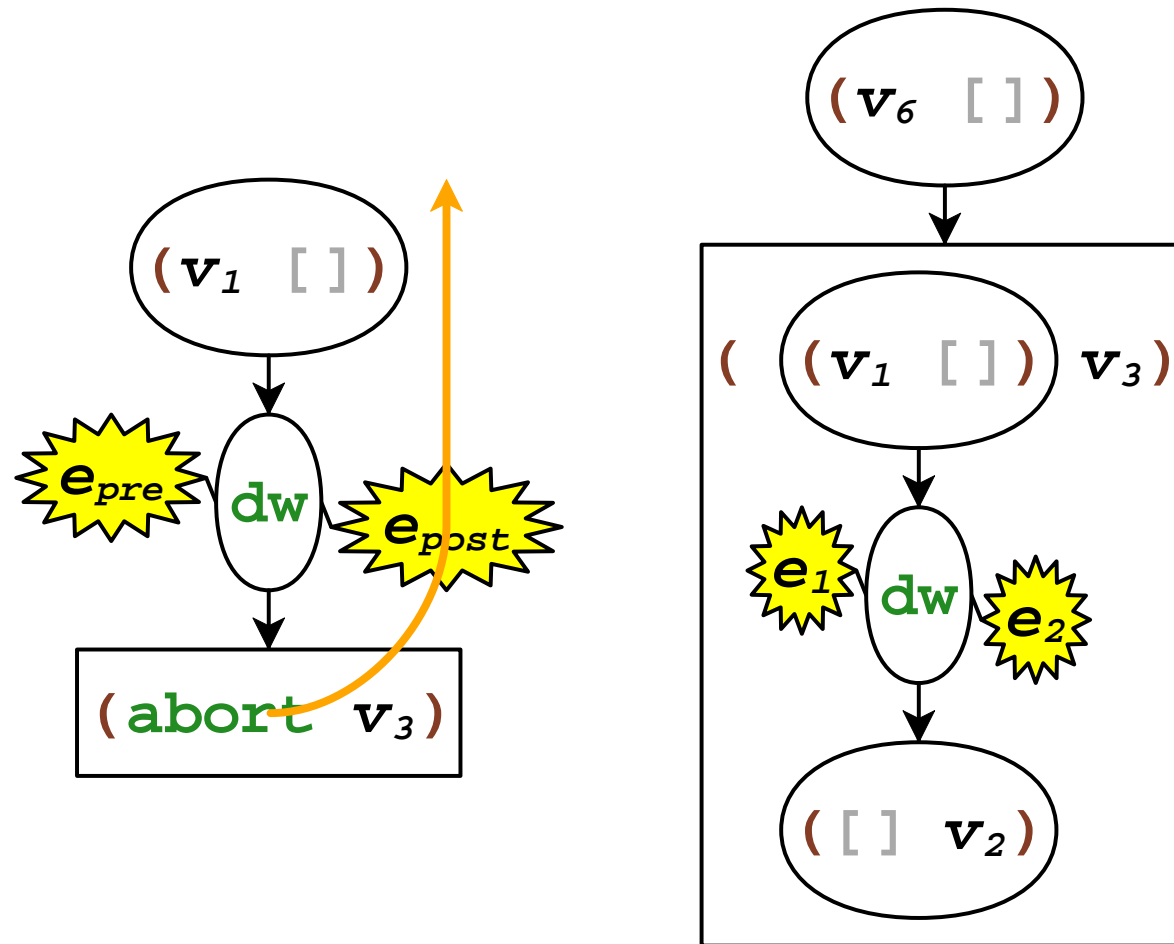
Dynamic Wind and Jumps



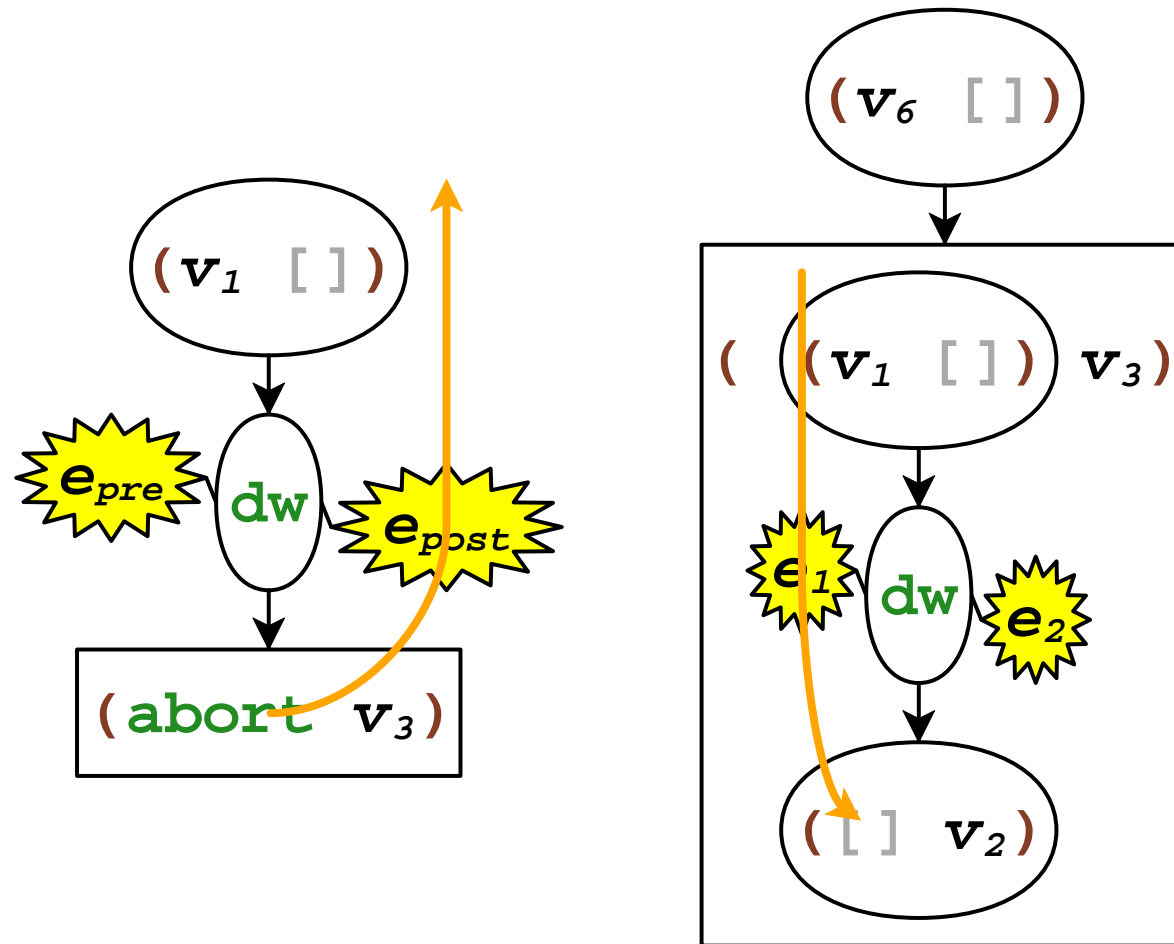
Dynamic Wind and Jumps



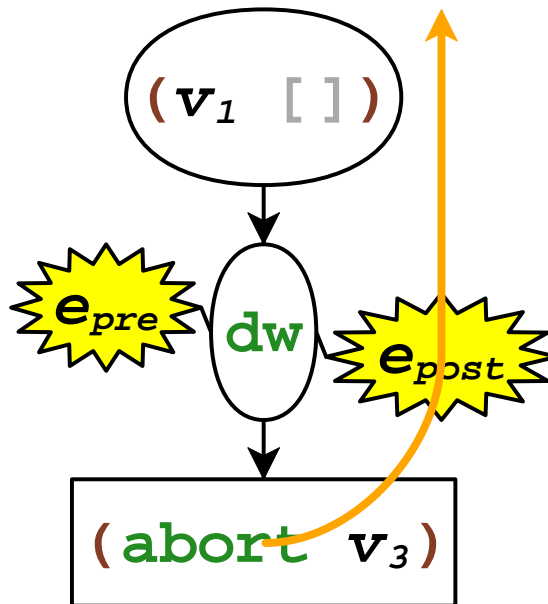
Dynamic Wind and Jumps



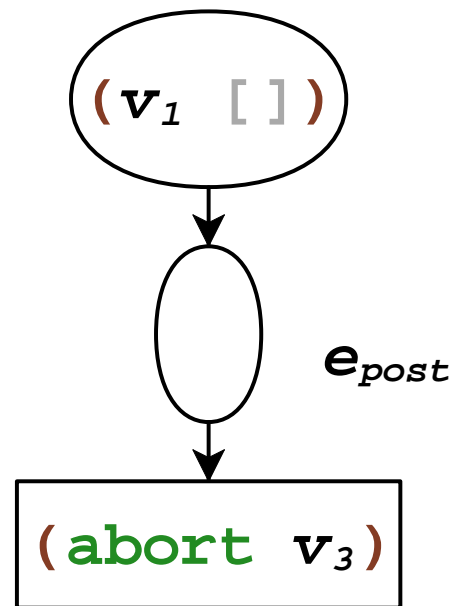
Dynamic Wind and Jumps



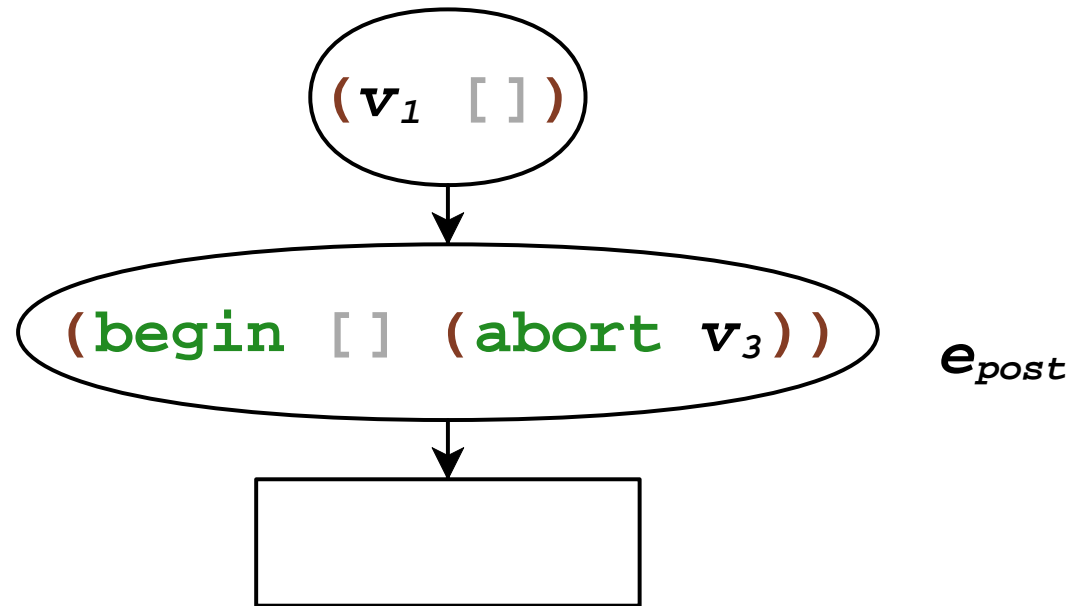
Dynamic Wind and Abort Evaluation



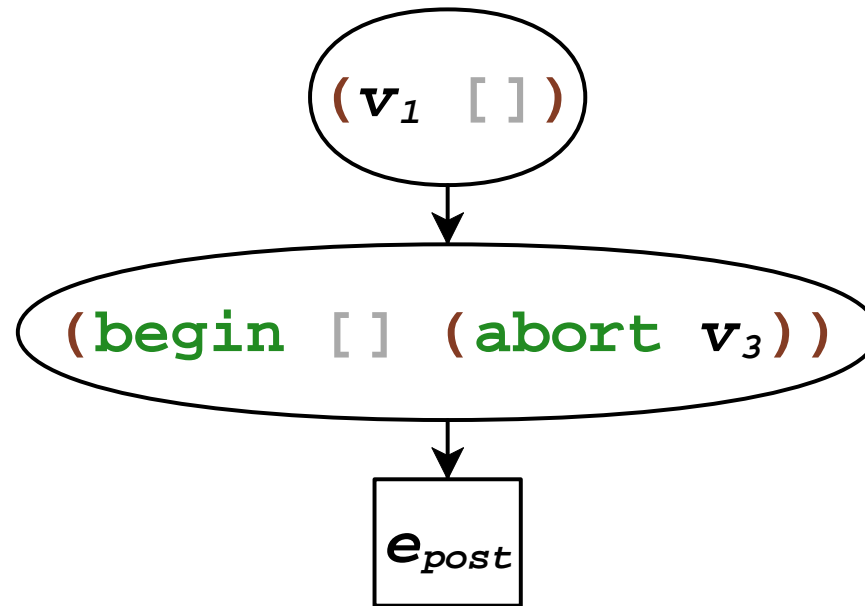
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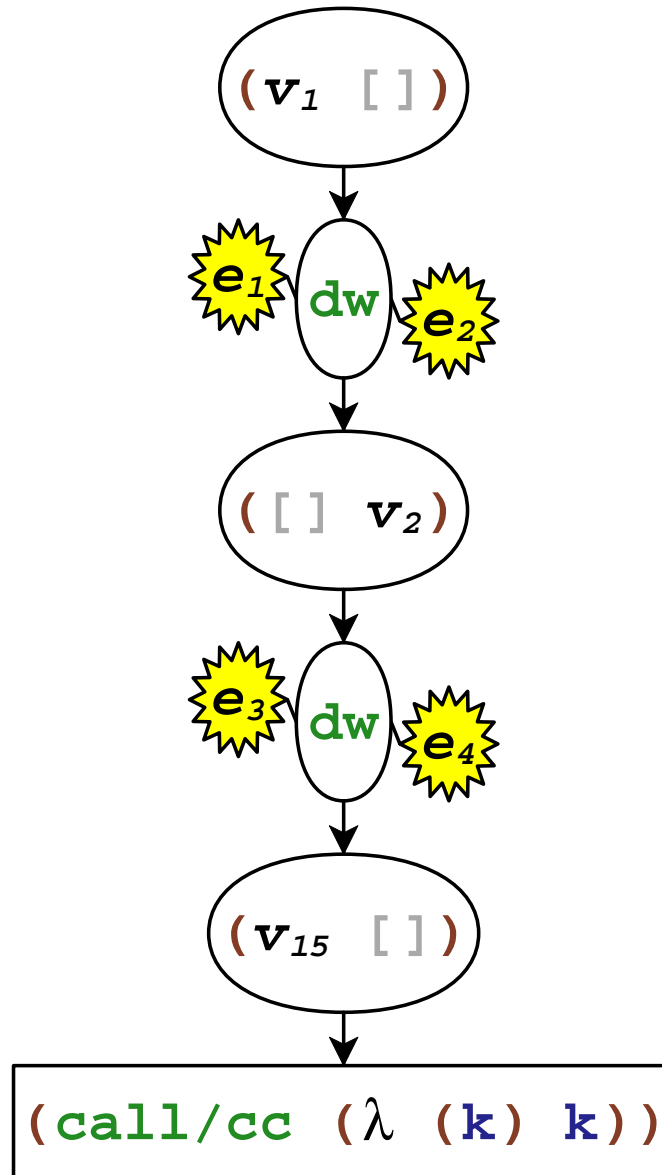
Dynamic Wind and Abort Evaluation



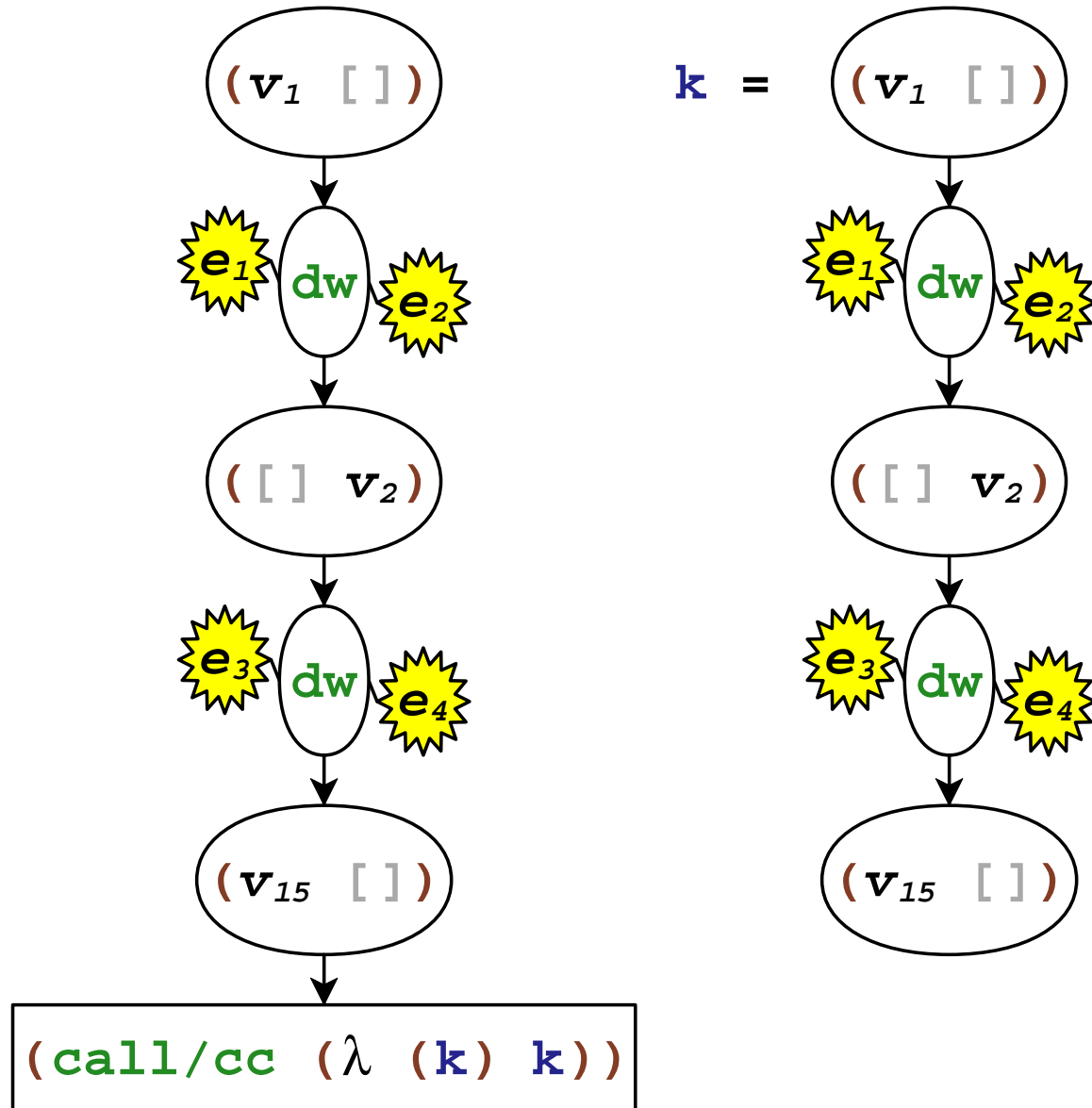
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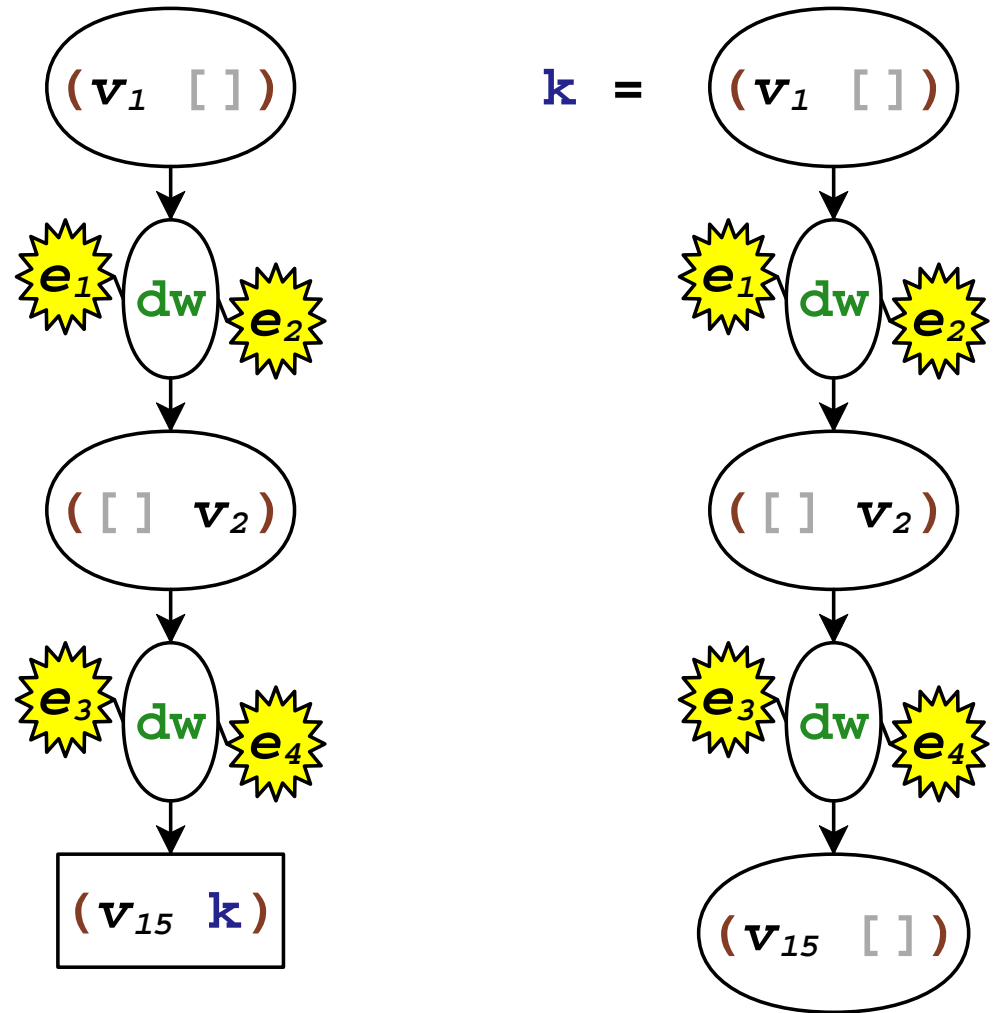
Dynamic Wind and Call/cc



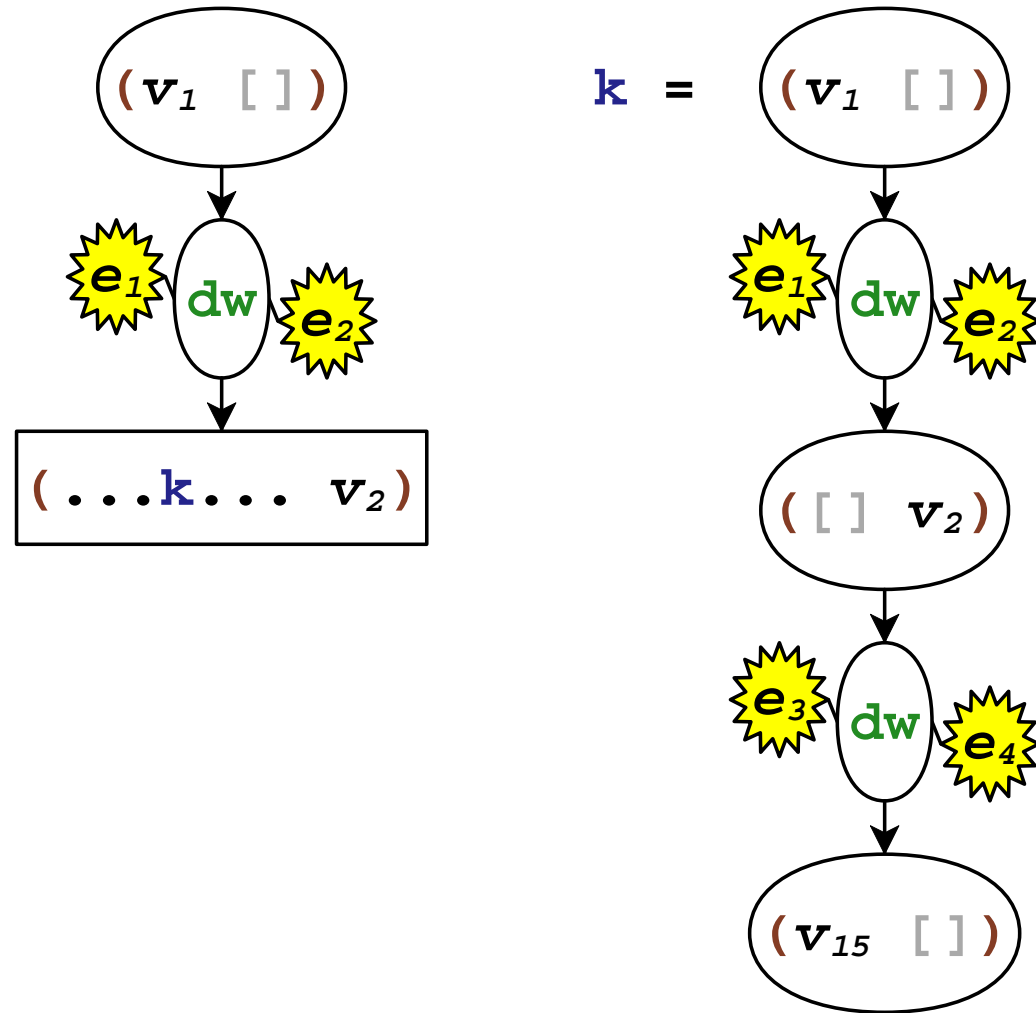
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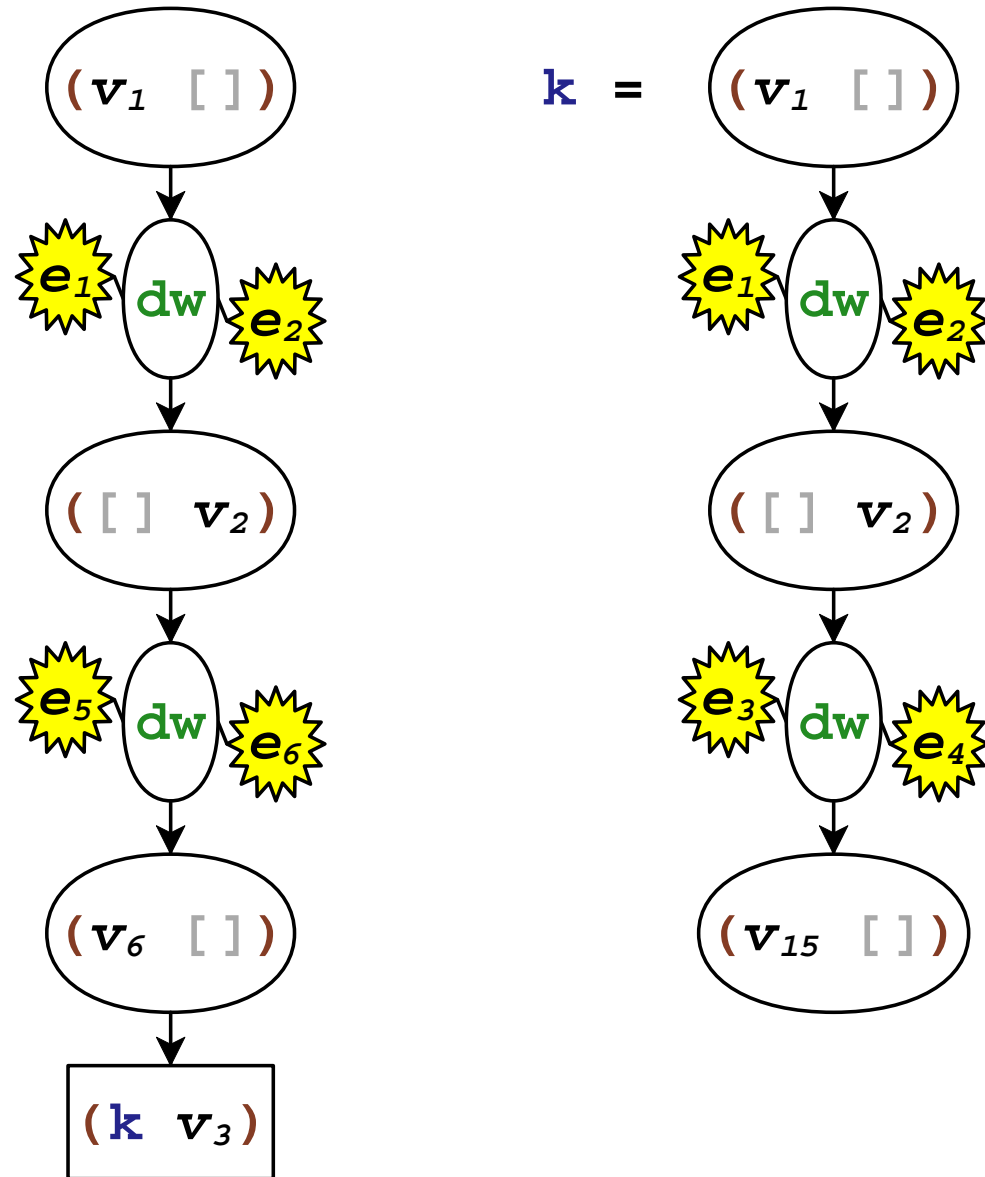
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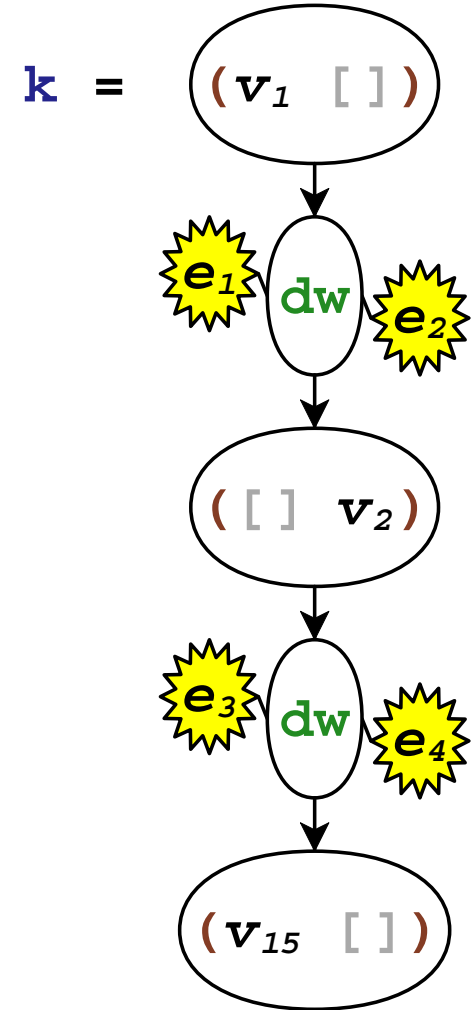
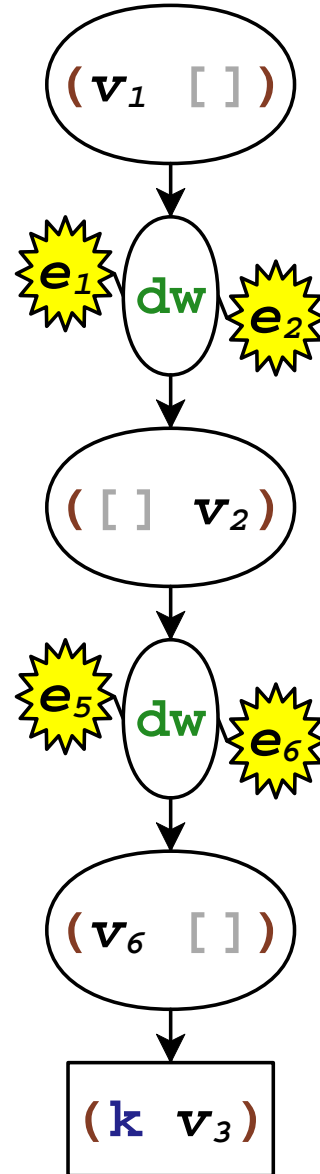
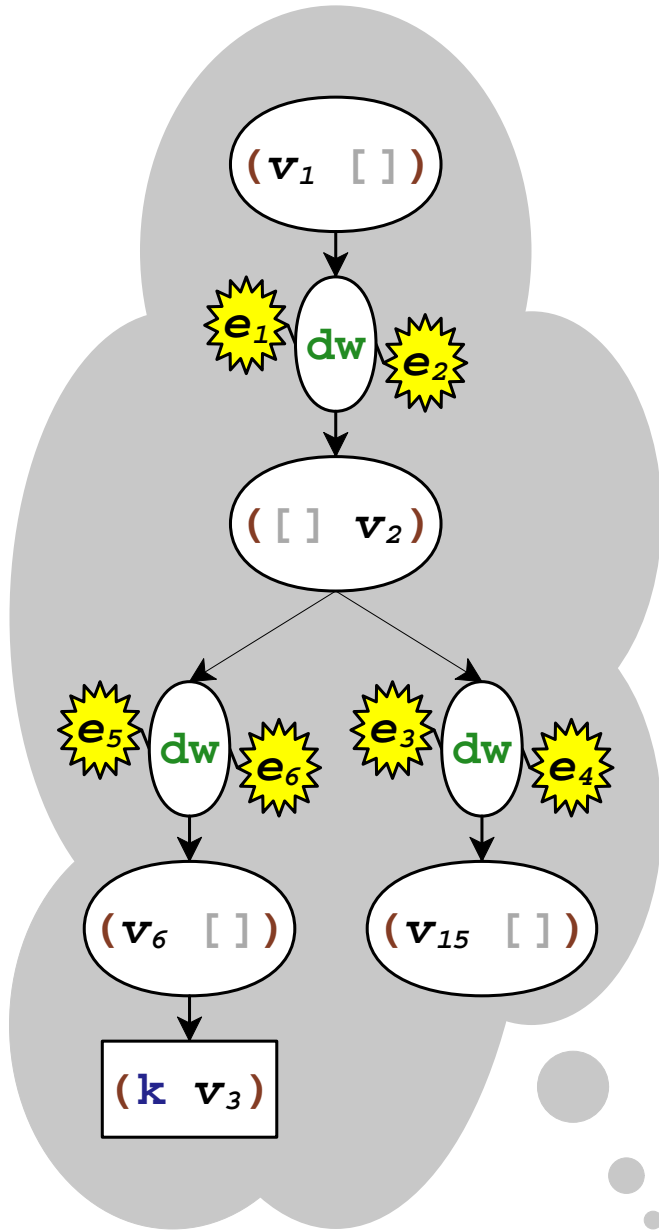
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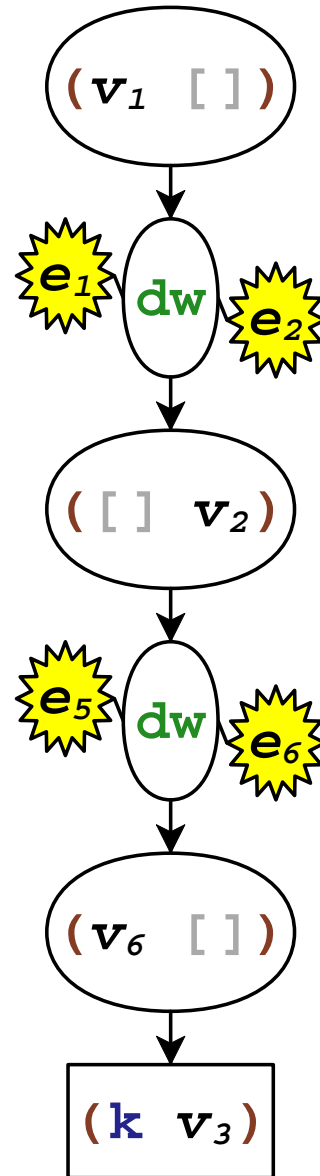
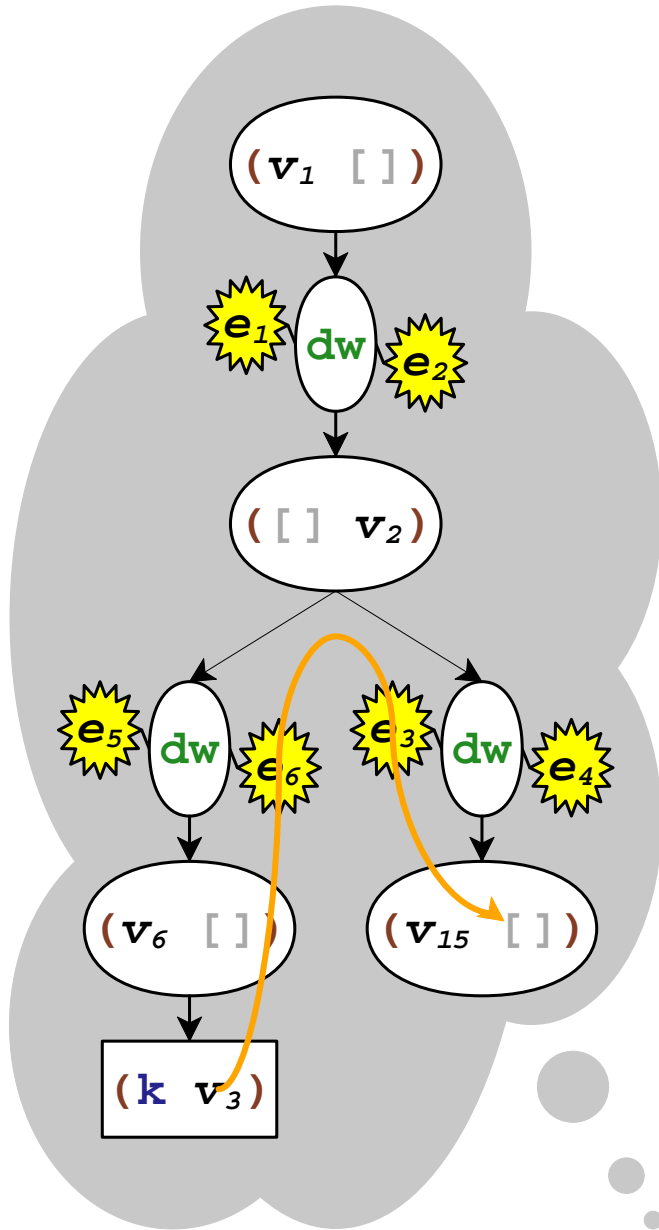
Dynamic Wind and Call/cc



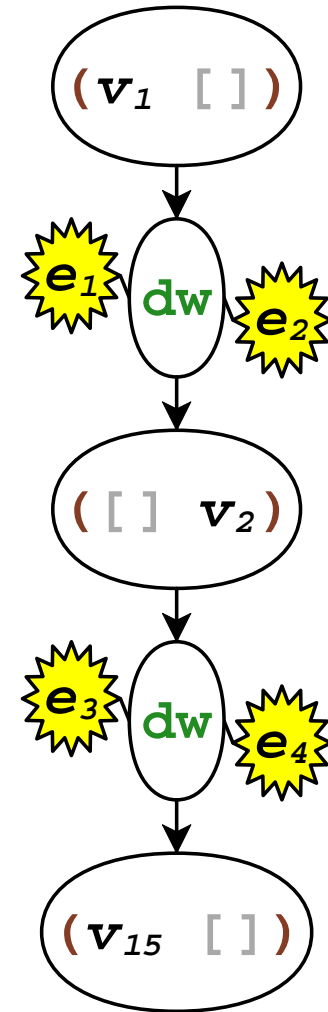
Dynamic Wind and Call/cc



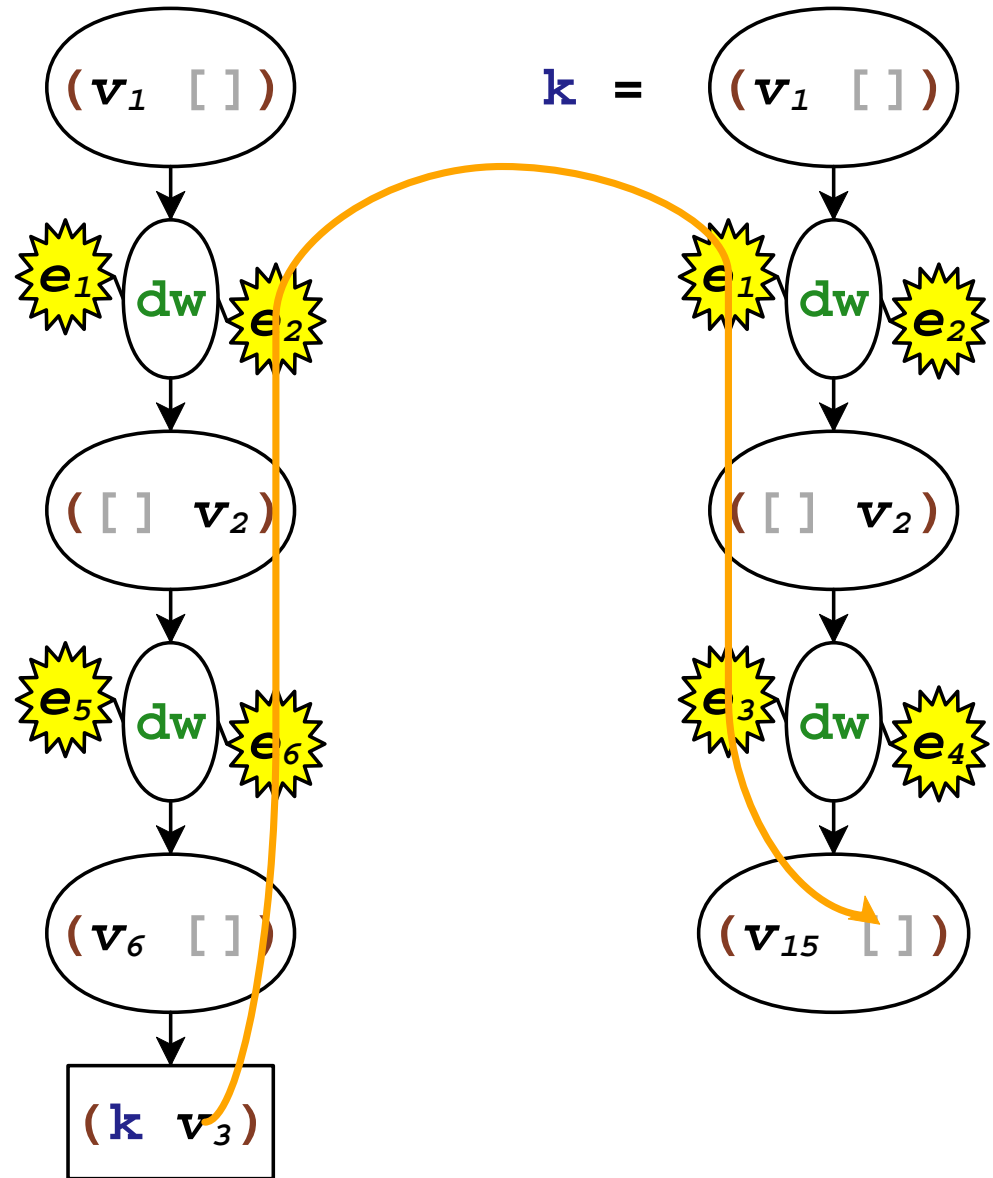
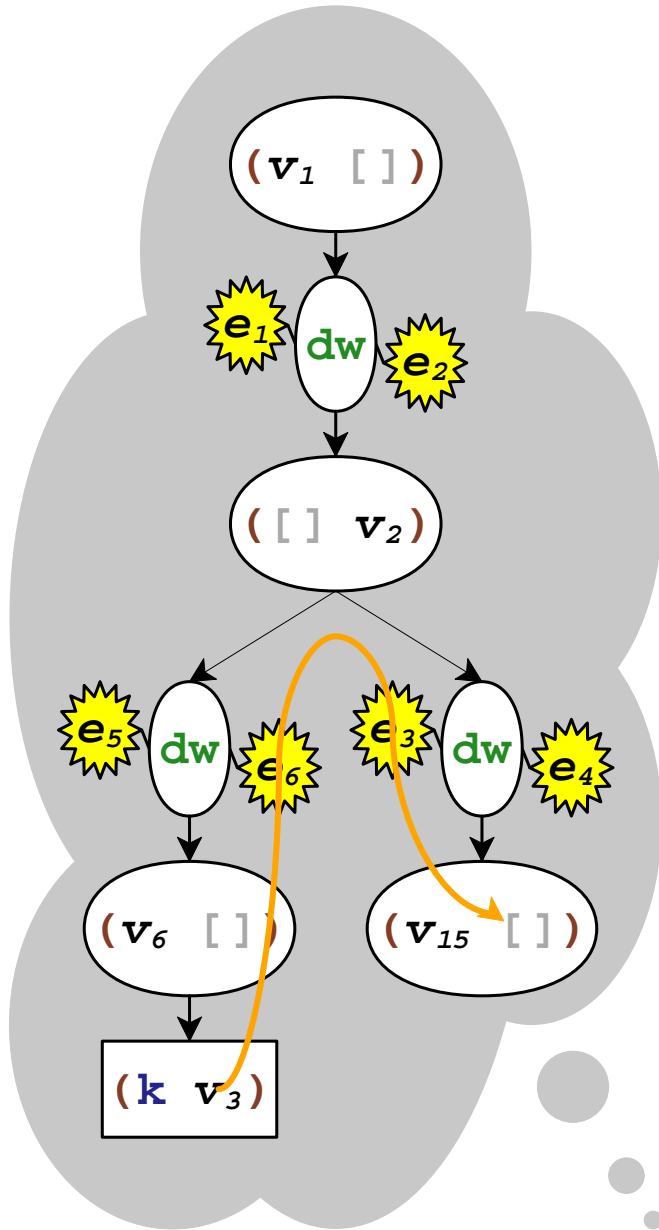
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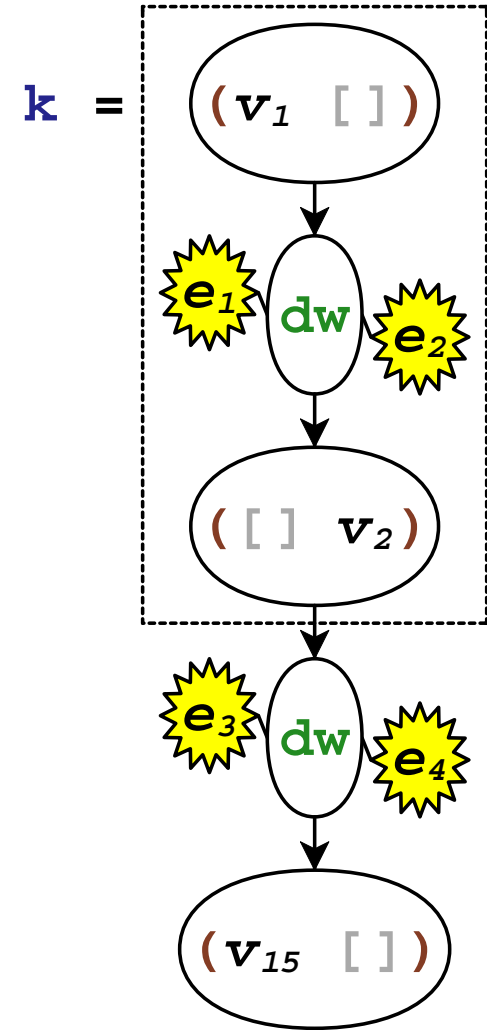
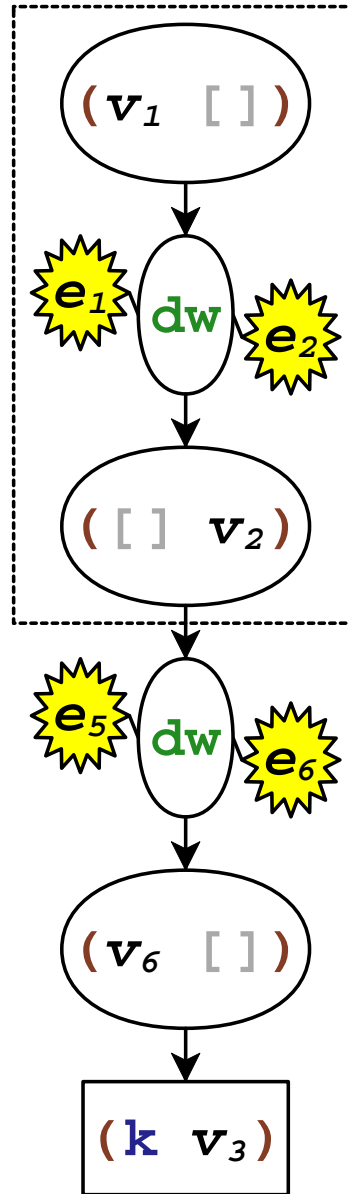
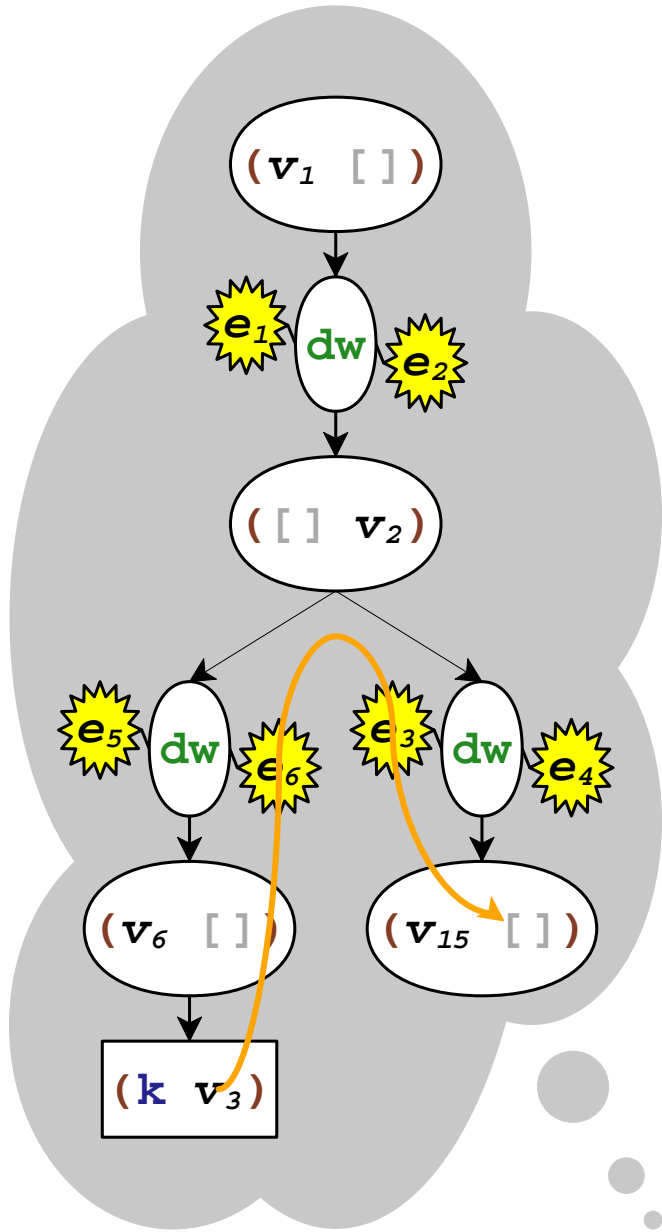
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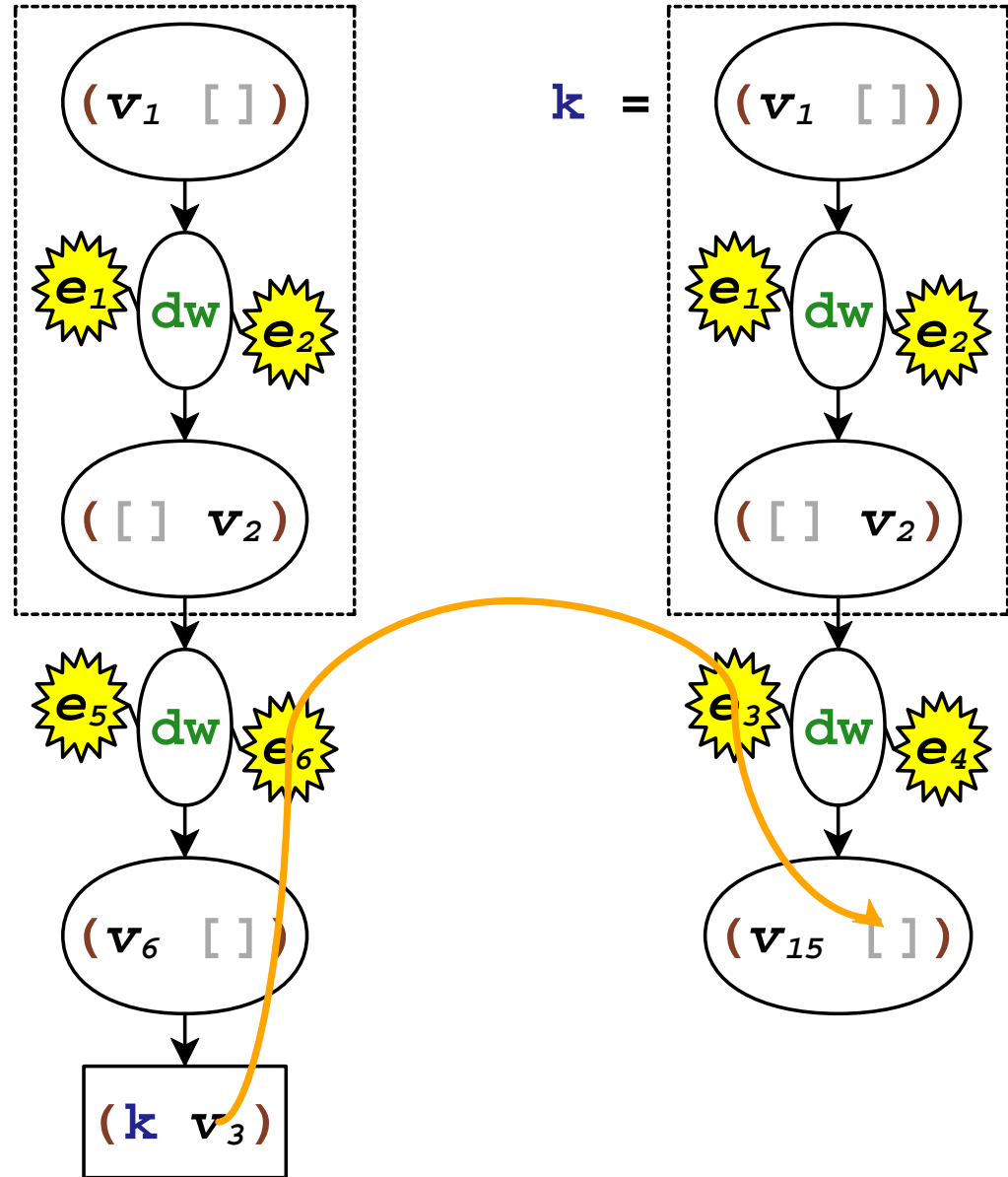
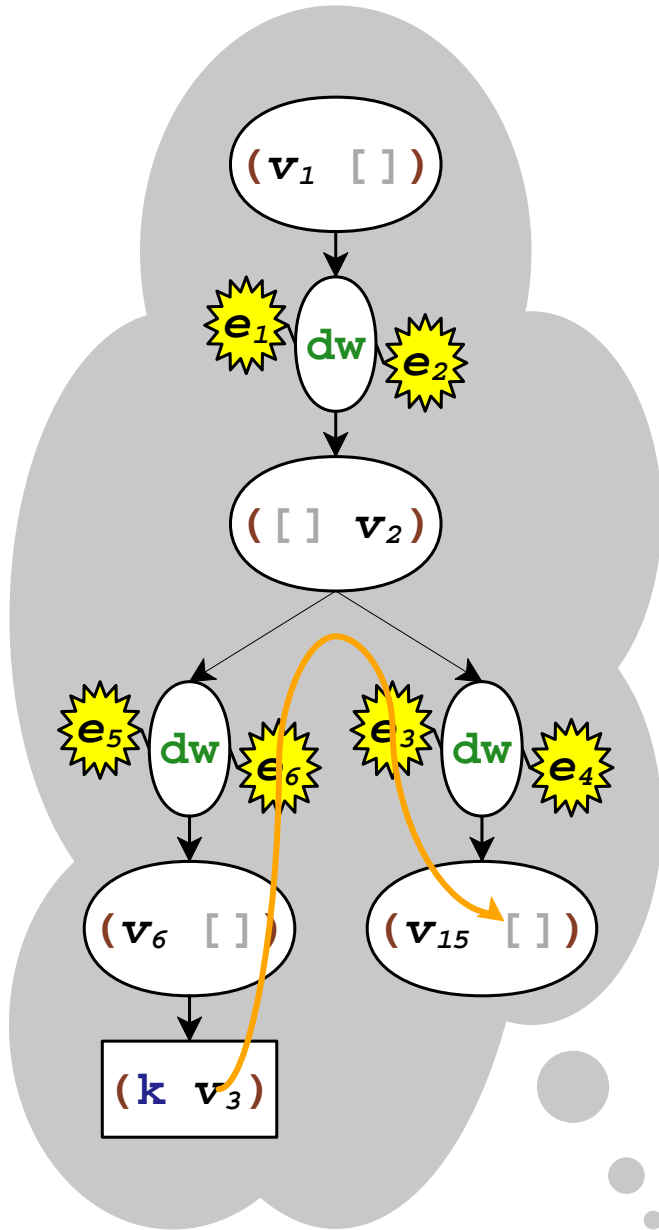
Dynamic Wind and Call/cc



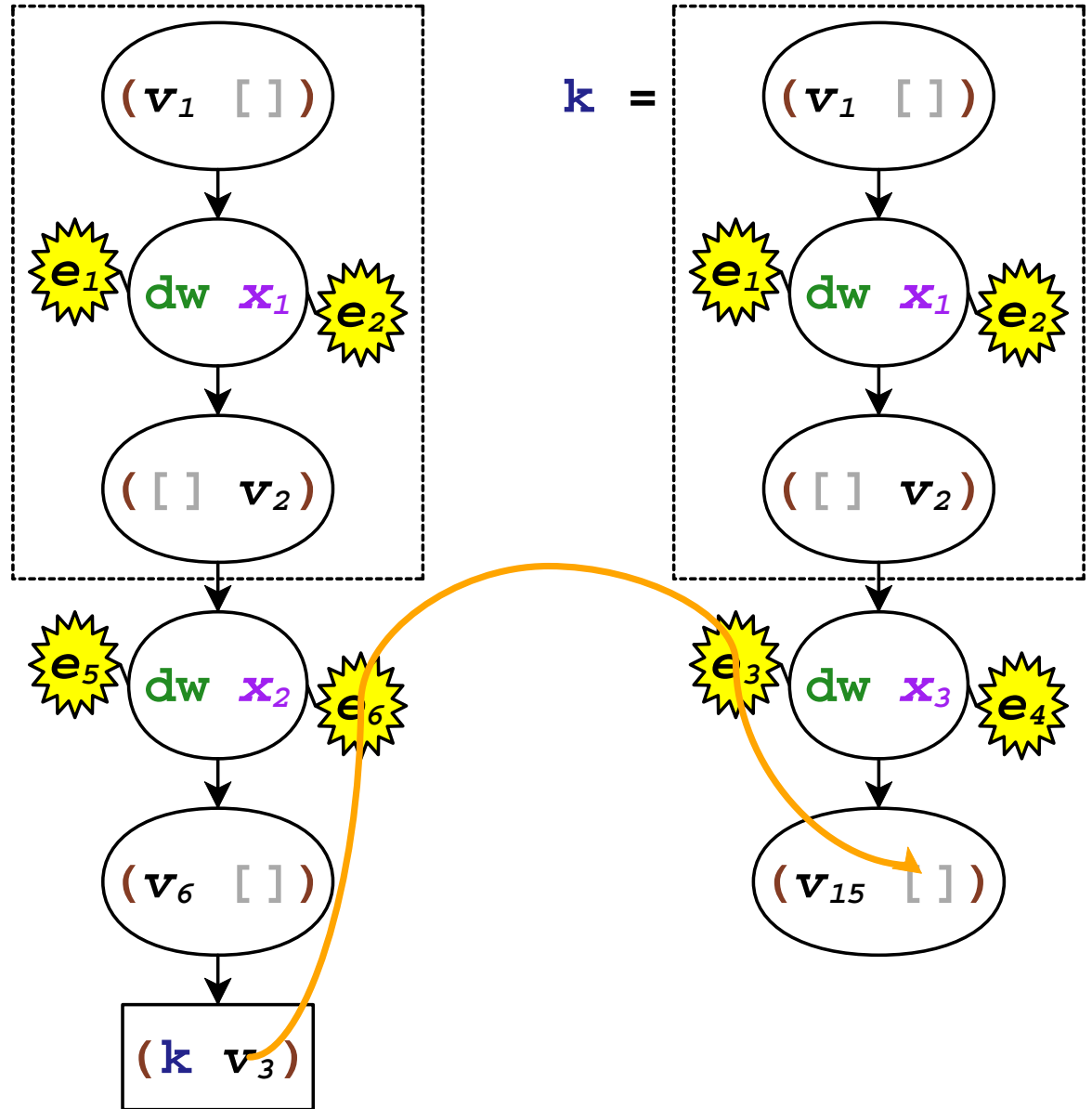
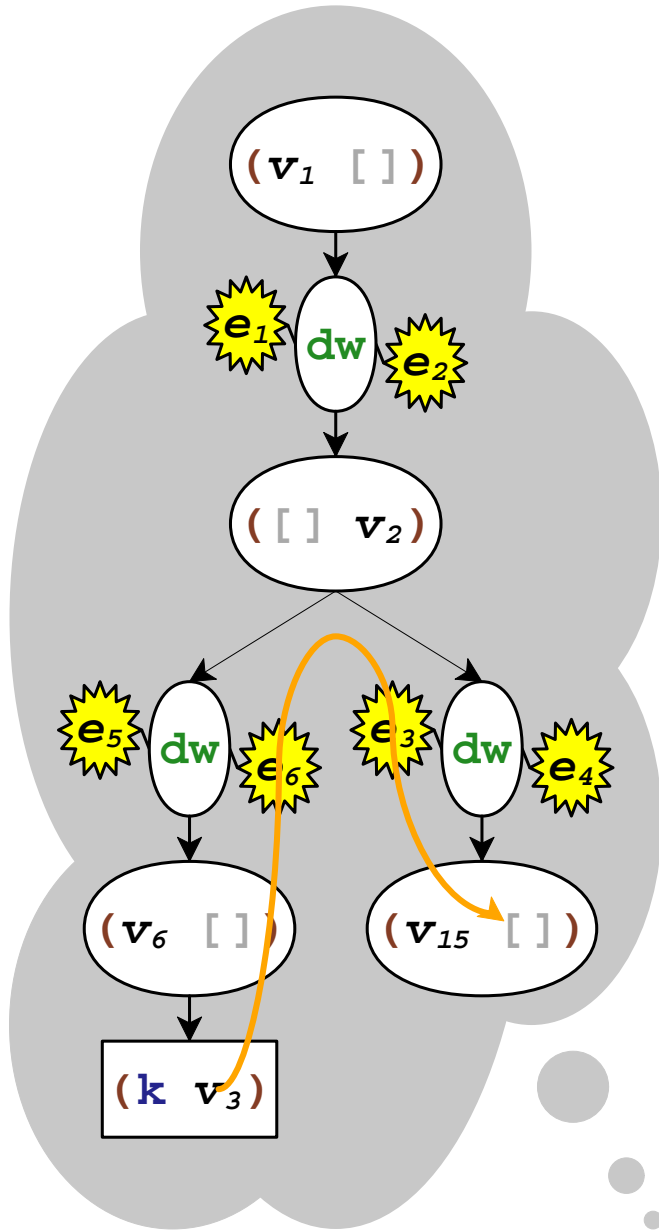
Dynamic Wind and Call/cc



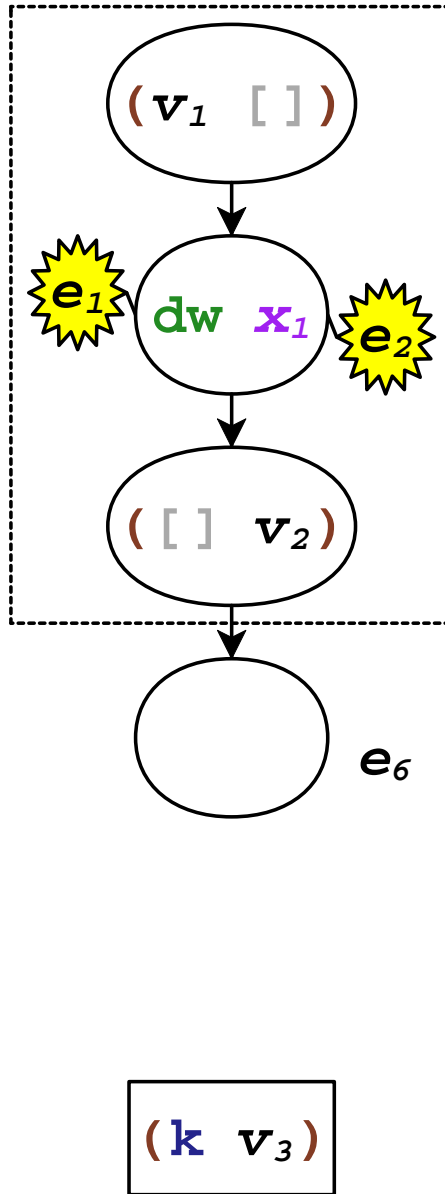
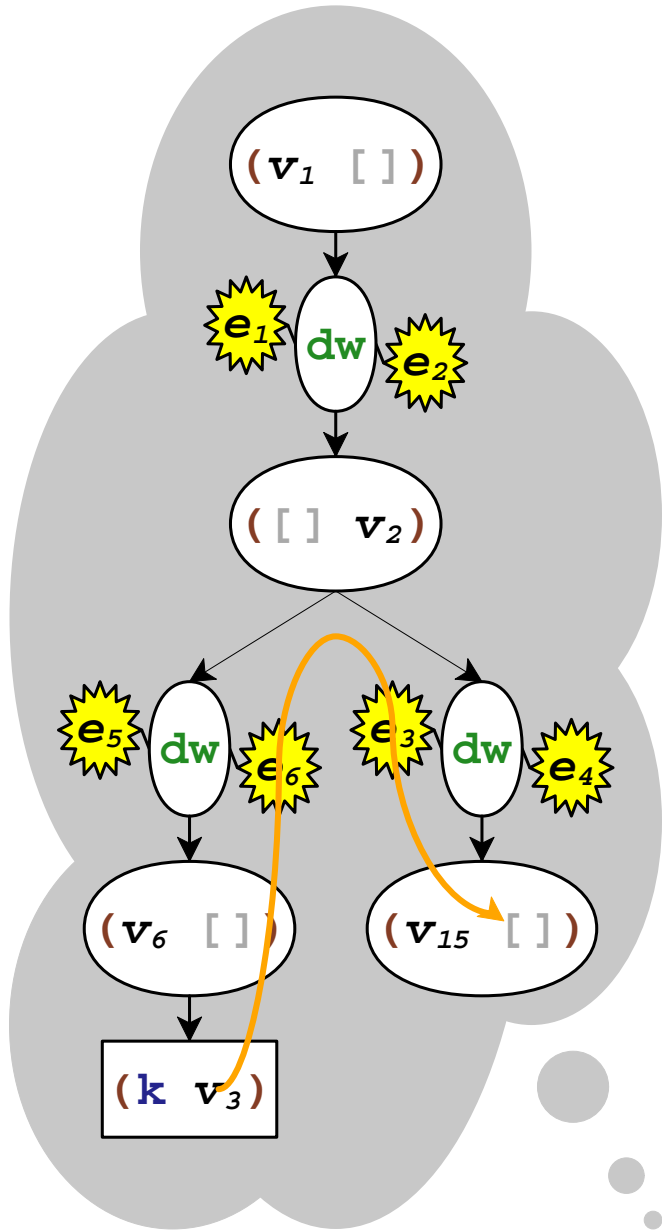
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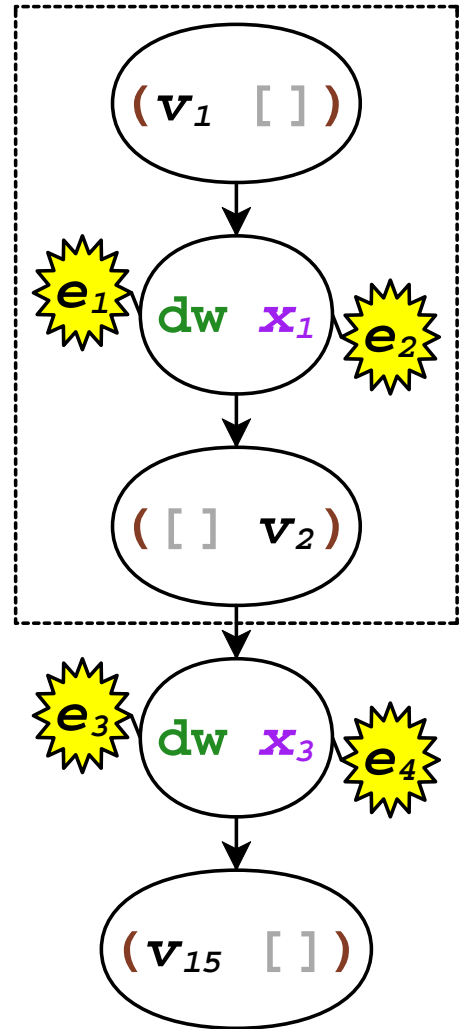
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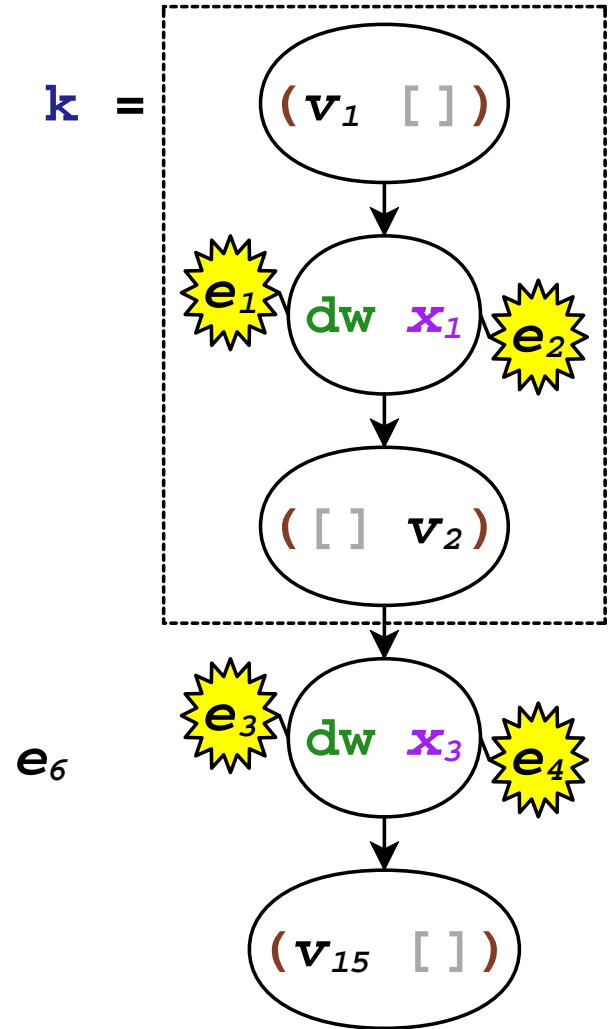
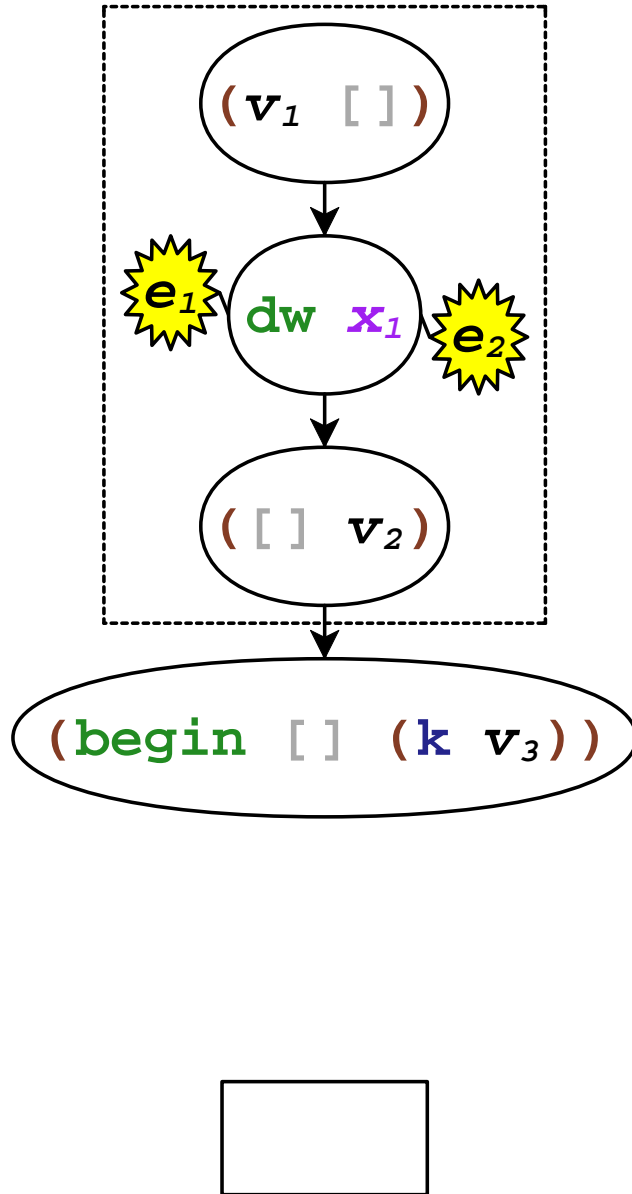
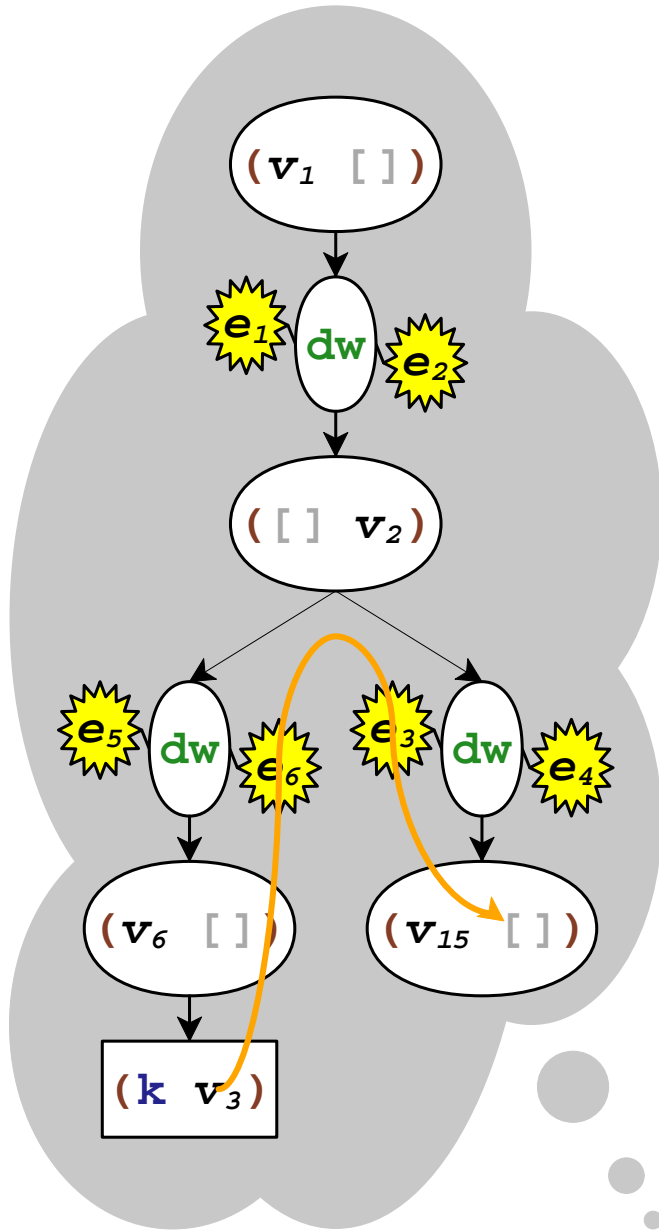
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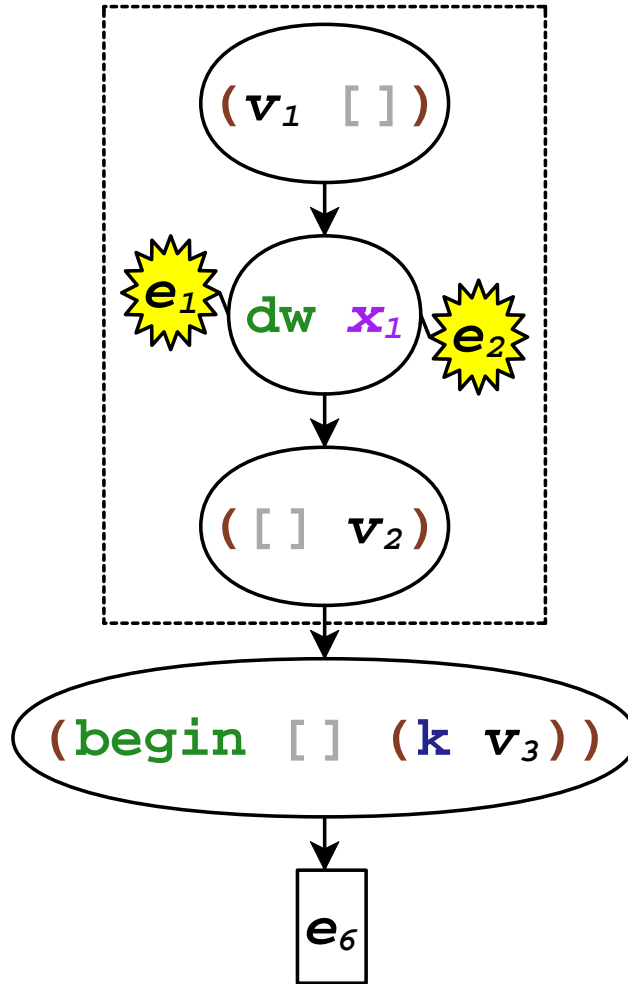
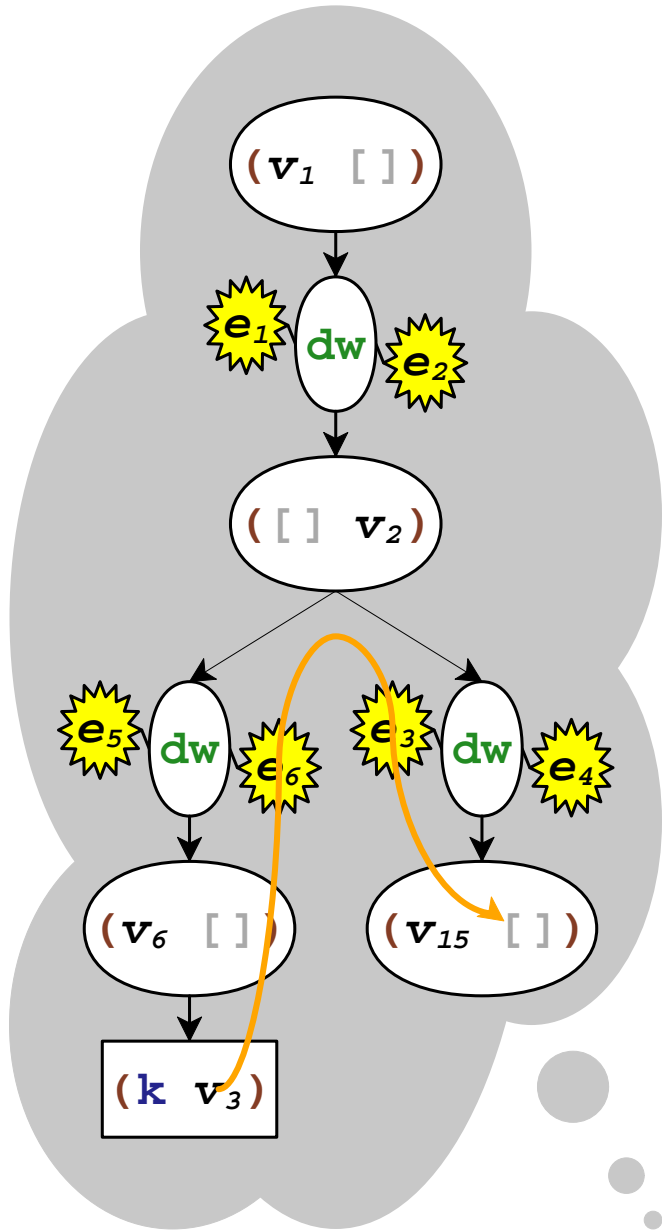
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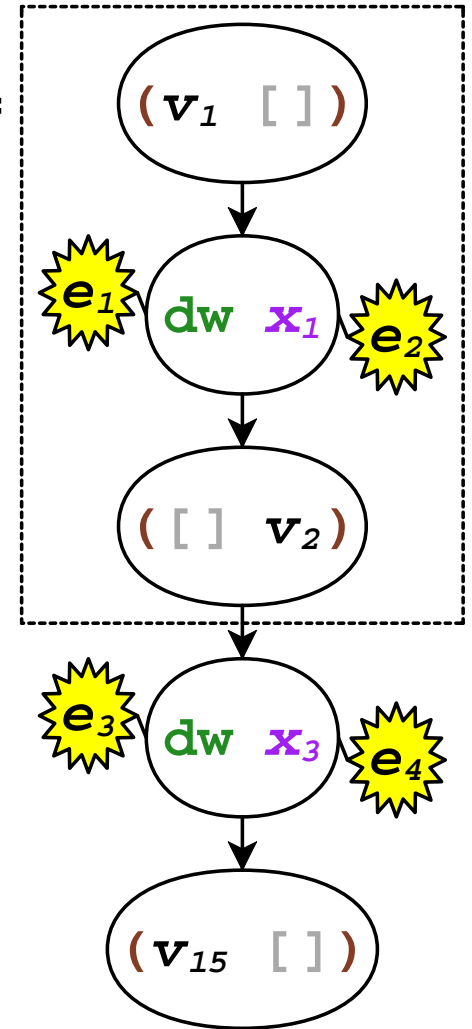
Dynamic Wind and Call/cc



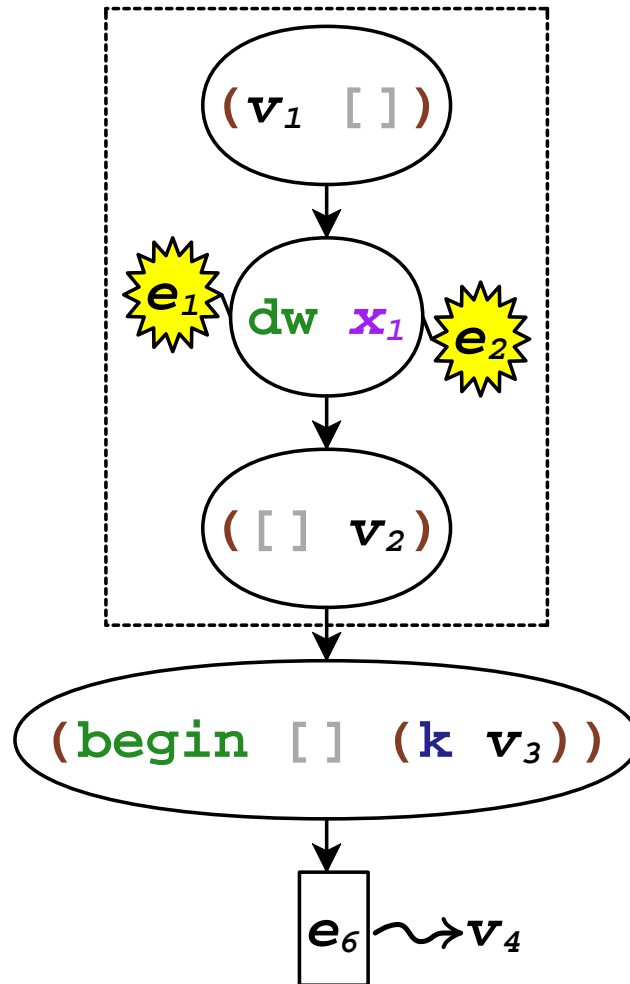
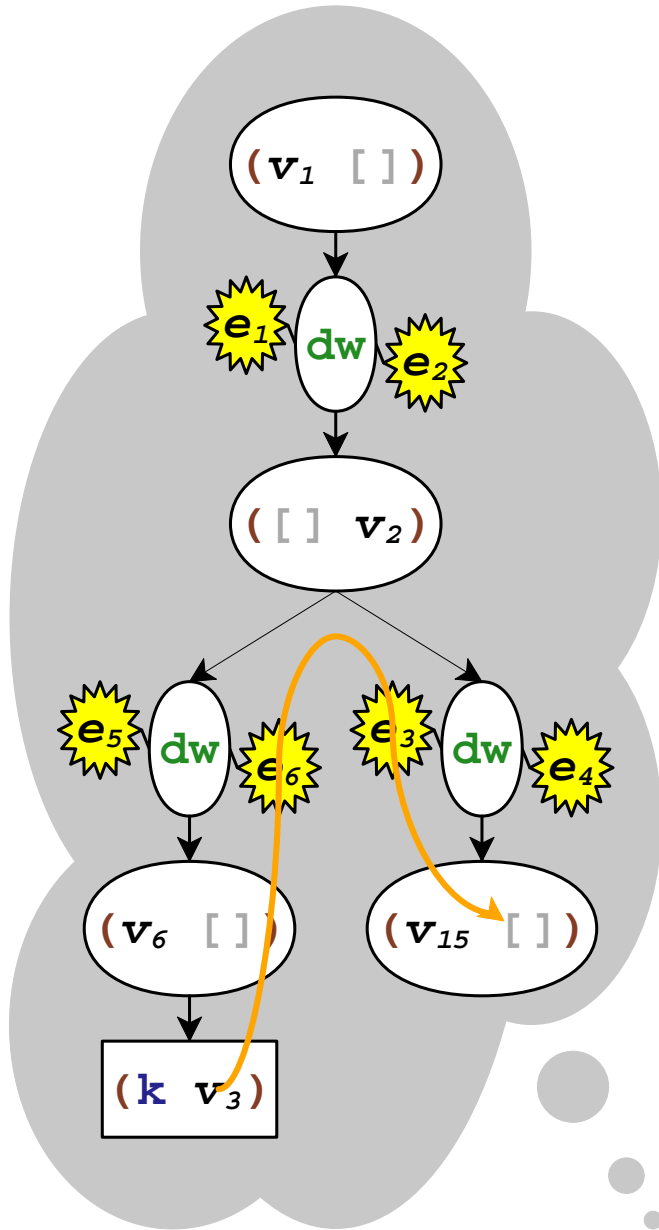
Dynamic Wind and Call/cc



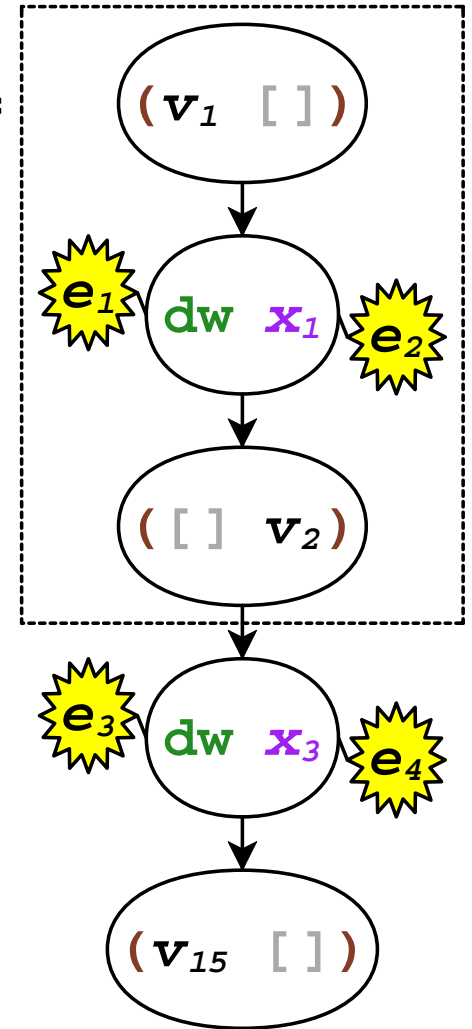
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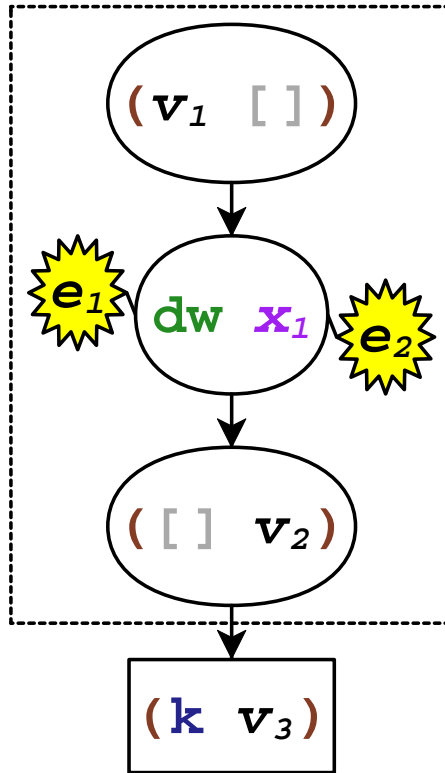
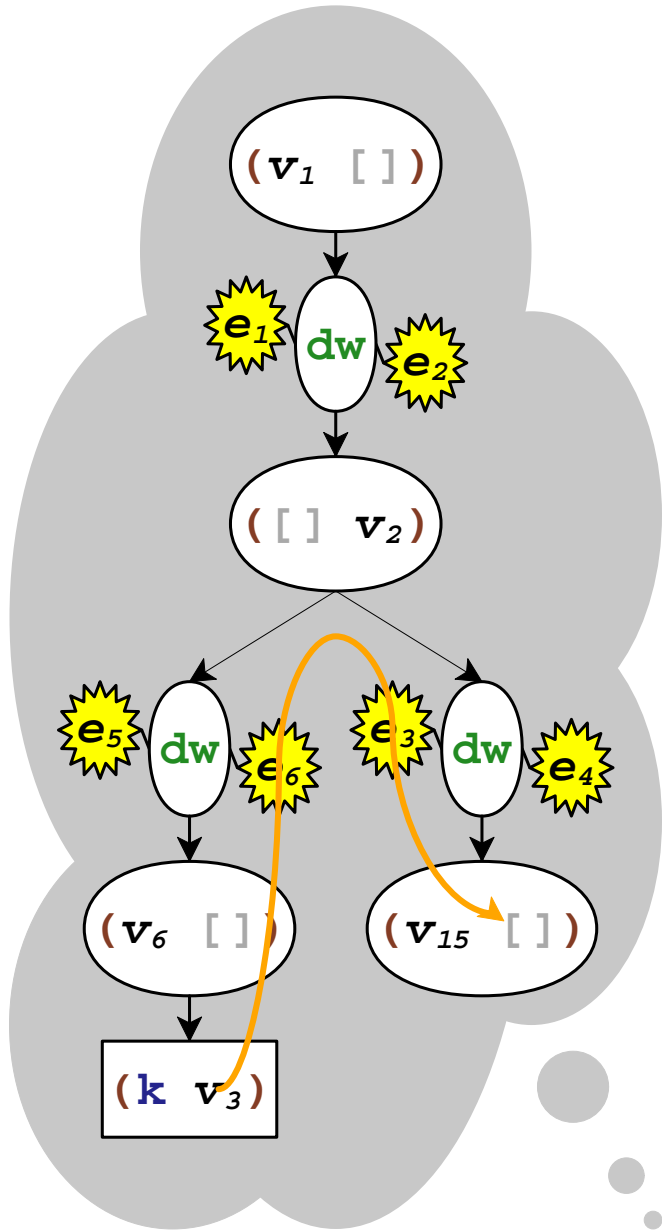
Dynamic Wind and Call/cc



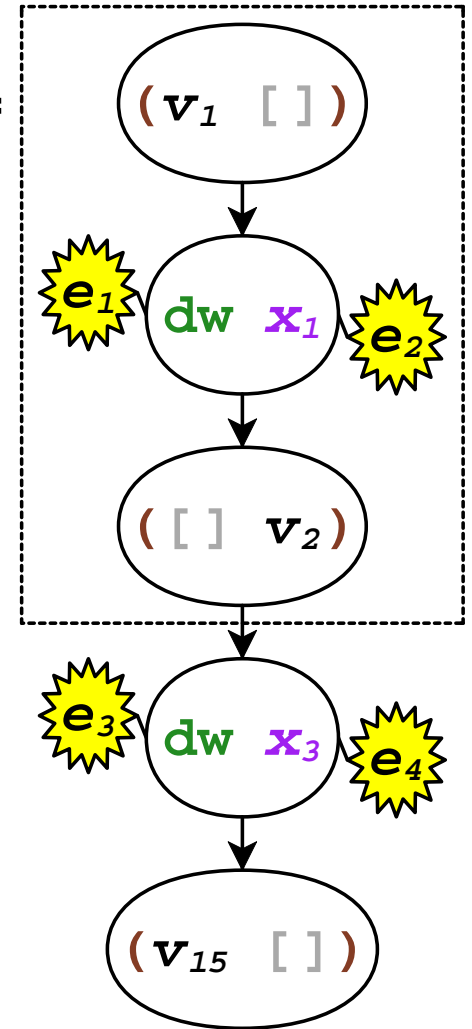
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Dynamic Wind and Call/cc



$k =$



Dynamic Wind Summary

- `dynamic-wind` generates `dw`
- `call/cc` detects sharing in continuation jumps
- capture `dw` thunks in `call/comp` and `call/cc`
- run post thunks in `abort`
- run pre thunks in continuation composition

To Continue...

- Read the paper

ICFP'07 Flatt, Yu, Findler, and Felleisen

- Run the Redex model:

`http://www.cs.utah.edu/plt/delim-cont/`

- Download the implementation:

`http://www.plt-scheme.org/`