

Chapter 1: introduction

our goal:

- ❖ get “feel” and terminology
- ❖ more depth, detail *later* in course
- ❖ approach:
 - use Internet as example

overview:

- ❖ what’s the Internet?
- ❖ what’s a protocol?
- ❖ network edge; hosts, access net, physical media
- ❖ network core: packet/circuit switching, Internet structure
- ❖ performance: loss, delay, throughput
- ❖ security
- ❖ protocol layers, service models
- ❖ history

Introduction 1-1

Chapter 1: roadmap

1.1 *what is the Internet?*

1.2 network edge

- end systems, access networks, links

1.3 network core

- packet switching, circuit switching, network structure

1.4 delay, loss, throughput in networks

1.5 protocol layers, service models

1.6 networks under attack: security

1.7 history

Introduction 1-2

What's the Internet: "nuts and bolts" view



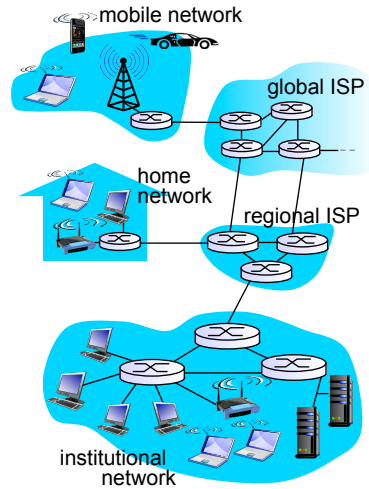
- ❖ millions of connected computing devices: *hosts*
 - running *network apps*



- ❖ *communication links*
 - fiber, copper, radio, satellite
 - transmission rate: *bandwidth*



- ❖ *routers*: forward packets (chunks of data)



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"Fun" internet appliances



IP picture frame
<http://www.ceiva.com/>



Web-enabled toaster + weather forecaster



Tweet-a-watt:
monitor energy use



Internet refrigerator



Slingbox: watch, control cable TV remotely

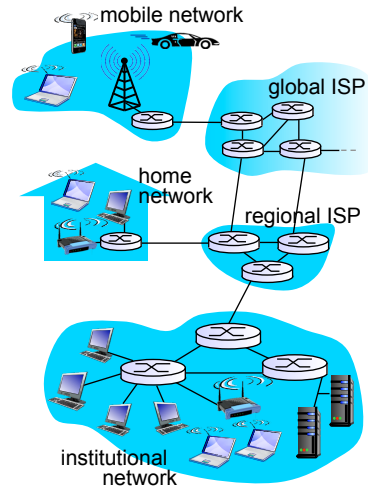


Internet phones

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What's the Internet: "nuts and bolts" view

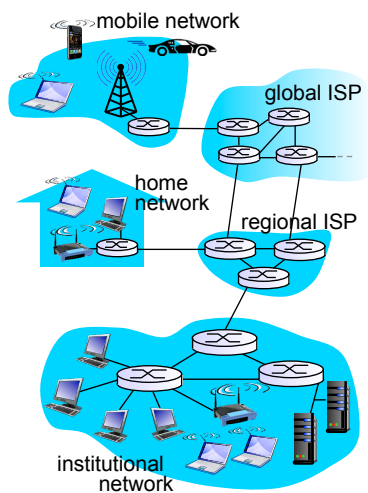
- ❖ **protocols** control sending, receiving of msgs
 - e.g., TCP, IP, HTTP, Skype, Ethernet
- ❖ **Internet: "network of networks"**
 - loosely hierarchical
 - public Internet versus private intranet
- ❖ **Internet standards**
 - RFC: Request for comments
 - IETF: Internet Engineering Task Force



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What's the Internet: a service view

- ❖ **communication infrastructure** enables distributed applications:
 - Web, VoIP, email, games, e-commerce, file sharing
- ❖ **communication services provided to apps:**
 - reliable data delivery from source to destination
 - "best effort" (unreliable) data delivery



Introduction 1-6

What's a protocol?

human protocols:

- ❖ “what’s the time?”
- ❖ “I have a question”
- ❖ introductions

... specific msgs sent

... specific actions taken when msgs received, or other events

network protocols:

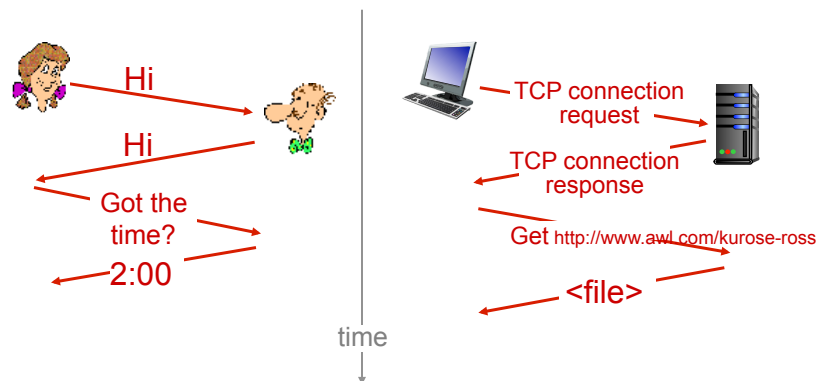
- ❖ machines rather than humans
- ❖ all communication activity in Internet governed by protocols

protocols define format, order of msgs sent and received among network entities, and actions taken on msg transmission, receipt

Introduction 1-7

What's a protocol?

a human protocol and a computer network protocol:



Q: other human protocols?

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Chapter 1: roadmap

1.1 what is the Internet?

1.2 network edge

- end systems, access networks, links

1.3 network core

- packet switching, circuit switching, network structure

1.4 delay, loss, throughput in networks

1.5 protocol layers, service models

1.6 networks under attack: security

1.7 history

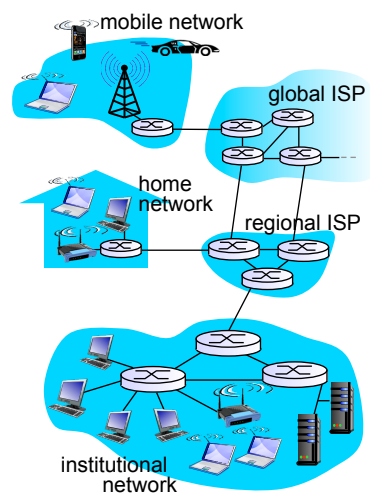
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A closer look at network structure:

❖ *network edge:*
applications and hosts

❖ *access networks, physical media:* wired, wireless communication links

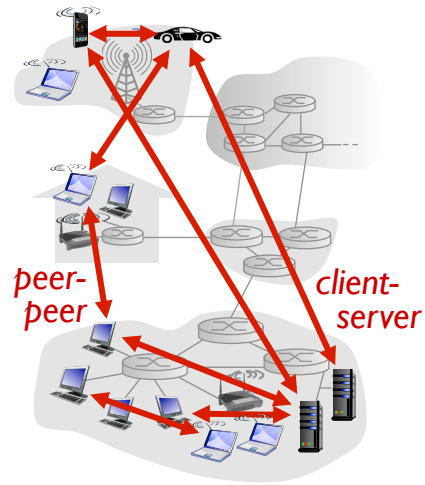
❖ *network core:*
▪ interconnected routers
▪ network of networks



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The network edge

- ❖ **end systems (hosts):**
 - run application programs
 - e.g. Web, email
 - at “edge of network”
- ❖ **client/server model**
 - client host requests, receives service from always-on server
 - e.g., Web browser/server; email client/server
- ❖ **peer-peer model:**
 - minimal (or no) use of dedicated servers
 - e.g. Skype, BitTorrent



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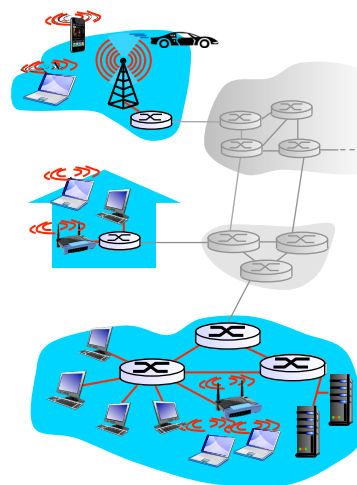
Access networks and physical media

Q: How to connect end systems to edge router?

- ❖ residential access nets
- ❖ institutional access networks (school, company)
- ❖ mobile access networks

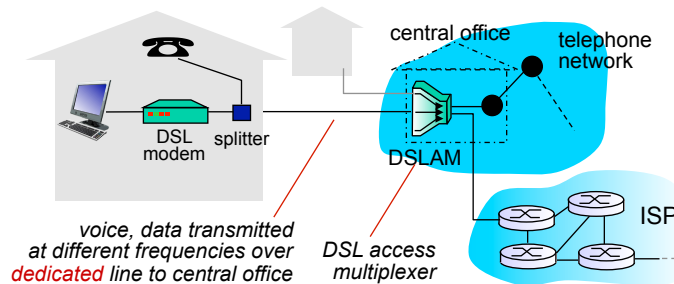
keep in mind:

- ❖ bandwidth (bits per second) of access network?
- ❖ shared or dedicated?



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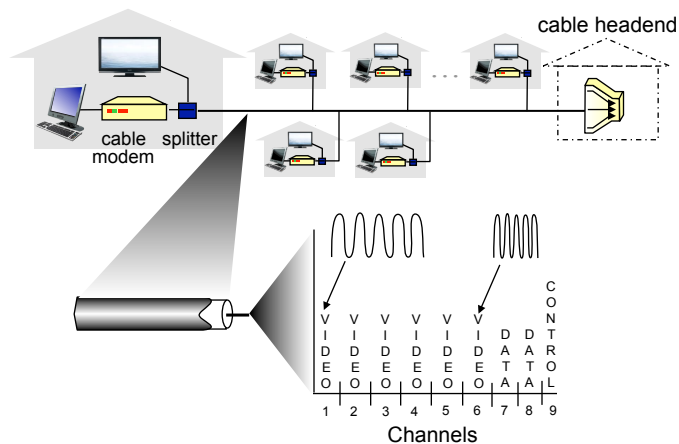
Access net: digital subscriber line (DSL)



- ❖ use **existing** telephone line to central office DSLAM
 - data over DSL phone line goes to Internet
 - voice over DSL phone line goes to telephone net
- ❖ < 2.5 Mbps upstream transmission rate (typically < 1 Mbps)
- ❖ < 24 Mbps downstream transmission rate (typically < 10 Mbps)

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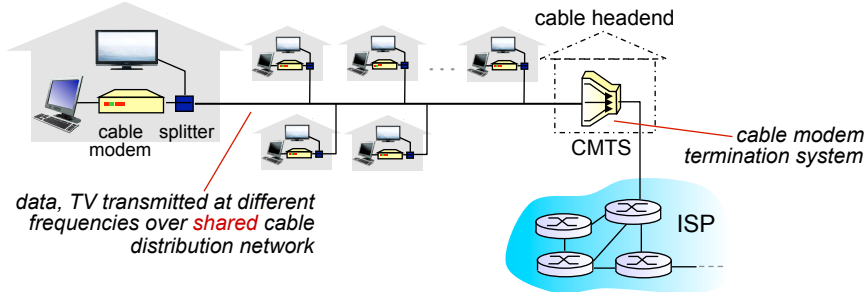
Access net: cable network



frequency division multiplexing: different channels transmitted in different frequency bands

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Access net: cable network



❖ HFC: hybrid fiber coax

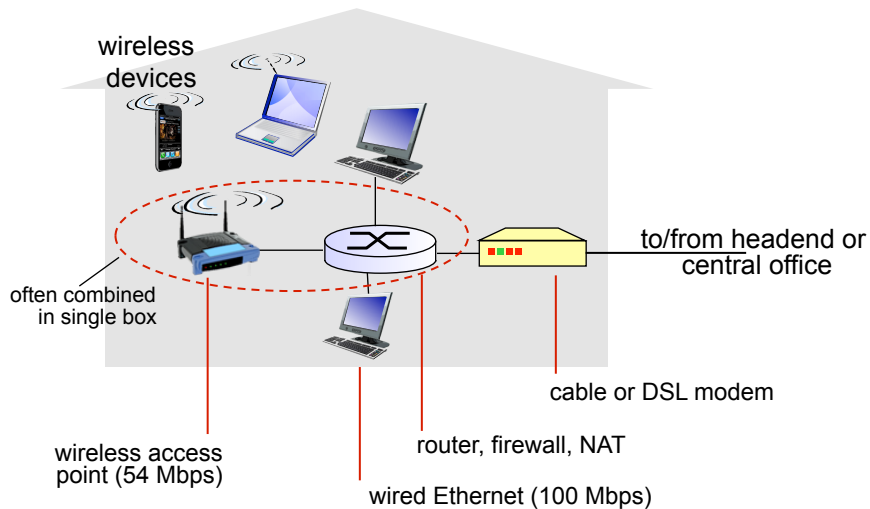
- asymmetric: up to 30Mbps downstream transmission rate, 2 Mbps upstream transmission rate

❖ network of cable, fiber attaches homes to ISP router

- homes *share access network* to cable headend
- unlike DSL, which has dedicated access to central office

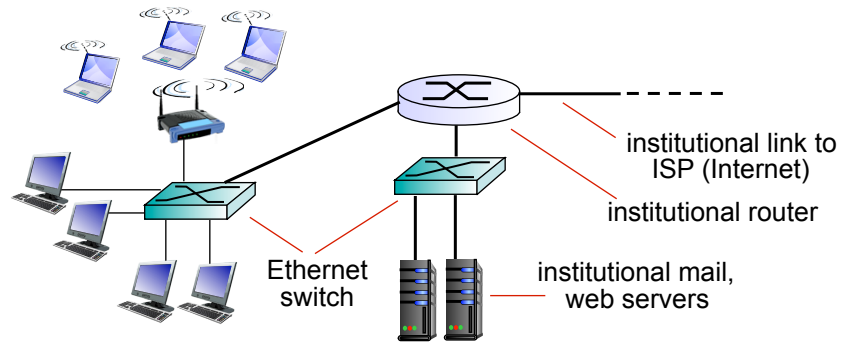
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Access net: home network



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Enterprise access networks (Ethernet)



- ❖ typically used in companies, universities, etc
- ❖ 10 Mbps, 100Mbps, 1Gbps, 10Gbps transmission rates
- ❖ today, end systems typically connect into Ethernet switch

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Wireless access networks

- ❖ shared *wireless* access network connects end system to router
 - via base station aka “access point”

wireless LANs:

- within building (100 ft)
- 802.11b/g (WiFi): 11, 54 Mbps transmission rate



to Internet

wide-area wireless access

- provided by telco (cellular) operator, 10's Km
- between 1 and 10 Mbps
- 3G, 4G: LTE, WiMax



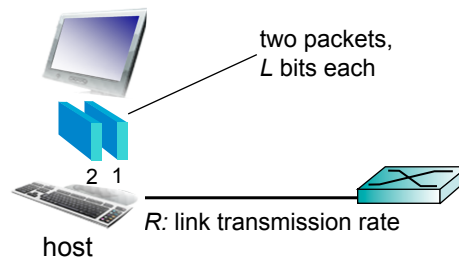
to Internet

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Host: sends packets of data

host sending function:

- ❖ takes application message
- ❖ breaks into smaller chunks, known as *packets*, of length L bits
- ❖ transmits packet into access network at *transmission rate R*
 - link transmission rate, aka link *capacity*, aka *link bandwidth*



$$\text{packet transmission delay} = \text{time needed to transmit } L\text{-bit packet into link} = \frac{L \text{ (bits)}}{R \text{ (bits/sec)}}$$

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Physical media

- ❖ **bit**: propagates between transmitter/receiver pairs
- ❖ **physical link**: what lies between transmitter & receiver
- ❖ **guided media**:
 - signals propagate in solid media: copper, fiber, coax
- ❖ **unguided media**:
 - signals propagate freely, e.g., radio

twisted pair (TP)

- ❖ two insulated copper wires
 - Category 5: 100 Mbps, 1 Gbps Ethernet
 - Category 6: 10Gbps



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Physical media: coax, fiber

coaxial cable:

- ❖ two concentric copper conductors
- ❖ bidirectional
- ❖ broadband:
 - multiple channels on cable
 - HFC



fiber optic cable:

- ❖ glass fiber carrying light pulses, each pulse a bit
- ❖ high-speed operation:
 - high-speed point-to-point transmission (e.g., 10's-100's Gpbs transmission rate)
- ❖ low error rate:
 - repeaters spaced far apart
 - immune to electromagnetic noise



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Physical media: radio

- ❖ signal carried in electromagnetic spectrum
- ❖ no physical "wire"
- ❖ bidirectional
- ❖ propagation environment effects:
 - reflection
 - obstruction by objects
 - interference

radio link types:

- ❖ **terrestrial microwave**
 - e.g. up to 45 Mbps channels
- ❖ **LAN** (e.g., WiFi)
 - 11 Mbps, 54 Mbps
- ❖ **wide-area** (e.g., cellular)
 - 3G cellular: ~ 1 Mbps
- ❖ **satellite**
 - Kbps to 45Mbps channel (or multiple smaller channels)
 - 270 msec end-end delay
 - geosynchronous versus low altitude

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Chapter 1: roadmap

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- packet switching, circuit switching, network structure

1.4 delay, loss, throughput in networks

1.5 protocol layers, service models

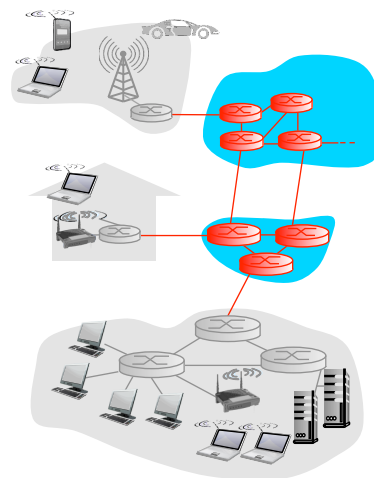
1.6 networks under attack: security

1.7 history

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The network core

- ❖ mesh of interconnected routers
- ❖ packet-switching: hosts break application-layer messages into *packets*
 - forward packets from one router to the next, across links on path from source to destination
 - each packet transmitted at full link capacity



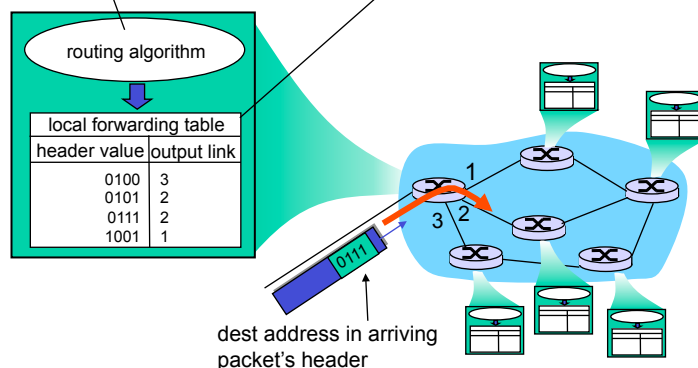
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Two key network-layer functions

routing: determines source-destination route taken by packets

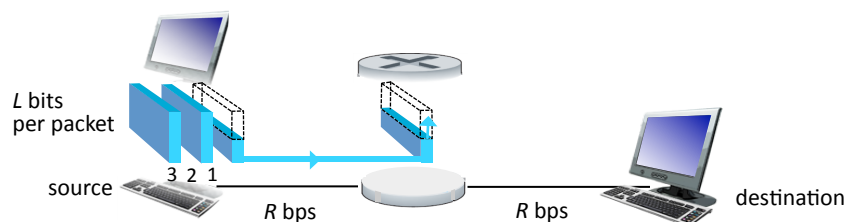
- routing algorithms

forwarding: move packets from router's input to appropriate router output



Network Layer 4-25

Packet-switching: store-and-forward



- ❖ takes L/R seconds to transmit (push out) L -bit packet into link at R bps
- ❖ **store and forward:** entire packet must arrive at router before it can be transmitted on next link
- ❖ end-end delay = $2L/R$ (assuming zero propagation delay)

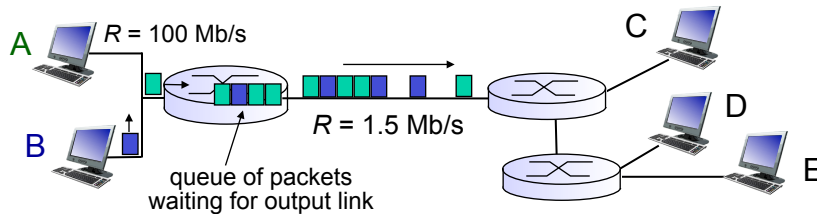
one-hop numerical example:

- $L = 7.5$ Mbits
- $R = 1.5$ Mbps
- one-hop transmission delay = 5 sec

} more on delay shortly ...

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Packet Switching: queuing delay, loss



resource contention:

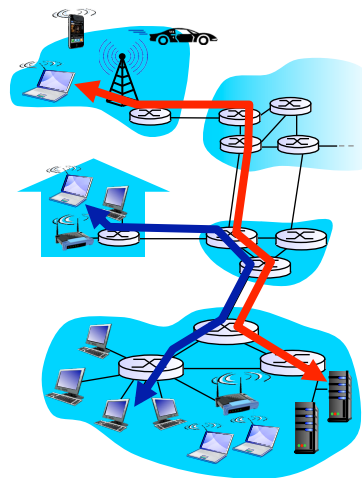
- ❖ aggregate resource demand (use of transmission link) can exceed amount available
- ❖ **congestion:**
 - packets will queue, wait for link use
 - packets can be dropped (lost) if no memory to store them

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Alternative core: circuit switching

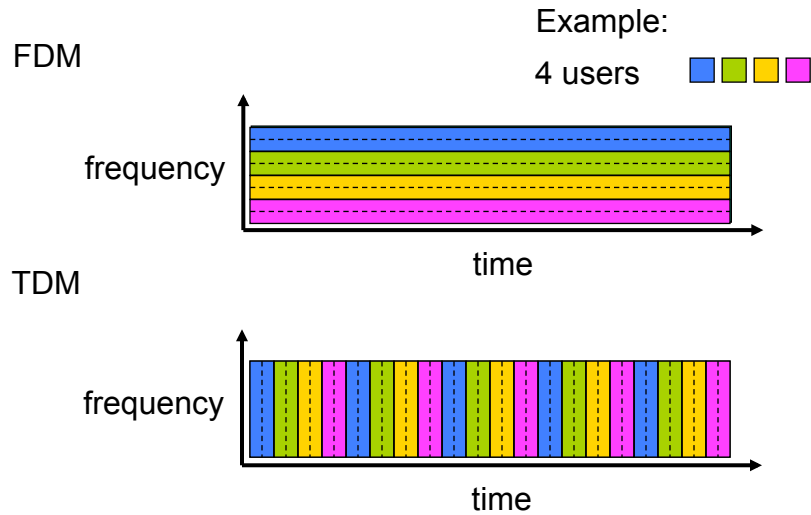
end-end resources allocated to, reserved for, "call" between source, dest.:

- ❖ dedicated resources: no sharing
 - circuit-like (guaranteed) performance
- ❖ resource piece idle if not used by owning call (*no sharing*)
- ❖ dividing link bandwidth into "pieces"
 - frequency division
 - time division



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Circuit switching: FDM versus TDM



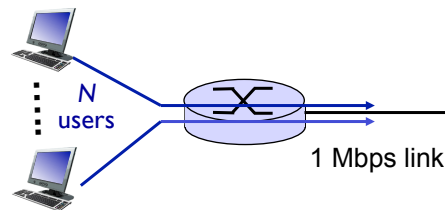
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Packet switching versus circuit switching

packet switching allows more users to use network!

example:

- 1 Mb/s link
- each user:
 - 100 kb/s when “active”
 - active 10% of time



❖ *circuit-switching:*

- 10 users

❖ *packet switching:*

- with 35 users, probability > 10 active at same time is less than .0004

Q: how did we get value 0.0004?

Q: what happens if > 35 users ?

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Packet switching versus circuit switching

is packet switching a “slam dunk winner?”

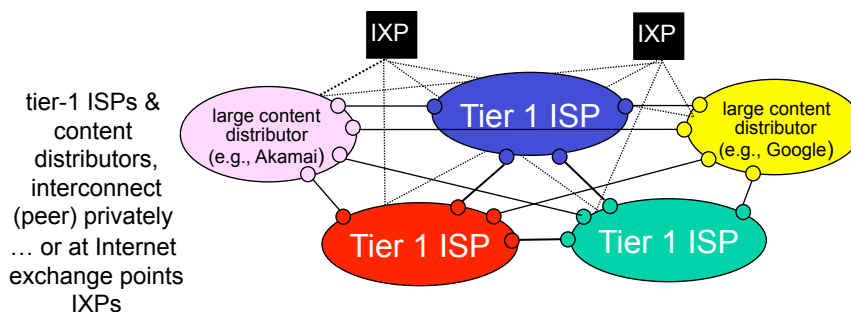
- ❖ great for bursty data
 - resource sharing
 - simpler, no call setup
- ❖ **excessive congestion possible:** packet delay and loss
 - protocols needed for reliable data transfer, congestion control
- ❖ **Q: How to provide circuit-like behavior?**
 - bandwidth guarantees needed for audio/video apps
 - still an unsolved problem (chapter 7)

Q: human analogies of reserved resources (circuit switching) versus on-demand allocation (packet-switching)?

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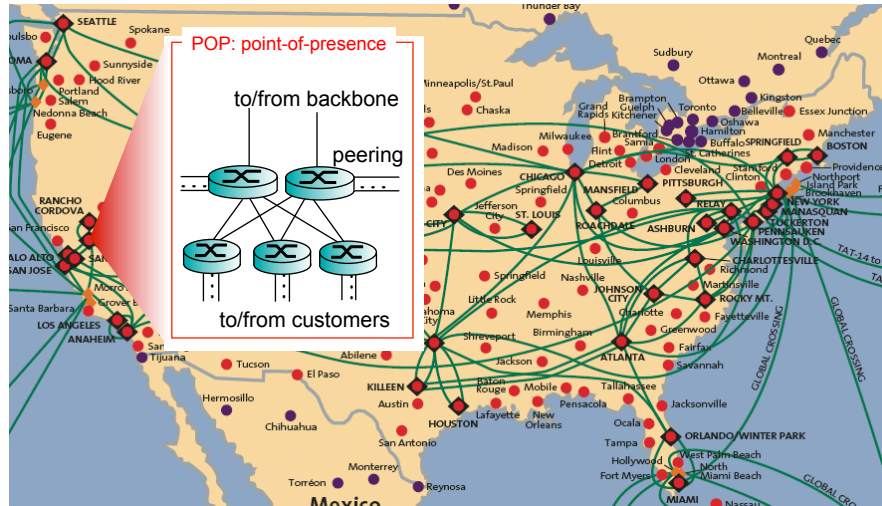
Internet structure: network of networks

- ❖ roughly hierarchical
- ❖ at center: small # of well-connected large networks
 - “**tier-1**” **commercial ISPs** (e.g., Verizon, Sprint, AT&T, Qwest, Level3), national & international coverage
 - **large content distributors** (Google, Akamai, Microsoft)
 - treat each other as equals (no charges)



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Tier-I ISP: e.g., Sprint

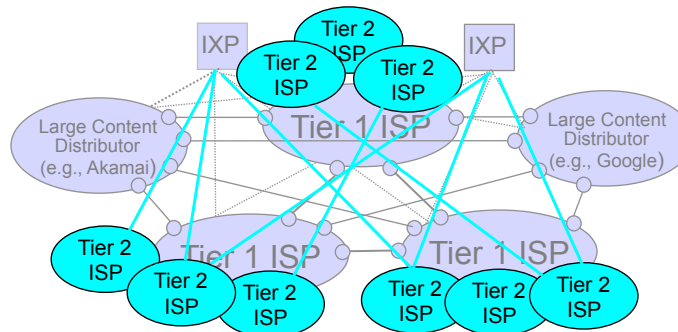


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Internet structure: network of networks

“tier-2” ISPs: smaller (often regional) ISPs

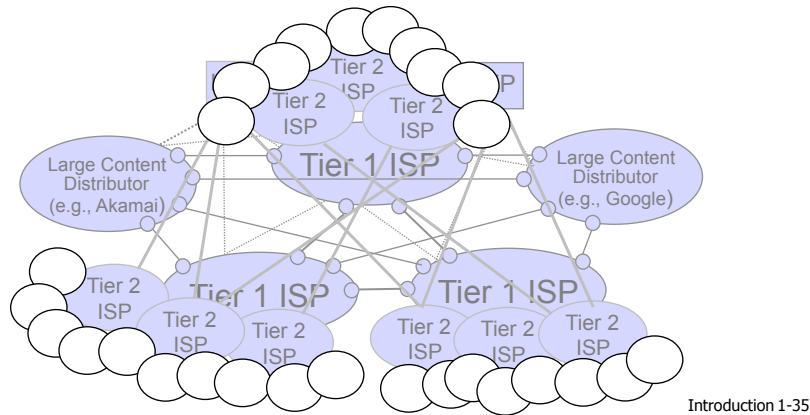
- connect to one or more tier-1 (*provider*) ISPs
 - each tier-1 has many tier-2 *customer nets*
 - tier 2 pays tier 1 provider
- tier-2 nets sometimes peer directly with each other (bypassing tier 1) , or at IXP



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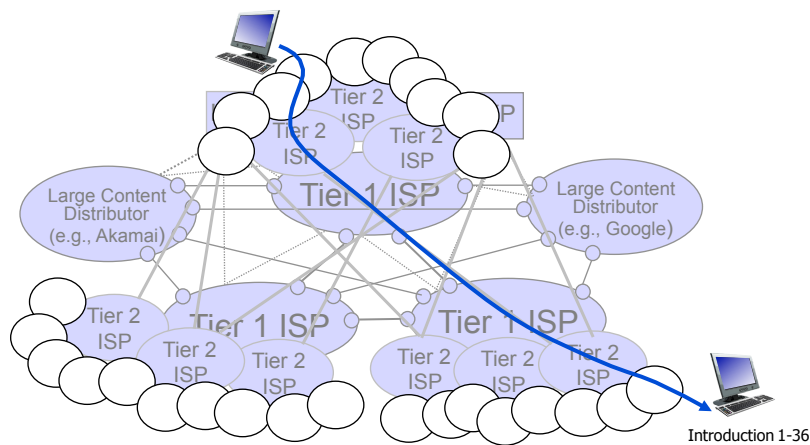
Internet structure: network of networks

- ❖ “tier-3” ISPs, local ISPs
- ❖ customer of tier 1 or tier 2 network
 - last hop (“access”) network (closest to end systems)



Internet structure: network of networks

- ❖ a packet passes through *many* networks from source host to destination host



Chapter 1: roadmap

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1.5 protocol layers, service models

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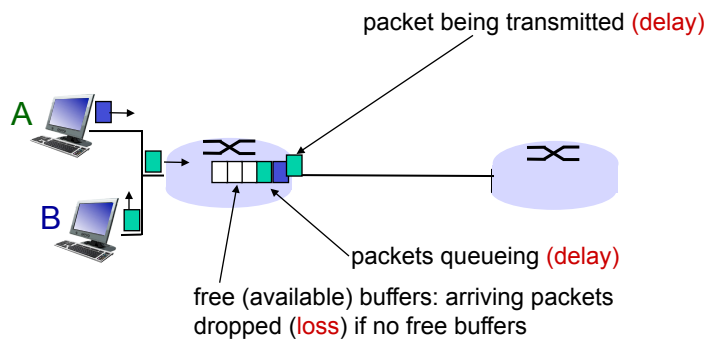
1.7 history

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How do loss and delay occur?

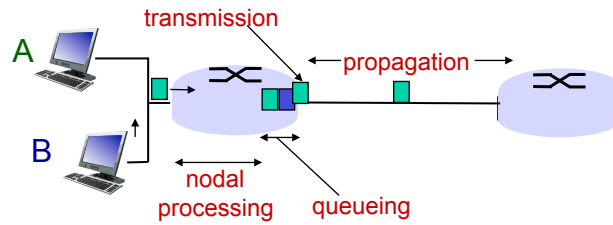
packets *queue* in router buffers

- ❖ packet arrival rate to link (temporarily) exceeds output link capacity
- ❖ packets queue, wait for turn



Introduction 1-38

Four sources of packet delay



$$d_{\text{nodal}} = d_{\text{proc}} + d_{\text{queue}} + d_{\text{trans}} + d_{\text{prop}}$$

d_{proc} : nodal processing

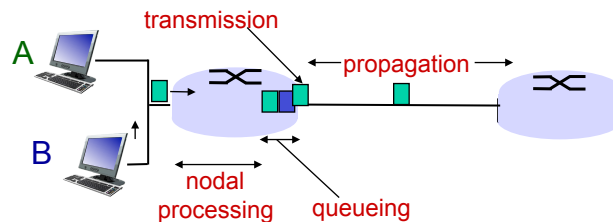
- check bit errors
- determine output link
- typically < msec

d_{queue} : queueing delay

- time waiting at output link for transmission
- depends on congestion level of router

Introduction 1-39

Four sources of packet delay



$$d_{\text{nodal}} = d_{\text{proc}} + d_{\text{queue}} + d_{\text{trans}} + d_{\text{prop}}$$

d_{trans} : transmission delay:

- L: packet length (bits)
- R: link bandwidth (bps)
- $d_{\text{trans}} = L/R$

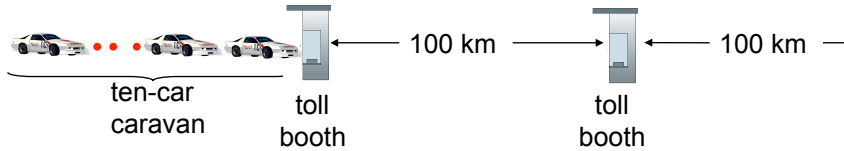
d_{prop} : propagation delay:

- d: length of physical link
- s: propagation speed in medium ($\sim 2 \times 10^8$ m/sec)
- $d_{\text{prop}} = d/s$

d_{trans} and d_{prop} very different

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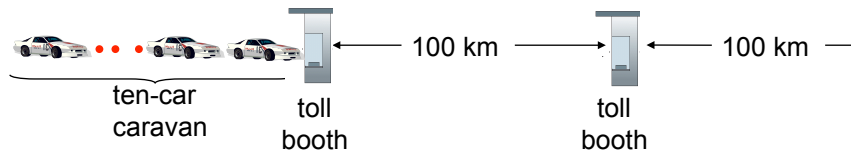
Caravan analogy



- ❖ cars “propagate” at 100 km/hr
- ❖ toll booth takes 12 sec to service car (transmission time)
- ❖ car~bit; caravan ~ packet
- ❖ **Q: How long until caravan is lined up before 2nd toll booth?**
- time to “push” entire caravan through toll booth onto highway = $12 \times 10 = 120$ sec
- time for last car to propagate from 1st to 2nd toll booth: $100 \text{ km} / (100 \text{ km/hr}) = 1$ hr
- **A: 62 minutes**

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Caravan analogy (more)

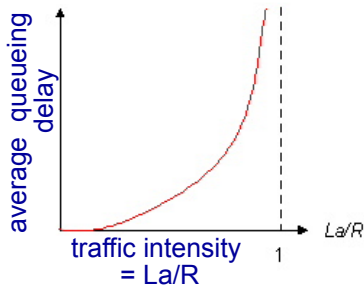


- ❖ cars now “propagate” at 1000 km/hr
- ❖ toll booth now takes 1 min to service a car
- ❖ **Q: Will cars arrive to 2nd booth before all cars serviced at first booth?**
 - **A: Yes!** after 7 min, 1st car arrives at second booth; three cars still at 1st booth.
 - first bit of packet can arrive at 2nd router before packet is fully transmitted at 1st router! (see Ethernet applet at AVL Web site)

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Queueing delay (revisited)

- ❖ R : link bandwidth (bps)
- ❖ L : packet length (bits)
- ❖ a : average packet arrival rate



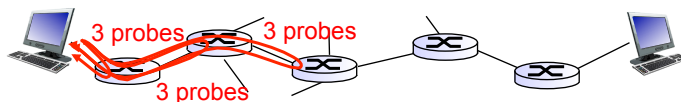
- ❖ $La/R \sim 0$: avg. queueing delay small
- ❖ $La/R \rightarrow 1$: avg. queueing delay large
- ❖ $La/R > 1$: more “work” arriving than can be serviced, average delay infinite!



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“Real” Internet delays and routes

- ❖ what do “real” Internet delay & loss look like?
- ❖ `traceroute` program: provides delay measurement from source to router along end-end Internet path towards destination. For all i :
 - sends three packets that will reach router i on path towards destination
 - router i will return packets to sender
 - sender times interval between transmission and reply.



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“Real” Internet delays, routes

traceroute: gaia.cs.umass.edu to www.eurecom.fr

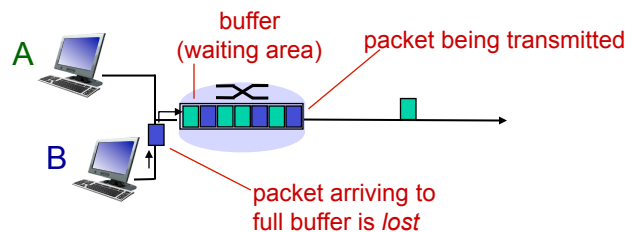
3 delay measurements from
gaia.cs.umass.edu to cs-gw.umass.edu

1 cs-gw (128.119.240.254) 1 ms 1 ms 2 ms
 2 border1-rt-fa5-1-0.gw.umass.edu (128.119.3.145) 1 ms 1 ms 2 ms
 3 cht-vbns.gw.umass.edu (128.119.3.130) 6 ms 5 ms 5 ms
 4 jn1-at1-0-0-19.wor.vbns.net (204.147.132.129) 16 ms 11 ms 13 ms
 5 jn1-so7-0-0-0.wae.vbns.net (204.147.136.136) 21 ms 18 ms 18 ms
 6 abilene-vbns.abilene.ucaid.edu (198.32.11.9) 22 ms 18 ms 22 ms
 7 nycm-wash.abilene.ucaid.edu (198.32.8.46) 22 ms 22 ms 22 ms
 8 62.40.103.253 (62.40.103.253) 104 ms 109 ms 106 ms ← trans-oceanic link
 9 de2-1.de1.de.geant.net (62.40.96.129) 109 ms 102 ms 104 ms
 10 de.fr1.fr.geant.net (62.40.96.50) 113 ms 121 ms 114 ms
 11 renater-gw.fr1.fr.geant.net (62.40.103.54) 112 ms 114 ms 112 ms
 12 nio-n2.cssi.renater.fr (193.51.206.13) 111 ms 114 ms 116 ms
 13 nice.cssi.renater.fr (195.220.98.102) 123 ms 125 ms 124 ms
 14 r3t2-nice.cssi.renater.fr (195.220.98.110) 126 ms 126 ms 124 ms
 15 eurecom-valbonne.r3t2.ft.net (193.48.50.54) 135 ms 128 ms 133 ms
 16 194.214.211.25 (194.214.211.25) 126 ms 128 ms 126 ms
 17 *** ← * means no response (probe lost, router not replying)
 18 ***
 19 fantasia.eurecom.fr (193.55.113.142) 132 ms 128 ms 136 ms

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Packet loss

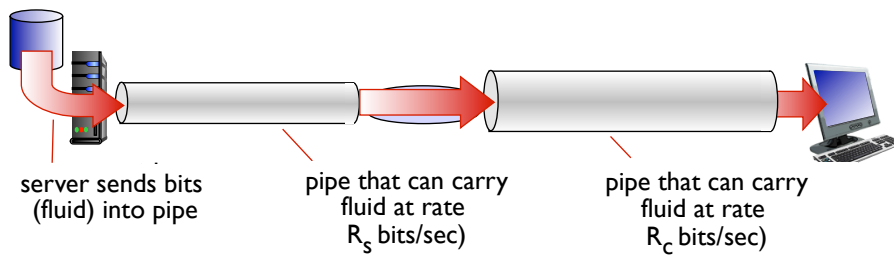
- ❖ queue (aka buffer) preceding link in buffer has finite capacity
- ❖ packet arriving to full queue dropped (aka lost)
- ❖ lost packet may be retransmitted by previous node, by source end system, or not at all



Introduction 1-46

Throughput

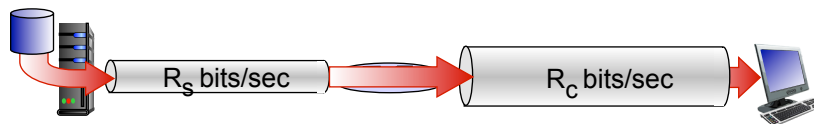
- ❖ **throughput**: rate (bits/time unit) at which bits transferred between sender/receiver
 - **instantaneous**: rate at given point in time
 - **average**: rate over longer period of time



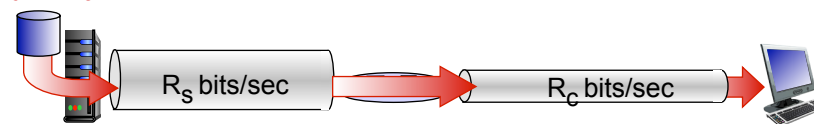
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Throughput (more)

- ❖ $R_s < R_c$ What is average end-end throughput?



- ❖ $R_s > R_c$ What is average end-end throughput?

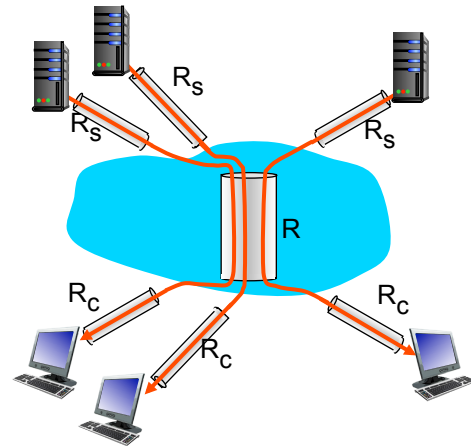


bottleneck link
link on end-end path that constrains end-end throughput

Introduction 1-48

Throughput: Internet scenario

- ❖ per-connection end-end throughput: $\min(R_c, R_s, R/10)$
- ❖ in practice: R_c or R_s is often bottleneck



10 connections (fairly) share backbone bottleneck link R bits/sec

Introduction 1-49

Chapter 1: roadmap

- 1.1 what is the Internet?
- 1.2 network edge
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- 1.3 network core
 - packet switching, circuit switching, network structure
- 1.4 delay, loss, throughput in networks
- 1.5 protocol layers, service models
- 1.6 networks under attack: security
- 1.7 history

Introduction 1-50

Protocol “layers”

*Networks are complex,
with many “pieces”:*

- hosts
- routers
- links of various media
- applications
- protocols
- hardware, software

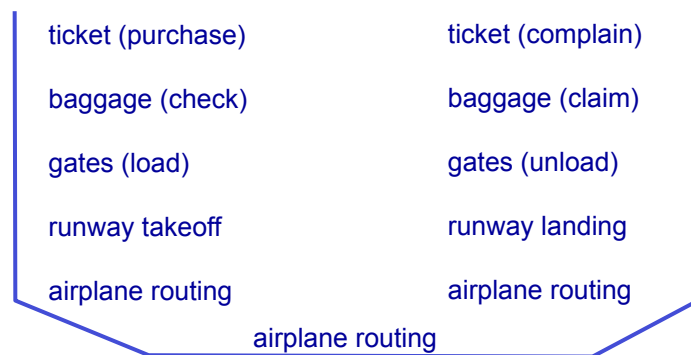
Question:

is there any hope of
organizing structure of
network?

.... or at least our
discussion of networks?

Introduction 1-51

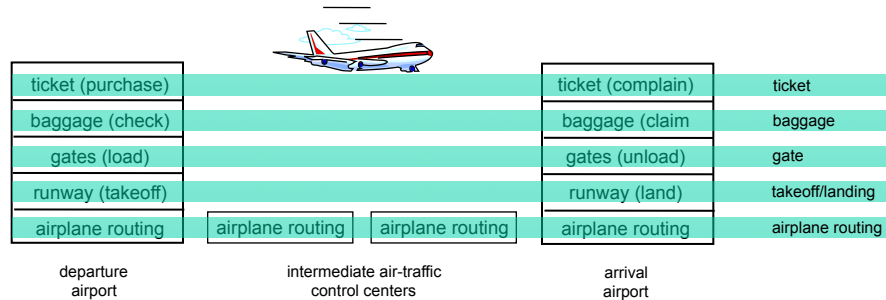
Organization of air travel



❖ a series of steps

Introduction 1-52

Layering of airline functionality



layers: each layer implements a service

- via its own internal-layer actions
- relying on services provided by layer below

Introduction 1-53

Why layering?

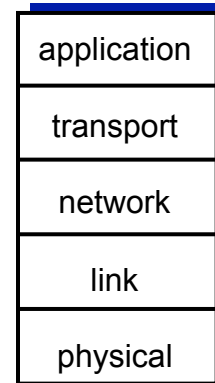
dealing with complex systems:

- ❖ explicit structure allows identification, relationship of complex system's pieces
 - layered *reference model* for discussion
- ❖ modularization eases maintenance, updating of system
 - change of implementation of layer's service transparent to rest of system
 - e.g., change in gate procedure doesn't affect rest of system
- ❖ layering considered harmful?

Introduction 1-54

Internet protocol stack

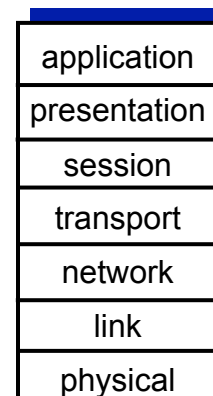
- ❖ **application**: supporting network applications
 - FTP, SMTP, HTTP
- ❖ **transport**: process-process data transfer
 - TCP, UDP
- ❖ **network**: routing of datagrams from source to destination
 - IP, routing protocols
- ❖ **link**: data transfer between neighboring network elements
 - Ethernet, 802.111 (WiFi), PPP
- ❖ **physical**: bits “on the wire”



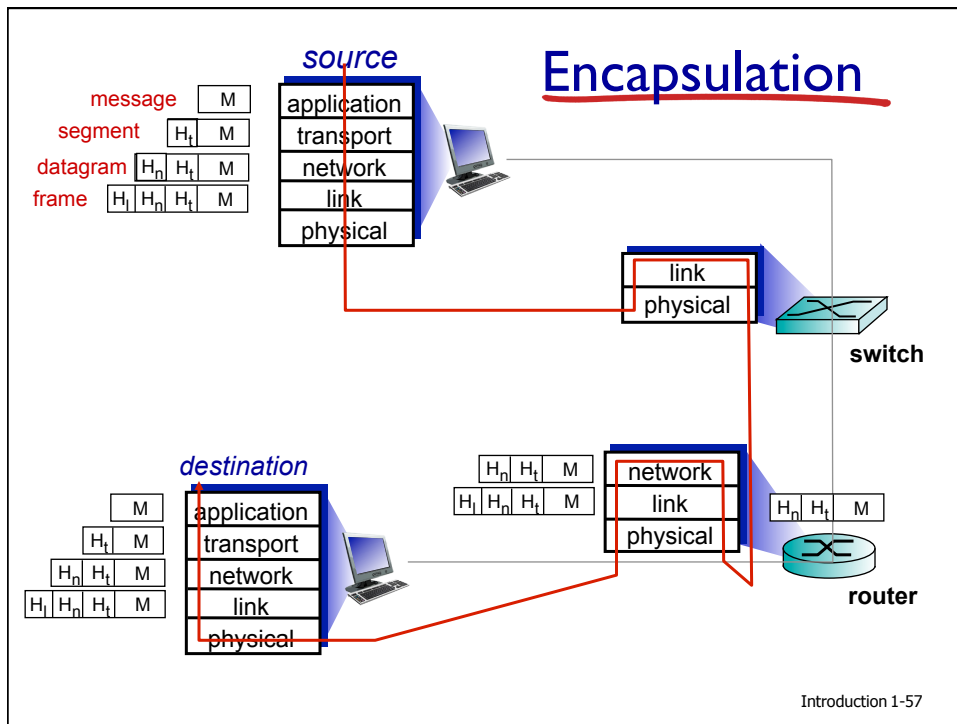
Introduction 1-55

ISO/OSI reference model

- ❖ **presentation**: allow applications to interpret meaning of data, e.g., encryption, compression, machine-specific conventions
- ❖ **session**: synchronization, checkpointing, recovery of data exchange
- ❖ Internet stack “missing” these layers!
 - these services, *if needed*, must be implemented in application
 - needed?



Introduction 1-56



Chapter I: roadmap

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- 1.7 history

Network security

- ❖ **field of network security:**
 - how bad guys can attack computer networks
 - how we can defend networks against attacks
 - how to design architectures that are immune to attacks
- ❖ **Internet not originally designed with (much) security in mind**
 - *original vision*: “a group of mutually trusting users attached to a transparent network” 😊
 - Internet protocol designers playing “catch-up”
 - security considerations in all layers!

Introduction 1-59

Bad guys: put malware into hosts via Internet

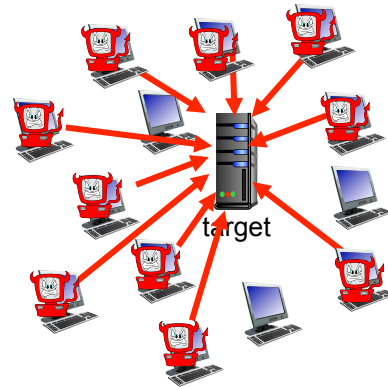
- ❖ malware can get in host from:
 - *trojan horse*: hidden in otherwise useful software
 - *virus*: self-replicating infection by executing object (e.g., e-mail attachment)
 - *worm*: self-replicating infection, actively seeks out other hosts to infect
- ❖ **spyware malware** can record keystrokes, web sites visited, upload info to collection site
- ❖ infected host can be enrolled in **botnet**, used for spam. DDoS attacks

Introduction 1-60

Bad guys: attack server, network infrastructure

Denial of Service (DoS): attackers make resources (server, bandwidth) unavailable to legitimate traffic by overwhelming resource with bogus traffic

1. select target
2. break into hosts around the network (see botnet)
3. send packets to target from compromised hosts

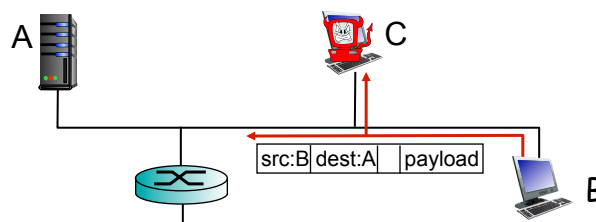


Introduction 1-61

Bad guys can sniff packets

packet “sniffing”:

- broadcast media (shared ethernet, wireless)
- promiscuous network interface reads/records all packets (e.g., including passwords!) passing by

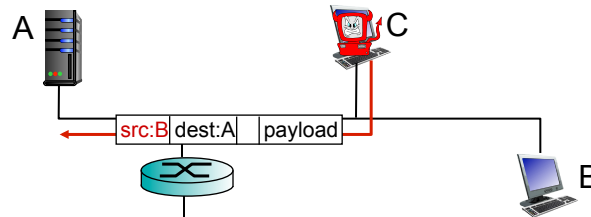


- ❖ wireshark software used for end-of-chapter labs is a (free) packet-sniffer

Introduction 1-62

Bad guys can use fake addresses

IP spoofing: send packet with false source address

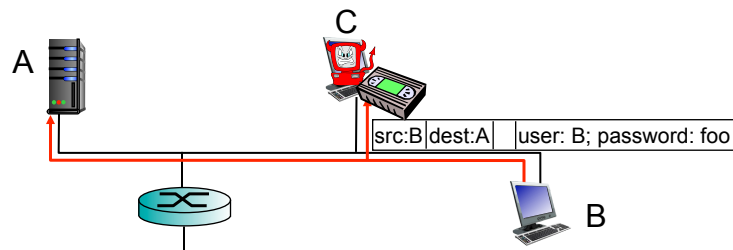


Introduction 1-63

Bad guys can record, playback

record-and-playback: sniff sensitive info (e.g., password), and use later

- password holder is that user from system point of view



... lots more on security (throughout, Chapter 8)

Introduction 1-64

Chapter 1: roadmap

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1.5 protocol layers, service models

1.6 networks under attack: security

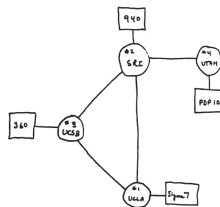
1.7 history

Introduction 1-65

Internet history

1961-1972: Early packet-switching principles

- ❖ 1961: Kleinrock - queueing theory shows effectiveness of packet-switching
- ❖ 1964: Baran - packet-switching in military nets
- ❖ 1967: ARPANet conceived by Advanced Research Projects Agency
- ❖ 1969: first ARPANet node operational
- ❖ 1972:
 - ARPANet public demo
 - NCP (Network Control Protocol) first host-host protocol
 - first e-mail program
 - ARPANet has 15 nodes



THE ARPA NETWORK

Introduction 1-66

Internet history

1972-1980: Internetworking, new and proprietary nets

- ❖ 1970: ALOHAnet satellite network in Hawaii
- ❖ 1974: Cerf and Kahn - architecture for interconnecting networks
- ❖ 1976: Ethernet at Xerox PARC
- ❖ late70's: proprietary architectures: DECnet, SNA, XNA
- ❖ late 70's: switching fixed length packets (ATM precursor)
- ❖ 1979: ARPAnet has 200 nodes

Cerf and Kahn's internetworking principles:

- minimalism, autonomy - no internal changes required to interconnect networks
- best effort service model
- stateless routers
- decentralized control

define today's Internet architecture

Introduction 1-67

Internet history

1980-1990: new protocols, a proliferation of networks

- ❖ 1983: deployment of TCP/IP
- ❖ 1982: smtp e-mail protocol defined
- ❖ 1983: DNS defined for name-to-IP-address translation
- ❖ 1985: ftp protocol defined
- ❖ 1988: TCP congestion control
- ❖ new national networks: Cernet, BITnet, NSFnet, Minitel
- ❖ 100,000 hosts connected to confederation of networks

Introduction 1-68

Internet history

1990, 2000's: commercialization, the Web, new apps

- ❖ early 1990's: ARPAnet decommissioned
- ❖ 1991: NSF lifts restrictions on commercial use of NSFnet (decommissioned, 1995)
- ❖ early 1990s: Web
 - hypertext [Bush 1945, Nelson 1960's]
 - HTML, HTTP: Berners-Lee
 - 1994: Mosaic, later Netscape
 - late 1990's: commercialization of the Web
- late 1990's – 2000's:
 - ❖ more killer apps: instant messaging, P2P file sharing
 - ❖ network security to forefront
 - ❖ est. 50 million host, 100 million+ users
 - ❖ backbone links running at Gbps

Introduction 1-69

Internet history

2010:

- ❖ ~750 million hosts
- ❖ voice, video over IP
- ❖ P2P applications: BitTorrent (file sharing) Skype (VoIP), PPLive (video)
- ❖ more applications: YouTube, gaming, Twitter
- ❖ wireless, mobility

Introduction 1-70

Introduction: summary

covered a “ton” of material!

- ❖ Internet overview
- ❖ what’s a protocol?
- ❖ network edge, core, access network
 - packet-switching versus circuit-switching
 - Internet structure
- ❖ performance: loss, delay, throughput
- ❖ layering, service models
- ❖ security
- ❖ history

you now have:

- ❖ context, overview, “feel” of networking
- ❖ more depth, detail to *follow!*