

# *Finite Automata*

*A finite automaton has a finite set of states with which it accepts or rejects strings.*

## A Finite Automaton

An FA has three components:

1. **input tape** contains single string;
2. **head** reads input string one symbol at a time;  
and
3. **Memory** is in one of a finite number of states.

## *Operating an FA*

### **Operating an FA.**

- 1) Set the machine to start state.*
- 2) If End-of-String then halt.*
- 3) Read a symbol.*
- 4) Update state according to current state and symbol read.*
- 5) Goto Step 2.*

## An FA Accepts Strings

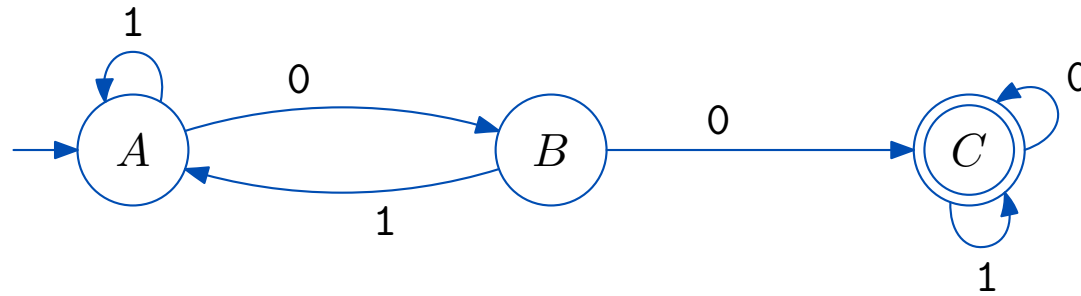
“Program” prescribes how symbols read affect current state.

**Final state** is state FA is in when finished reading the input string.

There are **accept** states (double circle) and **reject** states.

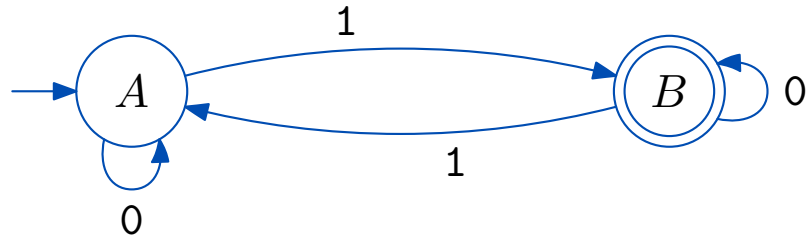
An FA **accepts** input string if final state is accept state; otherwise it rejects.

## An Example FA



Final state for **101001** is *C*, final state for **11101** is *A*.

## Example FA



Accepts all strings of 0's and 1's with odd number of 1's.

## Terminology

**alphabet** is a set of symbols (often denoted  $\Sigma$ )

**language** is a set of strings (unary language means  $|\Sigma| = 1$ )

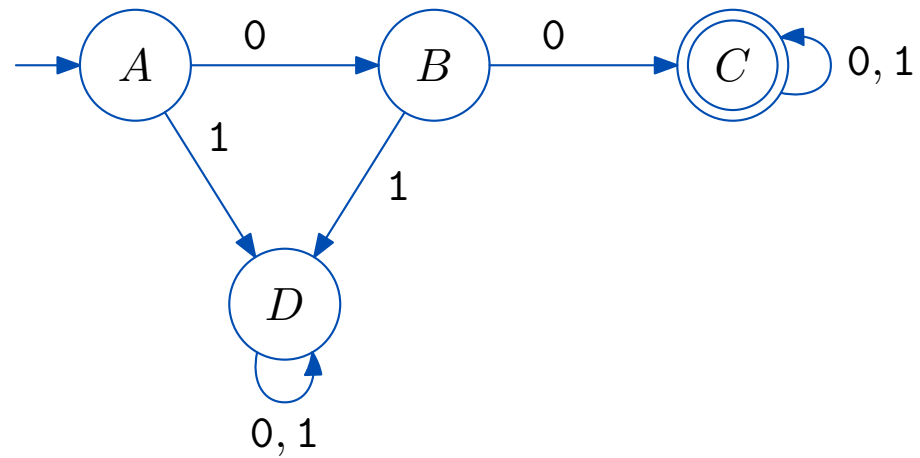
**language of FA** is the set of strings it accepts

**length** of a string is the number of symbols

**empty string** is denoted  $\varepsilon$ .

## *Building FAs: Do the Obvious*

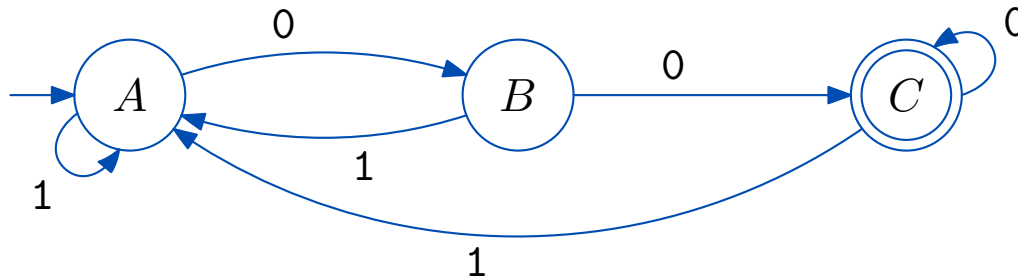
Starts with 00:





## *Building FAs: Recent Memory*

Ends with 00:

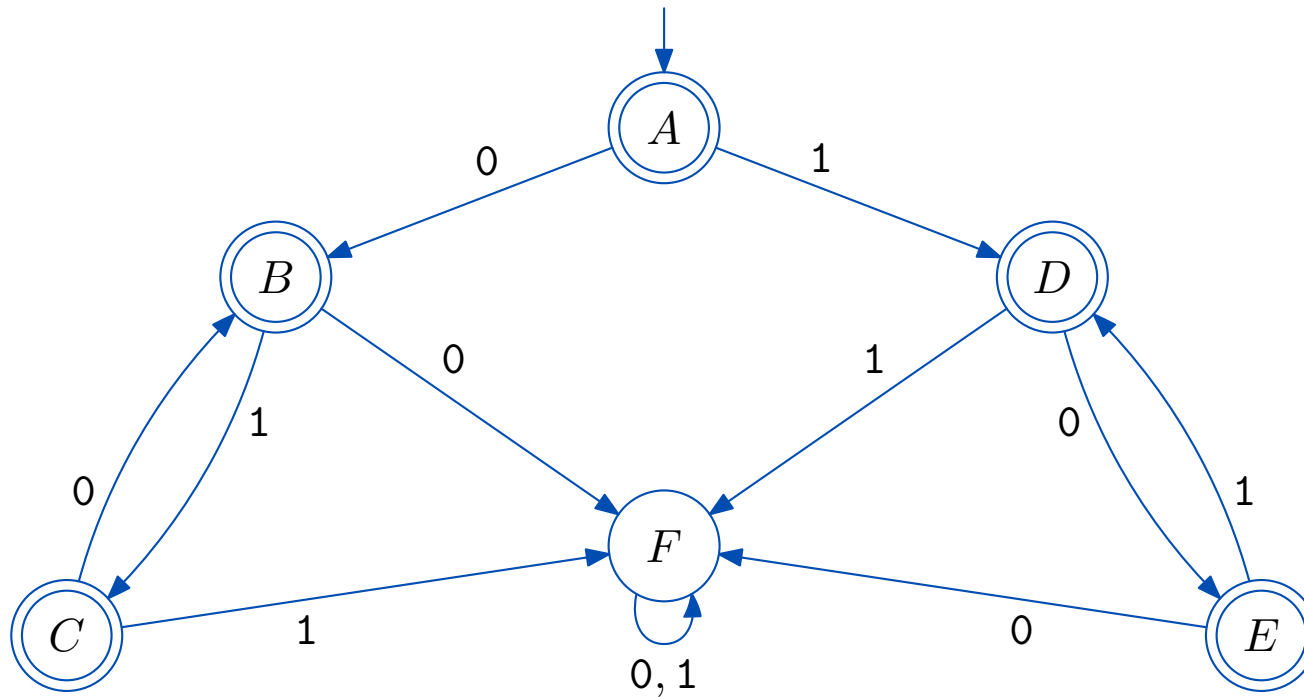


## *Building FAs: Traps*

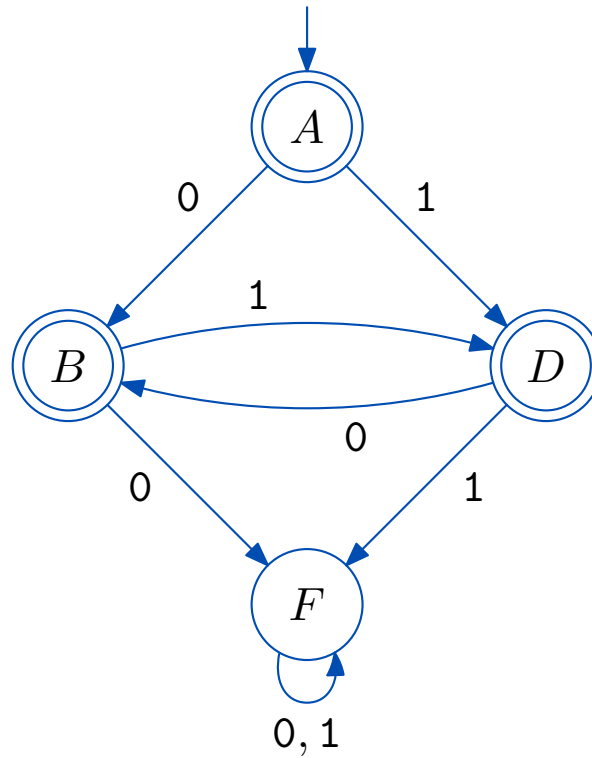
A ***trap*** is state that, once entered, one can never leave. Used to reject partly read strings that will never be accepted, or to accept partly read strings that will definitely be accepted.

## Example with a Trap

Alternating 0's and 1's:

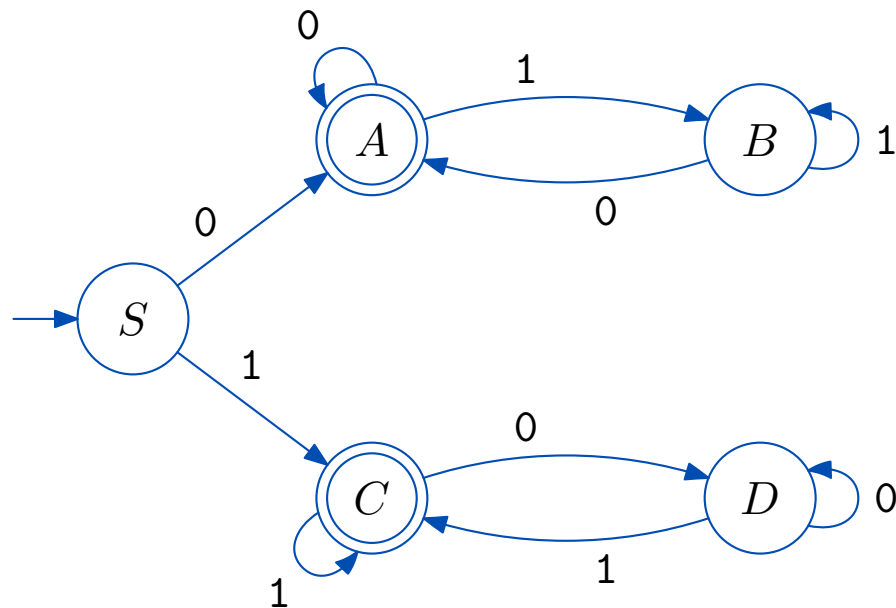


*Alternating 0's and 1's again*



## Building FAs: Permanent Memory

An FA remembers permanently by splitting into pieces:

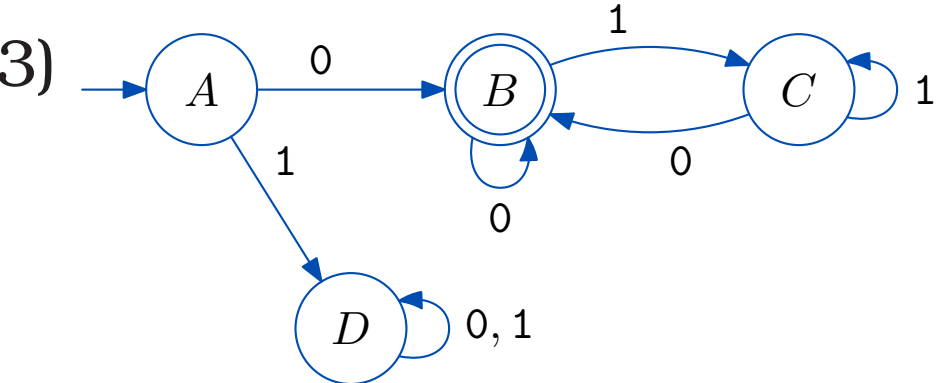
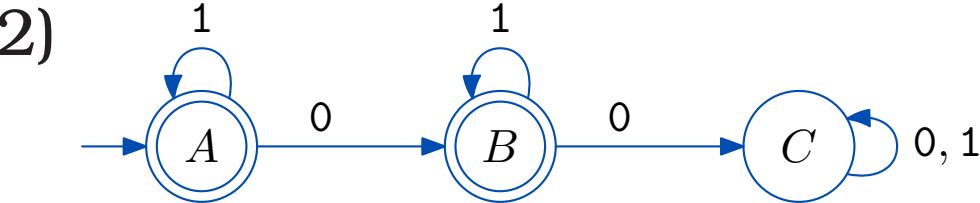
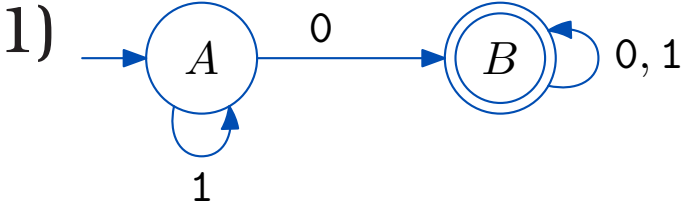


## *Practice*

Give FAs for each of the following three languages:

1. All binary strings with at least one 0
2. All binary strings with at most one 0
3. All binary strings starting and ending with 0  
(and single-0 string counts)

# Solutions to Practice



## Transition Table

A **transition table** is matrix that lists new state given current state and symbol read.

Here's transition table for FA for all binary strings that begin and end with same symbol.

		<i>Input</i>	
		0	1
<i>State</i>	S	A	C
	A	A	B
	B	A	B
	C	D	C
	D	D	C



## Formal Definition

A (deterministic) FA is 5-tuple

$(Q, \Sigma, q_0, T, \delta)$  where:

- $Q$  is finite set of states;
- $\Sigma$  is alphabet of input symbols;
- $q_0$  is start state;
- $T$  is subset of  $Q$  giving the accept states; and
- $\delta$  is **transition function** that maps state and symbol to state. (Mathematically,  $\delta: Q \times \Sigma \mapsto Q$ .)

## *Summary*

A finite automaton (FA) is a device that recognizes a language (set of strings). It has finite memory and an input tape; each input symbol that is read causes the machine to update its state based on its current state and the symbol read. The machine accepts the input if it is in an accept state at the end of the string; otherwise, the input is rejected.