Sample Mid-Term Exam 2

CS 5510/6510, Fall 2017

November 3

Name: ____

Instructions: You have eighty minutes to complete this open-book, open-note, closed-interpreter exam. Please write all answers in the provided space, plus the back of the exam if necessary.

Note on actual exam: The exam will refer to the lambda-k.rkt interpreter. If you need the interpreter for reference to answer the questions, please bring a copy (paper or electronic) with you.

- 1) [15 pts] Which of the following produce different results in a eager language and a lazy language? Both produce the same result if they both produce the same number or they both produce a procedure (even if the procedure doesn't behave exactly the same when applied), but they can differ in errors reported.
 - a) {{lambda {y} 12} {1 2}}
 - b) {lambda {x} {{lambda {y} 12} {1 2}}}
 - c) {+ 1 {lambda {y} 12}}
 - d) {+ 1 {{lambda {x} {+ 1 13}} {+ 1 {lambda {z} 12}}}
 - e) {+ 1 {{lambda {x} {+ x 13}} {+ 1 {lambda {z} 12}}}

2) [25 pts] Given the type rules

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$$\begin{split} [\dots \mathbf{x} \leftarrow \tau \dots] \vdash \mathbf{x} : \tau \quad \Gamma \vdash 1: \texttt{num} \quad \frac{\Gamma \vdash \mathbf{e}_1 : \texttt{num} \quad \Gamma \vdash \mathbf{e}_2 : \texttt{num}}{\Gamma \vdash \{\texttt{+} \mathbf{e}_1 \; \mathbf{e}_2\} : \texttt{num}} \\ \\ \frac{\Gamma[\mathbf{x} \leftarrow \tau_1] \vdash \mathbf{e} : \tau_2}{\Gamma \vdash \{\texttt{lambda} \; \{[\mathbf{x} \; : \; \tau_1]\} \; \mathbf{e}\} : (\tau_1 \to \tau_2)} \quad \frac{\Gamma \vdash \mathbf{e}_1 : (\tau_1 \to \tau_2) \quad \Gamma \vdash \mathbf{e}_2 : \tau_1}{\Gamma \vdash \{\mathbf{e}_1 \; \mathbf{e}_2\} : \tau_2} \end{split}$$

in one of the following expressions, the ____ can be filled in with a type so that the resulting expression has a type in the enmpty environment, while there is no type for the ____ that causes the other to have a type. Pick the right expression and show a derivation tree (which is a trace of typecheck that's written in the style as the type rules above) demonstrating that the chosen expression has a type.

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{{lambda {[x : ____]} {+ x 1}} x}
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{lambda {[x : ____]} {+ {x 1} 1}}
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Note that your answer should not include symbols like Γ , τ , or **e**, except when used as designated abbreviations, since those are meta-variables that are replaced by concrete environments, types, and expressions in the derivation tree.

3) [60 pts] Given the following expression:

{{lambda {x} {x x}} {lambda {y} 12}}

Describe a trace of the evaluation in terms of arguments to interp and continue functions for every call of each in the lambda-k.rkt interpreter. (There will be 7 calls to interp and 5 calls to continue.) The interp function takes three arguments — an expression, an environment, and a continuation — so show all three for each interp call. The continue function takes two arguments — a continuation and a value — so show both for each continue call. Represent continuations using records.

Answers

1) a and d.

 $\mathbf{2})$

$$\frac{ \begin{array}{c} \hline{\Gamma_1 \vdash \mathbf{x} : (\texttt{num} \to \texttt{num}) \quad \Gamma_1 \vdash \texttt{1} : \texttt{num}} \\ \hline{\Gamma_1 \vdash \{\texttt{x} \; \texttt{1}\} : \texttt{num}} \quad \Gamma_1 \vdash \texttt{1} : \texttt{num}} \\ \hline{\Gamma_1 = [\texttt{x} \leftarrow (\texttt{num} \to \texttt{num})] \vdash \{\texttt{+} \{\texttt{x} \; \texttt{1}\} \; \texttt{1}\} : \texttt{num}} \\ \hline{\emptyset \vdash \{\texttt{lambda} \{ [\texttt{x} : (\texttt{num} \to \texttt{num})) \} \} \{\texttt{+} \{\texttt{x} \; \texttt{1}\} \; \texttt{1}\} \} : ((\texttt{num} \to \texttt{num}) \to \texttt{num})} \end{array}}$$

 $\mathbf{3})$

interp expr =
$$\begin{bmatrix} \{ \{ ambda \{x\} \{x x\} \} \{ ambda \{y\} 12 \} \} \\ env = mt-env \\ k = (doneK) \\ \\ interp expr = \begin{bmatrix} \{ ambda \{x\} \{x x\} \} \} \\ env = mt-env \\ k = (appArgK \{ ambda \{y\} 12 \} mt-env (doneK)) = k_1 \\ \\ cont k = (appArgK \{ ambda \{y\} 12 \} mt-env (doneK)) or k_1 \\ val = (closV 'x \{x x\} mt-env) = v_1 \\ \\ interp expr = \{ ambda \{y\} 12 \} \\ env = mt-env \\ k = (doAppK v_1 (doneK)) = k_2 \\ \\ cont k = (doAppK v_1 (doneK)) or k_2 \\ val = (closV 'y 12 mt-env) = v_2 \\ \\ interp expr = \{ x x \} \\ env = (extend-env (bind 'x v_2) mt-env) = e_1 \\ k = (doneK) \\ \\ interp expr = e_1 \\ k = (appArgK x e_1 (doneK)) = k_3 \\ \\ cont k = (appArgK x e_1 (doneK)) or k_3 \\ val = v_2 \\ \\ interp expr = [x] \\ env = e_1 \\ k = (doAppK v_2 (doneK)) = k_4 \\ \\ cont k = (doAppK v_2 (doneK)) or k_4 \\ val = v_2 \\ \\ interp expr = [x] \\ env = e_1 \\ k = (doAppK v_2 (doneK)) or k_4 \\ val = v_2 \\ \\ interp expr = [12] \\ \\ \end{cases}$$

$$env = (extend-env (bind 'y v_2) mt-env)
k = (doneK)
cont k = (doneK)
val = (numV 12)$$